

Closure Homomorphisms

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The theory of closure homomorphisms will be patterned after the usual theory of homomorphisms for algebraic systems (cf. [3], [4]). In this paper we will see that the closure congruences submit to a simple axiomatization, and also if the closure space is algebraic, then we have a fundamental representation theorem.

A *closure space* (see [2]) is an ordered pair (\mathbf{C}, S) where \mathbf{C} is a closure operator on the set S (i.e., \mathbf{C} is an extensive, monotone, and idempotent set mapping). If θ is a homomorphism from an algebra (S, \mathcal{F}) into an algebra (S_1, \mathcal{F}_1) (for the definitions see [3], [4], etc.) and (\mathbf{C}, S) and (\mathbf{C}_1, S_1) are the *canonical* closure spaces associated with (S, \mathcal{F}) and (S_1, \mathcal{F}_1) respectively (i.e., $\mathbf{C}(P)$ is the smallest subalgebra containing P , etc.), then we have the identity

$$\theta[\mathbf{C}(P)] = \mathbf{C}_1[\theta(P)] \quad \text{for all } P \text{ in } S.$$

We abstract and define a *closure homomorphism* to be a mapping θ from a closure space (\mathbf{C}, S) to (\mathbf{C}_1, S_1) satisfying the identity $\mathbf{C}_1\theta = \theta\mathbf{C}$. [It is interesting to note that had we started out with relational structures (S, \mathcal{R}) (see [3]) instead of algebras (S, \mathcal{F}) , the corresponding condition would have been $\theta\mathbf{C} \subset \mathbf{C}_1\theta$, which for additive closures is precisely the study of continuous functions.]

I. AXIOMATIZATION OF CLOSURE CONGRUENCES

For any mapping θ from S into S_1 we know that $\theta^{-1}\theta$, considered as a subset of $S \times S$, is an equivalence relation. The equivalence relations on a closure space (\mathbf{C}, S) which we form in this manner from closure homomorphisms will be called *closure congruences*, and the family of closure congruences for (\mathbf{C}, S) will be denoted by $\mathfrak{R}(\mathbf{C}, S)$.

In the study of a general algebra (S, \mathcal{F}) we find that congruences are intrinsically characterized as equivalence relations which are subalgebras of

the product algebra $(S, \mathcal{F}) \times (S, \mathcal{F})$ (see [3]). The following theorem will give an intrinsic characterization of closure congruences. [Where we consider an equivalence relation E to induce a set mapping $E(P) = \{y : (x, y) \in E \text{ for some } x \text{ in } P\}$.]

LEMMA 1. *Let E be the equivalence relation associated with the closure homomorphism $\theta : (\mathbf{C}, S) \rightarrow (\mathbf{C}_1, S_1)$, (i.e., $E = \theta^{-1}\theta$). Then $ECE = EC$.*

Proof. $ECE = \theta^{-1}\theta\mathbf{C}\theta^{-1}\theta = \theta^{-1}\mathbf{C}_1\theta\theta^{-1}\theta = \theta^{-1}\mathbf{C}_1\theta = \theta^{-1}\theta\mathbf{C} = EC$.

By simple examples we can show that not every equivalence relation will satisfy such an identity, so the above lemma gives a genuine restriction on the equivalence relations which are congruences.

LEMMA 2. *If (\mathbf{C}, S) is a closure space and E is an equivalence relation on S , then the following are equivalent:*

- (i) $ECE = EC$
- (ii) $\mathbf{C}EC = EC$
- (iii) $\mathbf{C}E \subset EC$

Proof. From $ECE = EC$ follows $ECEC = ECC = EC$, and since $EC \subset \mathbf{C}EC \subset ECEC$, we see that (i) \Rightarrow (ii). Clearly (ii) \Rightarrow (iii), and from $\mathbf{C}E \subset EC$ follows $ECE \subset EEC = EC$, so (iii) \Rightarrow (i).

