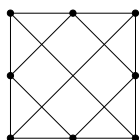


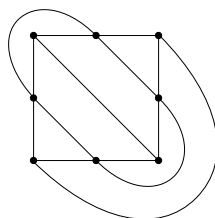
# MATH 239 Intro to Combinatorics, Solutions to Assignment 6

1: Determine which of the following graphs are planar.

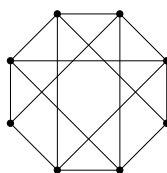
(a)



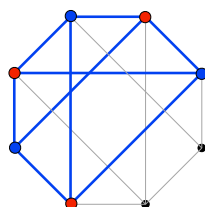
Solution: This graph is planar since we can redraw it as shown below.



(b)

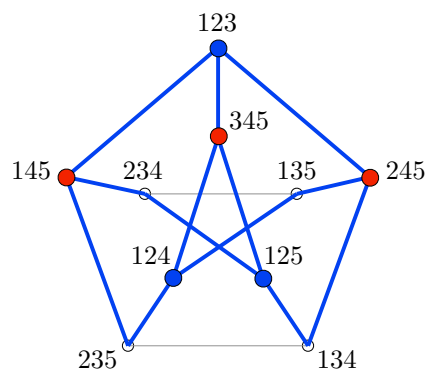


Solution: This graph is not planar since it contains a subgraph, as shown below, which is isomorphic to  $K_{3,3}$ .



(c)  $S_{5,3,1}$

Solution: Note that  $S_{5,3,1}$  is isomorphic to  $S_{5,2,0}$  (the Peterson graph) under the map which sends the 3-element subset to its 2-element complement. This observation helps to draw a symmetric picture of the graph, as shown below. We see that it is not planar since it contains a subgraph, as indicated in the picture, which is isomorphic to an edge subdivision of  $K_{3,3}$ .

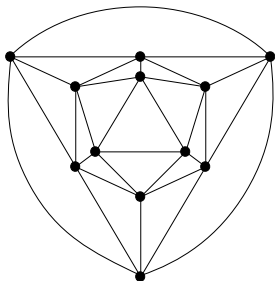


- 2: (a) Let  $G$  be a graph in which  $\deg(v) \geq 5$  for every vertex  $v$ . Suppose  $G$  has a planar embedding such that  $\deg(f) \geq 3$  for every face  $f$ . Show that  $G$  has at least 12 vertices. Find a planar embedding of such a graph with exactly 12 vertices.

Solution: We have  $2q = \sum_{v \in V} \deg(v) \geq 5p$  and we have

$$2q = \sum_{f \in F} \deg(f) \geq 3r = 3(q - p + c + 1) = 3q - 3p + 3(c + 1) \geq 3q - 3p + 6$$

so that  $q \leq 3p - 6$ . Combining these two inequalities gives  $5p \leq 2q \leq 6p - 12$  and so  $p \geq 12$ , as required. The picture below shows a planar embedding of the regular icosahedron, which satisfies  $p = 12$ ,  $\deg(v) = 5$  for all vertices  $v$ , and  $\deg(f) = 3$  for all faces  $f$ .



- (b) Let  $G$  be a connected graph which has a planar embedding in which  $\deg(f) \geq 4$  for every face  $f$ . Show that  $G$  has at least 3 vertices of degree at most 3. Find a planar embedding of such a graph with exactly 3 vertices of degree at most 3.

Solution: Note that  $2q = \sum_{f \in F} \deg(f) \geq 4r$ , so we have  $q \geq 2r = 2q - 2p + 2(c + 1) = 2q - 2p + 4$  and hence

$$q \leq 2p - 4 \quad (*)$$

If  $G$  had no vertices of degree at most 3, so that  $\deg(v) \geq 4$  for all  $v$ , then we would have  $2q = \sum \deg(v) \geq 4p$  so that  $q \geq 2p$ . But from (\*) this would imply that  $2p \leq q \leq 2p - 4$  so that  $0 \leq -4$ , which is not true. If  $G$  had exactly one vertex, say  $u$ , of degree at most 3, so that  $\deg(v) \geq 4$  for all  $v \neq u$ , then we would have  $2q = \deg(u) + \sum_{v \neq u} \deg(v) \geq 1 + 4(p - 1) = 4p - 3$ . But from (\*) this would imply that  $4p - 3 \leq 2q \leq 4p - 8$  so that  $-3 \leq -8$ , which is false. If  $G$  had exactly 2 vertices, say  $u_1$  and  $u_2$ , of degree at most 3, so that  $\deg(v) \geq 4$  for all  $v \neq u_1, u_2$ , then we would have  $2q = \deg(u_1) + \deg(u_2) + \sum_{v \neq u_1, u_2} \deg(v) \geq 1 + 1 + 4(p - 2) = 4p - 6$ .

But from (8) this would imply that  $4p - 6 \leq 2q \leq 4p - 8$  so that  $-6 \leq -8$ , which is not true. Thus  $G$  must have at least 3 vertices of degree at most 3. The (unique) tree on 3 vertices, embedded in the plane as shown below, has two vertices of degree 1 and one of degree 2, and it has one face, which is of degree 4.

