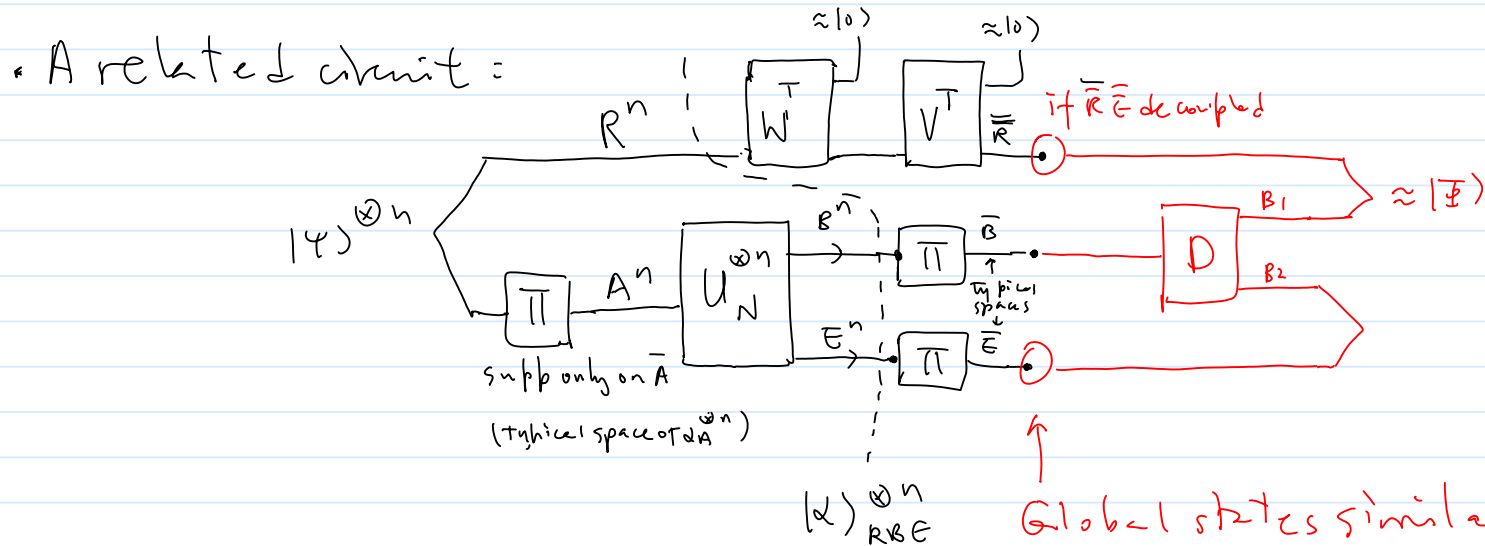


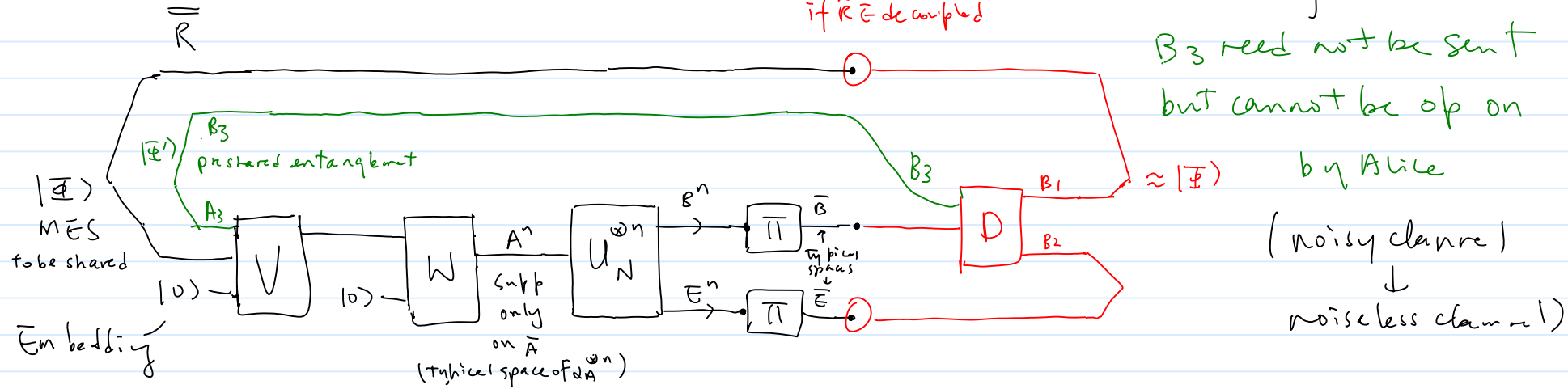
Note Title

- Recall the coding protocol for the LSD theorem:

