

# Boole's Equational Treatment of Particular Statements

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## Obstacles to Understanding Boole's Algebra of Logic

- Boole was using **ordinary algebra** with **idempotent symbols** ( $X^2 = X$ )

He was not using modern Boolean algebra nor Boolean rings.

His work was a blend of partial semantics with formal algebraic manipulation.

0	Empty class
1	Universe
$XY$	Intersection of $X$ and $Y$
$X + Y$	Union of $X$ and $Y$ if disjoint
$X - Y$	Difference of $X$ and $Y$ if $Y \subseteq X$

Thus  $1 - X$  is the **complement** of  $X$ .

- **Modern semantics** (for categorical propositions) was **not available** at the time of Boole's work.

Consider the following passage from  
De Morgan's *Formal Logic*, 1847, pp. 110:

On looking into any writer on logic,  
we shall see that *existence* is claimed  
for the significations of all the names.  
Never, in the statement of a  
proposition, do we find any reason left  
for the alternative, *suppose there  
should be no such thing*.

Regarding the meaning of 'All  $X$  is  $Y$ '  
De Morgan goes on to say,

If neither  $X$  nor  $Y$  exist, I will not  
... attempt to settle what nonexistent  
things agree or disagree.

Modern semantics was explicitly introduced  
by Peirce in 1880 when he stated that 'All  $X$   
is  $Y$ ' holds when  $X$  does not exist.

Peirce notes that the rule of limitation fails in  
his system.

Peirce's semantics was adopted by Schröder  
in his famous 3 volumes: *Algebra der Logik*.

# Boole’s Translation of Categorical Statements into Equations

## Affirmative Propositions

	1847	1848/1854
<b>A</b> All $X$ is $Y$	$X = XY$	$X = VY$
<b>I</b> Some $X$ is $Y$	$V = XY$	$VX = VY$

For the Negative Propositions just replace  $Y$  by  $1 - Y$  in the above.

	1847	1848/1854
<b>E</b> No $X$ is $Y$	$X = X(1 - Y)$	$X = V(1 - Y)$
<b>O</b> Some $X$ is not $Y$	$V = X(1 - Y)$	$VX = V(1 - Y)$

**EXAMPLE:** (Conversion by Limitation)

$$\frac{\text{All } X \text{ is } Y}{\text{Some } Y \text{ is } X.}$$

### **Boole's 1847 Derivation**

Given:  $X = XY$  (equivalently,  $X(1 - Y) = 0$ ).

Solve this for  $X$  to obtain  $X = VY$ .

[Verify this by substituting into the previous equation  $X(1 - Y) = 0$ .]

### **Boole's 1854 Derivation**

Given:  $X = VY$ .

Multiply both sides by  $V$  to obtain  $VX = VVY$ .

Use the idempotence of  $V$  to obtain  $VX = VY$ .

# The Theme of This Talk

Boole's use of **equations** for the **particular** propositions

	1847	1848/1854
<b>I</b>	$V = XY$	$VX = VY$
<b>O</b>	$V = X(1 - Y)$	$VX = V(1 - Y)$

has been unjustly dismissed by logicians,  
starting with the rejection by C.S. Peirce.

Boole's system can be easily patched to give  
a robust system (that will be called  $BF_3$  ).

# Boole Freed From every Flaw

The system  $BF_3$  :

<b>A</b>	All $X$ is $Y$	$X = XY$
<b>I</b>	Some $X$ is $Y$	$V = VXY$

For the negative propositions replace  $Y$  by  $1 - Y$ .

The choice for **A** is the same as Boole 1847.

The choice for **I** lies between Boole 1847 and Boole 1848/1854 since:

$$\begin{array}{lll} V = XY & \text{implies} & V = VXY \\ & \text{implies} & VX = VY. \end{array}$$



Kneale and Kneale, *The Development of Logic*, 1962, p. 411

Hailperin, *Boole's Logic and Probability*, 1986, p. 112

**OBJECTION:** Boole has special rules for interpreting  $V$  as 'some', for example:

Some  $X$  is not  $Y$                        $VX = V(1 - Y)$

algebraically leads to

Some  $Y$  is not  $X$                        $VY = V(1 - X)$

$VY$  is not to be interpreted as 'some  $Y$ '.

**ANSWER:** In  $BF_3$  there are no special rules limiting when one can interpret the symbols as 'some'.

In  $\text{BF}_3$  one has

Some  $X$  is not  $Y$        $V = VX(1 - Y)$

does not lead to

Some  $Y$  is not  $X$        $V = VY(1 - X)$

Furthermore we note that in  $\text{BF}_3$  any letter can play the role of  $V$ .

**EXAMPLE:** (Limitation)

All  $X$  is  $Y$        $X = XY$

leads to

Some  $X$  is  $Y$        $X = X(XY)$

C.S. Peirce, *On the algebra of logic*, 1880,  
[p. 22]

**OBJECTION:** ‘Some  $X$  is not  $Y$ ’ yields, by  
Boole’s algebraic system, ‘Some  $Y$  is not  $X$ ’.

**ANSWER:** (See previous two slides):

(1) Peirce does not realize that Boole has  
restrictions on when one can interpret  $V$  as  
‘some’.

In Peirce’s defence: In *Laws* Boole does not  
formulate his restrictions on using  $V$  as ‘some’  
until Chapter XV, the last chapter on logic.

(2) Cannot make this inference in  $BF_3$  .

E. Schröder, *Algebra der Logik Vol. II*, 1891,  
[p. 92]

**OBJECTION:** He gives a rigorous proof that one cannot use a system of equations to express ‘Some  $X$  is  $Y$ ’.

**ANSWER:**

(1): He is using modern semantics. His proof uses 0,1 as possible values for  $X, Y, \dots$

In a footnote he says that further details are needed if the value 0 is excluded.

His proof uses substitution—he actually proves no single equation  $E(X, Y) = 0$  can uniformly express ‘Some  $X$  is  $Y$ ’ using modern semantics.

(2) A minor modification of  $BF_3$  works for modern semantics.

T. Hailperin, *Boole's Logic and Probability*, 1986

He notes (p. 110) that Boole always converts  $X = VY$  into  $X(1 - Y) = 0$ , but he goes on to say that

**OBJECTION:** “in the case of the particular categorical there is no easy way out.”

**ANSWER:**

(1) He is using the modern semantics.

(2) There is an easy way out:  $V = VXY$

(3) Actually Boole does not always eliminate the  $V$  in the translations of universal statements — see Chap. *XV* of *Laws of Thought* where Boole gives a uniform treatment of syllogisms.