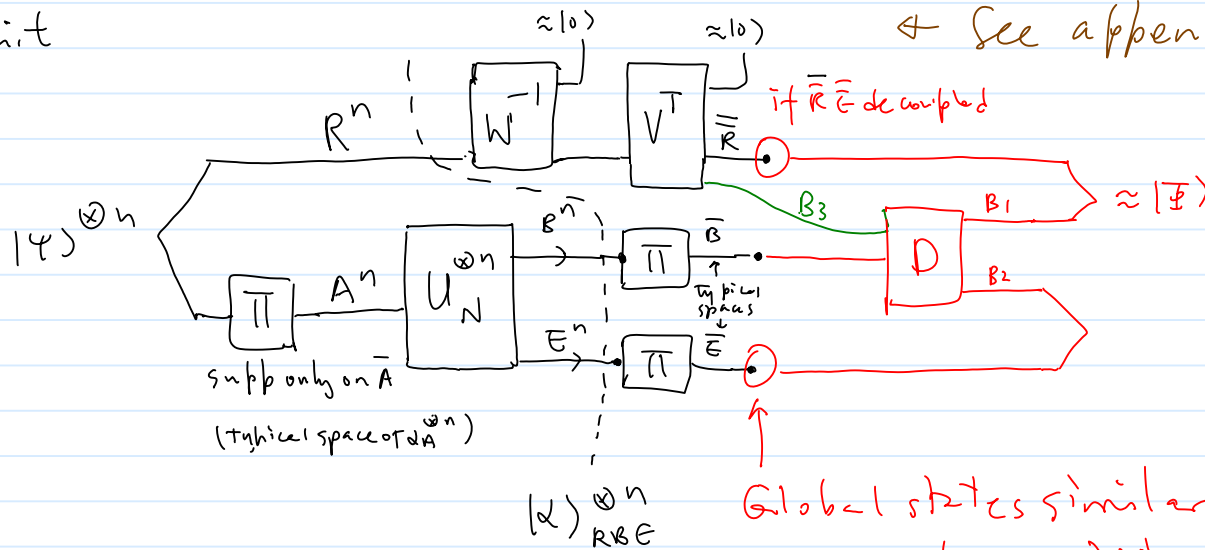


Global states similar on these 2 circuits
can analyze 2nd one

- The father protocol (direct coding): charged-entanglement assisted Q comm by Q channel



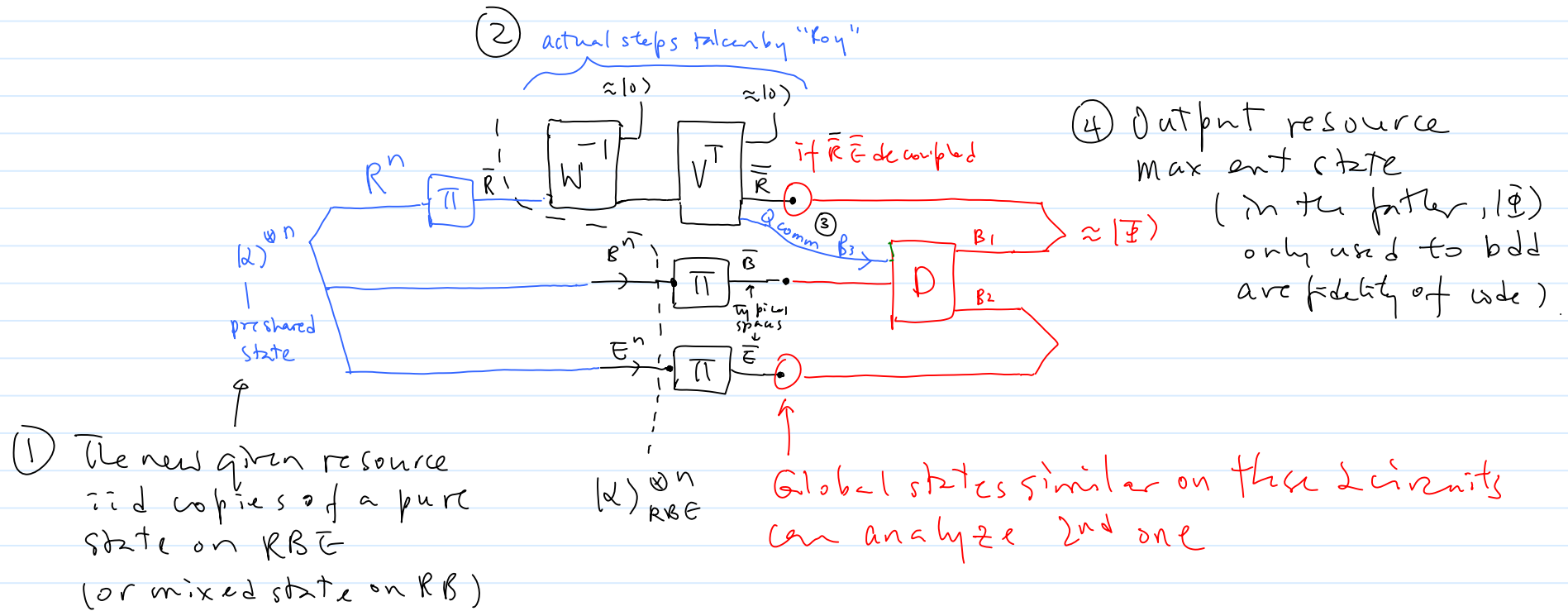
- A related circuit



[Choose basis for R
s.t. $W^T = W^\dagger$]

Global states similar on these 2 circuits can analyze 2nd one

- A protocol "mother" suggested by the related circuit:



Mother = Mixed state entanglement purification (noisy ent \rightarrow noiseless ent)
using charged α communication

Suffices to analyze mother (father follows with $|\alpha\rangle = I \otimes I \otimes |\Phi\rangle$
same dim for \bar{R} & B_3).

Goal = $\max \dim(\overline{R})$ (Q comm or
distillable entanglement obtained)

$\min \dim(B_3)$ (Noiseless ent or
 Q communication spent)

while decoupling $\overline{R} \in$ (to guarantee the job done right)

Recall approx decoupling lemma: (lecture 10 p 11)

$$\text{If } \| \rho_{\bar{R}\bar{E}} - \left(\frac{I}{2^{nr}} \right)_{\bar{R}} \otimes \rho_{\bar{E}} \|_{\text{tr}} \leq \epsilon'$$

then $\exists |\psi\rangle_{\bar{R}\bar{E}B_1B_2}$ purifying $\rho_{\bar{R}\bar{E}}$

$$\text{s.t. } \| \text{tr}_{\bar{E}B_2} |\psi\rangle\langle\psi|_{\bar{R}\bar{E}B_1B_2} - |\Phi\rangle\langle\Phi|_{\bar{R}B_1} \| \leq 2\sqrt{\epsilon'}$$

$$\text{So we bound } \underbrace{\| \rho_{\bar{R}\bar{E}} - \left(\frac{I}{2^{nr}} \right)_{\bar{R}} \otimes \rho_{\bar{E}} \|_{\text{tr}}}_M$$

Useful lemmas:

(L1) Cauchy-Schwarz inequality

$$\|M\|_{\text{tr}}^2 \leq \text{rank}(M) \text{tr}(M^\dagger M) = \text{rank}(M) \|M\|_2^2$$

(L2) $\text{tr}(M^2) = \text{tr}(\underbrace{\text{SWAP}} M \otimes M)$

the operator taking $|ij\rangle$ to $|ji\rangle$ eg $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$

p.f: $\text{tr}(M^2) = \sum_j \langle j | M I M | j \rangle$

Alternative
proof in
appendix 2

$$\begin{aligned} &= \sum_{i,j} \langle j | M | i \rangle \langle i | M | j \rangle \\ &= \sum_{i,j} \langle j | \langle i | M \otimes M | i \rangle | j \rangle \end{aligned}$$

$$= \text{tr} \left(\underbrace{\sum_{i,j} |i\rangle\langle j| \otimes |j\rangle\langle i|}_{\text{SWAP}} M \otimes M \right)$$

(L3) Let $S_1 = S_{11} S_{12} \dots S_{1s}$ be s 1-qubit systems
 $S_2 = S_{21} S_{22} \dots S_{2s}$ - - - - -

$$\text{Then } \text{SWAP}_{S_1 S_2} = \frac{1}{2^s} \sum_{P \in \mathcal{P}_s = 4^s \text{ Pauli matrices on } s \text{ qubits}} P \otimes P$$

$$\text{Pf: } \text{SWAP}_{S_{1i} S_{2i}} = \frac{1}{2} (I + XX + YY + ZZ) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\text{SWAP}_{S_1 S_2} = \bigotimes_{i=1}^s \text{SWAP}_{S_{1i} S_{2i}} = \frac{1}{2^s} \sum_{P \in \mathcal{P}_s} P \otimes P$$

state matrix extension of ρ_1 i.e. $\text{tr}_2 \rho_{12} = \rho_1$

$$(L4) \quad \text{Tr}(\overset{\text{state matrix}}{\rho_1} \overset{\text{extension of } \rho_1}{Q \otimes I_2}) = \text{Tr}(\rho_{12} Q \otimes I_2)$$

each with t qubits

$$(L5) \quad \mathbb{E} \left(V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \right) \left(I_{T_1 T_2} \otimes \text{SWAP}_{S_1 S_2} \right) \left(V_{T_1 S_1} \otimes V_{T_2 S_2} \right)$$

$V \in C_{s+t}$ = Clifford group on $s+t$ qubits

$$= \alpha I_{(T_1 S_1)(T_2 S_2)} + \beta \text{SWAP}_{(T_1 S_1)(T_2 S_2)}$$

$$\text{for } \alpha = 2^s \left[\frac{4^t - 1}{4^{s+t} - 1} \right] \leq \frac{1}{2^s} = \frac{1}{|S_1|}$$

$$\beta = 2^t \left[\frac{4^s - 1}{4^{s+t} - 1} \right] \leq \frac{1}{2^t} = \frac{1}{|T_1|}$$

$\mathcal{H} = B_{\gamma}(\mathcal{L}_3)$, LHS of (\mathcal{L}_5)

$$= \overline{\bigoplus_{V \in C_{s+t}}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\mathbb{I}_{T_1 T_2} \otimes \left(\sum_{P \in P_S} P_{S_1} \otimes P_{S_2} \right) \frac{1}{2^S} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$$

• If $P = \mathbb{I}$, $\overline{\bigoplus_{V \in C_{s+t}}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\mathbb{I}_{T_1 T_2} \otimes \mathbb{I}_{S_1 S_2} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$

$$= \mathbb{I}_{T_1 S_1 T_2 S_2}$$

• If $P \neq \mathbb{I}$, $\overline{\bigoplus_{V \in C_{s+t}}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[(\mathbb{I}_{T_1} P_{S_1}) \otimes (\mathbb{I}_{T_2} P_{S_2}) \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$

$$= \overline{\bigoplus_{V \in C_{s+t}}} V_{T_1 S_1}^\dagger (\mathbb{I}_{T_1} P_{S_1}) V_{T_1 S_1} \otimes V_{T_2 S_2}^\dagger (\mathbb{I}_{T_2} P_{S_2}) V_{T_2 S_2}$$

see appendix 3

The Clifford group permutes Pauli matrices by conjugation
 Δ acts transitively on all the non-identity Pauli's.

(for any $P_1 \neq I$, $P_2 \neq I$, $\exists V$ s.t. $VP_1V^\dagger = P_2$)

$$J = \frac{1}{\begin{matrix} s+t \\ 4 & -1 \end{matrix}} \sum_{Q \in P_{s+t}, Q \neq I} Q \otimes Q$$

of possible $Q = V^\dagger (I \otimes P) V \neq I$, $Q \in P_{s+t}$

$$\textcircled{L3} = \left(\frac{1}{\begin{matrix} s+t \\ 4 & -1 \end{matrix}} \right) \left(2^{s+t} \text{SWAP}_{(T_1 S_1)(T_2 S_2)} - I \otimes I \right)$$

So LHS of (LS)

$$= \mathbb{E}_{V \in C_{s+t}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\mathbb{I}_{T_1 T_2} \otimes \left(\sum_{P \neq P_S} P_{S_1} \otimes P_{S_2} \right) \frac{1}{2^s} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$$

$$= \frac{1}{2^s} (P = \mathbb{I} \text{ case}) + \frac{4^s - 1}{2^s} (P \neq \mathbb{I} \text{ case})$$

$$= \frac{1}{2^s} \mathbb{I}_{T_1 S_1} \mathbb{I}_{T_2 S_2} + \frac{4^s - 1}{2^s} \frac{1}{4^{s+t} - 1} \left(2^{s+t} \text{SWAP}_{(T_1 S_1)(T_2 S_2)} - \mathbb{I}_{T_1 S_1} \mathbb{I}_{T_2 S_2} \right)$$

$$= \frac{1}{2^s} \left(1 - \frac{4^s - 1}{4^{s+t} - 1} \right) \mathbb{I}_{T_1 S_1} \mathbb{I}_{T_2 S_2} + 2^t \frac{4^s - 1}{4^{s+t} - 1} \text{SWAP}_{(T_1 S_1)(T_2 S_2)}$$

$$\frac{1}{2^s} \left(\frac{4^{s+t} - 4^s}{4^{s+t} - 1} \right) = 2^s \left(\frac{4^t - 1}{4^{s+t} - 1} \right) = 2 \quad \beta$$

NB The average over Clifford group in (LS) same as average over V drawn over the Haar meas.

The above allows stabilizer codes (not random codes) to be used.

A slightly modified proof (appendix 3) shows that the Clifford group is a "2-design".

Back to $\| f_{\bar{R}\bar{E}} - \left(\frac{I}{2^{nr}}\right)_{\bar{R}} \otimes f_{\bar{E}} \|_2^2$

$$= \text{tr} \left[\left(f_{\bar{R}\bar{E}} - \left(\frac{I}{2^{nr}}\right)_{\bar{R}} \otimes f_{\bar{E}} \right) \left(f_{\bar{R}\bar{E}} - \left(\frac{I}{2^{nr}}\right)_{\bar{R}} \otimes f_{\bar{E}} \right) \right]$$

$$= \text{tr} \left(f_{\bar{R}\bar{E}}^2 \right) - 2 \text{tr} \left(f_{\bar{R}\bar{E}} \left(\frac{I}{2^{nr}}\right)_{\bar{R}} \otimes f_{\bar{E}} \right) + \text{tr} \left(\left(\frac{I}{2^{nr}}\right)_{\bar{R}}^2 \otimes f_{\bar{E}}^2 \right)$$

$$= \text{tr} \left(f_{\bar{R}\bar{E}}^2 \right) - 2 \left(\frac{1}{2^{nr}} \right) \underbrace{\text{tr} \left(f_{\bar{R}\bar{E}} \left(I_{\bar{R}} \otimes f_{\bar{E}} \right) \right)}_{\downarrow \textcircled{L4}} + \frac{1}{2^{nr}} \text{tr} f_{\bar{E}}^2$$

$$\text{tr} \left(\text{tr}_{\bar{R}} \left(f_{\bar{R}\bar{E}} \right) \cdot f_{\bar{E}} \right) = \text{tr} \left(f_{\bar{E}} \cdot f_{\bar{E}} \right)$$

$$= \text{tr} \left(f_{\bar{R}\bar{E}}^2 \right) - \frac{1}{2^{nr}} \text{tr} f_{\bar{E}}^2$$

Now bounding $\mathbb{E}_V \text{tr} (P_{\bar{R}\bar{E}})^2$

(L2)

$$= \mathbb{E}_V \text{tr} \left[\left(P_{\bar{R}_1 \bar{E}_1} \otimes P_{\bar{R}_2 \bar{E}_2} \right) \text{SWAP}_{(\bar{R}_1 \bar{E}_1) (\bar{R}_2 \bar{E}_2)} \right]$$

(L4)

$$= \mathbb{E}_V \text{tr} \left[\underbrace{\left(P_{B_{31} \bar{R}_1 \bar{E}_1} \otimes P_{B_{32} \bar{R}_2 \bar{E}_2} \right)}_{V_{B_{31} \bar{R}_1} \otimes I_{\bar{E}_1} (\alpha_{\bar{R}_1 \bar{E}_1}) V_{B_{31} \bar{R}_1}^\dagger \otimes I_{\bar{E}_1}} \left(I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right) \right]$$

$$V_{B_{31} \bar{R}_1} \otimes I_{\bar{E}_1} (\alpha_{\bar{R}_1 \bar{E}_1}) V_{B_{31} \bar{R}_1}^\dagger \otimes I_{\bar{E}_1}$$

similarly

$$\boxed{\bar{R}_1 = \bar{R}_1 B_3}$$

cyclic tr

$$= \mathbb{E}_V \text{tr} \left[\left(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \right) \times \left(V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger \otimes I_{\bar{E}_1 \bar{E}_2} \left(I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes I_{\bar{E}_1 \bar{E}_2} \right) \right]$$

$$= \text{tr} \left[(\mathcal{L}_{\bar{R}_1 \bar{E}_1} \otimes \mathcal{L}_{\bar{R}_2 \bar{E}_2}) \times \left[\frac{1}{V} V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger \otimes I_{\bar{E}_1 \bar{E}_2} (I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2}) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes I_{\bar{E}_1 \bar{E}_2} \right] \right]$$

$$= \text{tr} \left[(\mathcal{L}_{\bar{R}_1 \bar{E}_1} \otimes \mathcal{L}_{\bar{R}_2 \bar{E}_2}) \times \left[\frac{1}{V} V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger (I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2}) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right] \right]$$

(L5)

$$\leq \text{tr} \left[(\mathcal{L}_{\bar{R}_1 \bar{E}_1} \otimes \mathcal{L}_{\bar{R}_2 \bar{E}_2}) \times \left[\left(\frac{1}{|\bar{R}|} I_{B_{31} \bar{R}_1 B_{32} \bar{R}_2} + \frac{1}{|B_3|} \text{SWAP}_{(B_{31} \bar{R}_1)(B_{32} \bar{R}_2)} \right) \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right] \right]$$

$$= \frac{1}{|\bar{R}|} \text{tr} (\mathcal{L}_{\bar{E}_1} \otimes \mathcal{L}_{\bar{E}_2} \text{SWAP}_{\bar{E}_1 \bar{E}_2}) + \frac{1}{|B_3|} \text{tr} (\mathcal{L}_{\bar{R}_1 \bar{E}_1} \otimes \mathcal{L}_{\bar{R}_2 \bar{E}_2} \text{SWAP}_{\bar{R}_1 \bar{E}_1 \bar{R}_2 \bar{E}_2})$$

$$= \frac{1}{|\vec{R}|} \text{tr}(\alpha \vec{E}^2) + \frac{1}{|B_3|} \text{tr}(\alpha \vec{R} \vec{E}^2)$$

$$\left\| \rho_{\vec{R}\vec{E}} - \left(\frac{\mathbb{I}}{2^{nr}} \right)_{\vec{R}} \otimes \rho_{\vec{E}} \right\|_2^2 \leq \underbrace{\left(\frac{1}{|\vec{R}|} \text{tr}(\alpha \vec{E}^2) + \frac{1}{|B_3|} \text{tr}(\alpha \vec{R} \vec{E}^2) \right)}_{\text{}} - \frac{1}{2^{nr}} \text{tr} \rho_{\vec{E}}^2$$

$$= \frac{1}{|B_3|} \text{tr}(\alpha \vec{R} \vec{E}^2)$$

$$\text{By } \textcircled{1}, \quad \left\| \rho_{\vec{R}\vec{E}} - \left(\frac{\mathbb{I}}{2^{nr}} \right)_{\vec{R}} \otimes \rho_{\vec{E}} \right\|_{\text{tr}}^2 \leq \frac{|\vec{R} \vec{E}|}{|B_3|} \underbrace{\text{tr}(\alpha \vec{R} \vec{E}^2)}_{\text{}}$$

$\alpha \vec{R} \vec{E} = (\alpha \vec{R} \vec{E})^{\otimes n}$ projected onto
typical space of $\rho_{\vec{E}}$. Each eig value
 $\approx \frac{1}{2^{nS(\rho_{\vec{E}})}}$ $\therefore \text{tr}(\alpha \vec{R} \vec{E}^2) = \frac{1}{2^{nS(\rho_{\vec{E}})}}$

If we choose $|\bar{R}| = 2^{\frac{n}{2} [S(R=B)_2 - \delta]}$

$|B_3| = 2^{\frac{n}{2} [S(R:Z)_2 + \delta]}$

then $\| \cdot \|_{tr} \leq \frac{2^{\frac{n}{2} [S(R=B)_2 - \delta]}}{2^{\frac{n}{2} [S(R:Z)_2 + \delta]}} \times \frac{2^{n(S(Z)_2 + \epsilon)}}{2^{n(S(R:Z)_2 - \epsilon)}}$

$= 2^{\frac{n}{2} [(S(R) + S(B) - S(R:Z))_2 - (S(R) + S(Z) - S(R:Z))_2 - 2\delta]} 2^{n(S(Z)_2 - S(R:Z)_2 + \epsilon)}$

$= 2^{n(S(B)_2 - S(R:Z)_2 - \delta)} 2^{n(S(Z)_2 - S(R:Z)_2 + 2\epsilon)}$

$= 2^{-n(\delta - 2\epsilon)}$

choose $\delta = 3\epsilon$

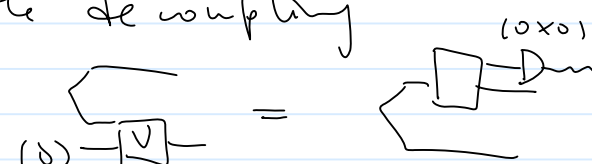
↑
fixed by how good
the typical spaces are

$= 2^{-n\epsilon} \rightarrow 0 \text{ as } n \rightarrow \infty$

NB the unassisted case (the LSP thm) has $|B_3| = 1$

$$\begin{aligned} \text{We choose } |\bar{R}| &= \sum \frac{n}{2} (S(R=B)_\alpha - S(R=E)_\alpha - 2\delta) \\ &= \sum n [I_c(R > B) - \delta] \end{aligned}$$

and the same decoupling condition for $\bar{R} \bar{E}$ holds.

The mapping from the direct coding scheme to the decoupling condition based on $\alpha^{\otimes n}$ has a glitch: 

For other ways out wires, decoupling still works but the proof is involved (see 0702005 p9-10 ptof Thm IV), & omitted

Alt: just use the father to get LSP.

① mother:

$$n \{ \{ \} \} + \frac{n}{2} (S(R=\bar{E}) + \delta) \underbrace{[q \rightarrow f]}_{\text{q bit}} \geq \frac{n}{2} (S(R=B) - \delta) \underbrace{[ff]}_{\text{ebits}}$$

noisy \nearrow static 2-party quantum resource ($\downarrow R=B$ here)

evaluated on α \uparrow the more \bar{E} has the more assistance it takes to decompile her

noiseless dynamical \swarrow

Bob's full potential \uparrow

② father:

$$n \{ \{ q \rightarrow f \} \} + \frac{n}{2} (S(R=\bar{E}) + \delta) \underbrace{[ff]}_{\text{ebits}} \geq \frac{n}{2} (S(R=B) - \delta) \underbrace{[q \rightarrow f]}_{\text{q bit}}$$

noisy dynamical resource (N here)

eval on $I \otimes U_N | \phi \rangle = | \alpha \rangle$

Note that we have asymptotic approximate resource inequalities here. Say, $XXX \geqslant YYY$. We demand the output resource YYY (lesser side) to be close to the ideal resource in trace distance or diamond norm. This ensures the protocol underlying the resource inequality can be used as a subroutine in any other protocol to produce YYY (using XXX) and when YYY is consumed, it is basically as good as ideal.

Say, $XXX + ZZZ \geqslant YYY + ZZZ \geqslant KKK$

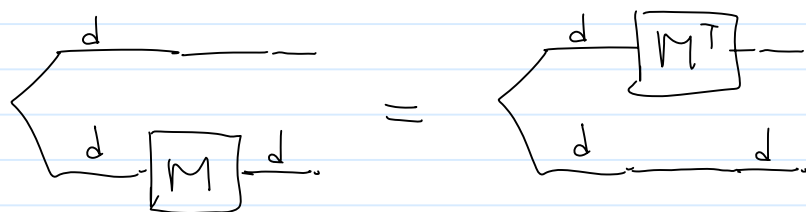
The first inequality only holds if $XXX \geqslant YYY$ is given by a protocol producing sufficiently

Appendix 1:

Let \sum_d denote $\sum_{i=1}^d |i\rangle\langle i| = \sqrt{d} \times \text{max ent state}$

The well-known transpose trick says the following:

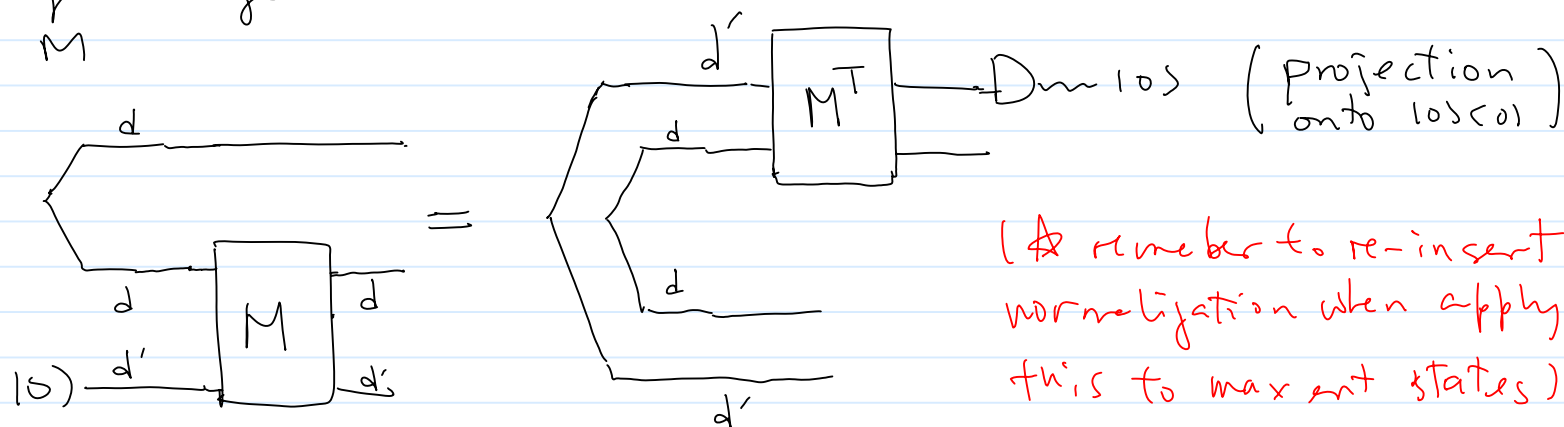
$\forall d \times d$ matrices M



(This holds also for MES)

This can be generalized:

$\forall d' \times d' \times d \times d$ M



(Remember to re-insert the normalization when applying this to max ent states)

Pf The output of the LHS

$$= I \otimes V \left(\sum_i |i\rangle |i\rangle \right) |0\rangle$$

$$\left[\text{Let } V|i\rangle = \sum_{rs} V_{ij} |rs\rangle \right]$$

$$= \sum_i |i\rangle \sum_{rs} V_{i0} |rs\rangle$$

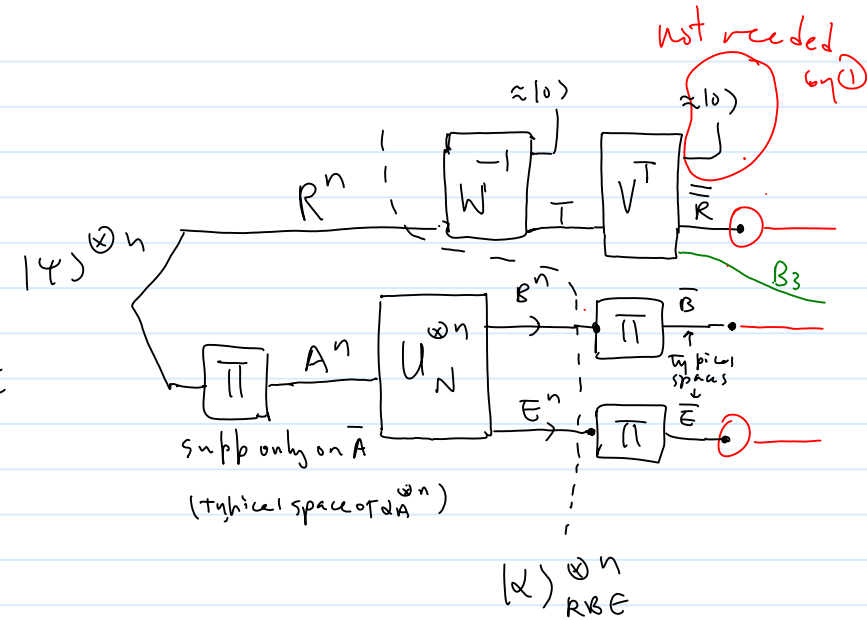
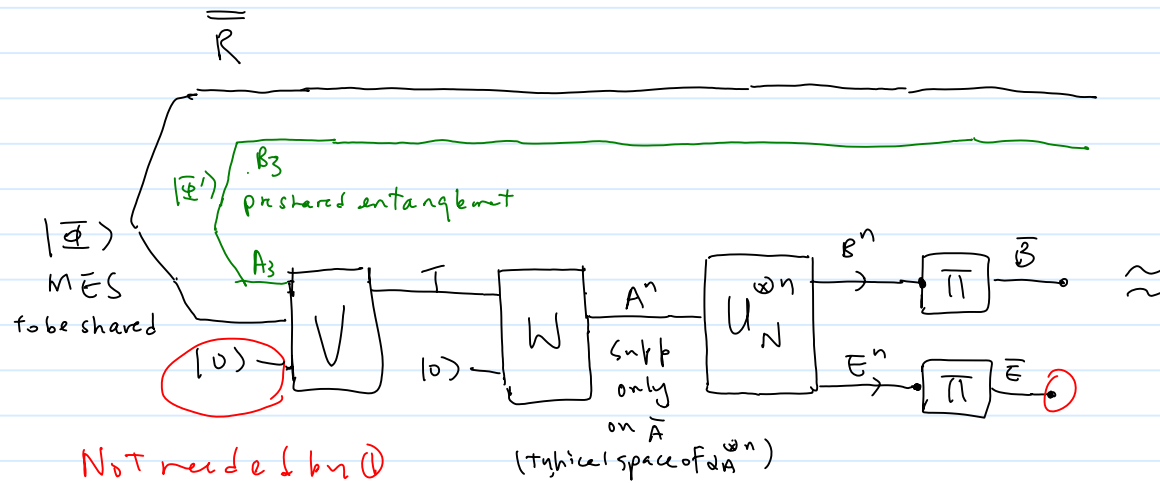
$$= \sum_{rs} \left(\sum_i V_{i0} |i\rangle \right) |rs\rangle$$

$$\left[\because V^T |rs\rangle = \sum_i V_{ij} |i\rangle \right]$$

$$= \sum_{rs} \left(I \otimes \langle 0| \right) \left(V^T |rs\rangle \right) \otimes |rs\rangle$$

$$= \left(I \otimes \langle 0| \otimes I \right) \left(V^T \otimes I \right) \sum_{rs} |rs\rangle |rs\rangle = \text{output of RHS}$$

for our purpose, ie to show



we don't need the above generalization of the transpose trick.

$$\textcircled{1} \text{ In the end, } \dim(\bar{R}) = 2^{\frac{n}{2} [S(R:B)_\alpha - 1]}$$

$$\dim(B_3) = 2^{\frac{n}{2} [S(R:E)_\alpha + 1]}$$

So $\dim(\bar{R} B_3) = \dim(T) = 2^{n(S(R) + \epsilon)}$ and $\begin{matrix} \bar{R} \\ |10\rangle \end{matrix} \boxed{V} = \boxed{V}$

(if necessary add ϵ terms to $\dim(B_3)$)

(For the unassisted case (the LSO then with $|B_3| = 0$)

moving V to the upper register requires an actual proj on the upper register. We can have a meas affecting one out of many possible projections.

Thm 14 of 0702005 is needed to ensure any outcome gives the same decoupling condition.)

② To move W : note initial state on RHS $\approx M E S$ on $2^{n S(K_L)}$ dims not $2^{n d_R}$ dims. So moving W to W^T only requires the original transpose trick.

Appendix 2:

Alt proof for L2: $\text{tr}(M^2) = \text{tr}(\text{SWAP } M \otimes M)$

$$\text{Pf: } \text{tr}(\text{SWAP } M \otimes M)$$

$$= \text{tr}(\text{SWAP } M \otimes I \cdot I \otimes M)$$

$$= \text{tr}(\underbrace{\text{SWAP } M \otimes I \text{ SWAP } I \otimes M}_{\text{I} \otimes M \text{ SWAP } I \otimes M})$$

$$= \text{tr}(\text{I} \otimes M \text{ SWAP } I \otimes M)$$

$$= \text{tr}(\text{SWAP } I \otimes M^2) \stackrel{\text{L4}}{=} \text{tr}\left[\underbrace{(\text{tr, SWAP})}_{\frac{1}{I}} M^2\right]$$

$$= \text{tr}(M^2)$$

$\frac{1}{I}$ (see next page)

NB $\text{Tr}_c(\text{SWAP}) = \mathbb{I}$

✓ Kronecker δ -fcn.

Df: $\text{SWAP} = \sum_{i,j} |j\rangle\langle i|$, $\text{tr}(|j\rangle\langle i|) = \delta_{ij}$

$\therefore \text{Tr}_c(\text{SWAP}) = \sum_{i,j} \delta_{ij} |i\rangle\langle j| = \mathbb{I}$ (on 2^n sys)

Appendix 3:

$$\text{Let } \gamma(M) = \sum_{V \in C_{s+t}} V \otimes V M V^\dagger \otimes V^\dagger$$

$V \in C_{s+t}$ (Weyl group on $s+t$ qubits)

① if $P \neq I$, $P \in P_{s+t}$ (Pauli group on $s+t$ qubits)

$$\text{then } \gamma(P \otimes P) = \frac{1}{4^{s+t} - 1} \sum_{\substack{Q \in P_{s+t} \\ Q \neq I}} Q \otimes Q$$

Pf: Note that C_{s+t} is transitive on $P_{s+t} - \{I\}$

ie for any $Q_1, Q_2 \in P_{s+t} - \{I\}$

$$\exists W \in C_{s+t} \text{ s.t. } W Q_1 W^\dagger = Q_2$$

Also, $\forall V \in C_{s+t}, V(P_{s+t} - \{I\})V^T$ only permutes elements of $P_{s+t} - \{I\}$.

$$\text{So } \tau(P \otimes P) = \sum_{Q \in P_{s+t} - \{I\}} \mu(Q) Q \otimes Q \text{ for some distribution } \mu(Q)$$

If $\mu(Q)$ not uniform, then $\exists Q_1, Q_2$ s.t. $\mu(Q_1) \neq \mu(Q_2)$

$$\text{let } W' Q_1 W'^T = Q_2$$

$$\text{then } \tau(P \otimes P) = W' \tau(P \otimes P) W'^T \left(\begin{array}{c} \text{this merely changes} \\ \sum_{V \in C_{s+t}} \text{ to } \sum_{WV \in C_{s+t}} \end{array} \right)$$

$$\left. \begin{array}{l} Q_2 \otimes Q_2 \text{ has weight } \mu(Q_2) \text{ on LHS} \\ Q_2 \otimes Q_2 = \dots \mu(Q_1) \text{ on RHS} \end{array} \right\} \otimes$$

($\because W'$ is a permutation on $P_{s+t} - \{I\}$, the $Q_2 \otimes Q_2$ term in the

2nd line can only come from the $Q \otimes Q$ term before conjugation by W .)

But $\{Q \otimes Q\}$ is trace orthonormal.

$\therefore (*)$ is a contradiction $\therefore \mu(Q)$ has to be uniform

$$(2) \quad \tau(P \otimes Q) = 0 \quad \forall P \neq Q, \quad P, Q \in P_{s+t}$$

Pf. WLOG $P \neq I$ (at least one of $P, Q \neq I$)

There are 4^{s+t-2} Pauli's anti commuting with P
& commuting with Q

Let R be one of them

$$\tau(P \otimes Q) = \overline{\mathbb{E}}_{V \in C_{s+t}} V \otimes V P \otimes Q V^* \otimes V^T$$

$$= \frac{1}{2} \left[\overline{\mathbb{E}}_{V \in C_{s+t}} V \otimes V P \otimes Q V^* \otimes V^T + \overline{\mathbb{E}}_{V \in C_{s+t}} VR \otimes VR P \otimes Q (VR)^* \otimes (VR)^T \right]$$

$$= 0.$$

So $\forall M$, $\tau(M)$ = linear combination of II & $SWAP$.

It's easy to show that the weights are same as that of

$$\int dU U \otimes U M U^* \otimes U^T \quad (\text{average } U \text{ over Haar meas.})$$

(see Q. data hiding paper DiVincenzo, L. Terhal for detail.)