LEMMA 3. *Let (\mathbf{C}, S) be a closure space and E an equivalence relation on S such that $ECE = EC$. Let $S_1 = S/E$ and define \mathbf{C}_1 to be $\mu\mathbf{C}\mu^{-1}$, where μ is the canonical map from S onto S_1 . Then \mathbf{C}_1 is the unique closure on S_1 such that μ is a closure homomorphism with congruence E .*

Proof. \mathbf{C}_1 is clearly extensive and isotone. Also $\mathbf{C}_1^2 = \mu\mathbf{C}\mu^{-1}\mu\mathbf{C}\mu^{-1} = \mu\mathbf{C}E\mathbf{C}\mu^{-1} = \mu EC\mu^{-1} = \mu\mu^{-1}\mu\mathbf{C}\mu^{-1} = \mu\mathbf{C}\mu^{-1} = \mathbf{C}_1$, so \mathbf{C}_1 is idempotent. To show that μ is a homomorphism we note that $\mu\mathbf{C} = \mu\mu^{-1}\mu\mathbf{C} = \mu EC = \mu ECE = \mu\mathbf{C}E = (\mu\mathbf{C}\mu^{-1})\mu = \mathbf{C}_1\mu$. The uniqueness of \mathbf{C}_1 is straightforward.

THEOREM 1. *If (\mathbf{C}, S) is a closure space and E is an equivalence relation on S , then $E \in \mathfrak{R}(\mathbf{C}, S)$ if and only if $ECE = EC$.*

II. A REPRESENTATION THEOREM

LEMMA 4. *Let $E \in \mathfrak{R}(\mathbf{C}, S)$ and P, Q be subsets of S satisfying $E(P) \subset E(Q)$; then for any $x \in \mathbf{C}(P)$ there is a $y \in \mathbf{C}(Q)$ such that xEy .*

Proof. Note that $\mathbf{C}(P) \subset \mathbf{CE}(P) \subset \mathbf{CE}(Q) \subset \mathbf{EC}(Q)$.

LEMMA 5. Suppose E is a congruence for the closure space (\mathbf{C}, S) where \mathbf{C} is algebraic (i.e., $\mathbf{C}(P) = \bigcup \{\mathbf{C}(Q) : Q \subset P, Q \text{ finite}\}$). Then there is an abstract algebra (S, \mathcal{F}) whose canonical closure is \mathbf{C} and such that E is a congruence for (S, \mathcal{F}) in the sense of general algebra, and \mathcal{F} consists solely of symmetrical operations. [An operator $f(x_1, \dots, x_n)$ is symmetrical if its value is invariant under permutations of the arguments.]

Proof. First, well-order S in such a manner that if $x < y$ and $x \not E y$, then for $x_1 E x$ and $y_1 E y$ we have $x_1 < y_1$. For each finite set Q in S and each $x \in \mathbf{C}(Q)$ define the map $f_{Q,x}$ by:

$$\begin{aligned} f_{Q,x}(x_1, \dots, x_n) &= x \text{ if } \{x_1, \dots, x_n\} = Q \quad [\text{where } n = \text{Card}(Q)] \\ &= y \text{ if } \{x_1, \dots, x_n\} \neq Q, x_i E q_i, \text{ where} \\ &\quad \{q_1, \dots, q_n\} = Q, \text{ and } y \text{ is an element} \\ &\quad \text{in } \mathbf{C}(\{x_1, \dots, x_n\}) \text{ which is equivalent to } x \\ &= \inf\{x_1, \dots, x_n\} \text{ otherwise.} \end{aligned}$$

In this definition we consider $\{x_1, \dots, x_n\}$ to be a set (rather than an n -tuple). Also there is usually more than one candidate for y —select any one (the existence of such a y is guaranteed by Lemma 4). Clearly each $f_{Q,x}$ is a symmetric operator, and if each argument is replaced by an equivalent element, the values of the operator are equivalent (with respect to E). Then, considering the elements of $\mathbf{C}(\phi)$ as nullary operations, we obtain the desired algebra by letting

$$\mathcal{F} = \{f_{Q,x} : Q \subset S, x \in Q\} \cup \mathbf{C}(\phi).$$

THEOREM 2. Let θ be a homomorphism from the algebraic closure space (\mathbf{C}, S) onto the closure space (\mathbf{C}_1, S_1) . Then we can “fit” the spaces (\mathbf{C}, S) and (\mathbf{C}_1, S_1) with abstract algebras (S, \mathcal{F}) and (S_1, \mathcal{F}_1) such that (i) θ is a homomorphism from (S, \mathcal{F}) to (S_1, \mathcal{F}_1) , and (ii) the algebraic structures induce the respective closures.

Proof. Consider the congruence $\theta^{-1}\theta$ on (\mathbf{C}, S) and apply Lemma 5 to obtain an algebra (S, \mathcal{F}) which induces (\mathbf{C}, S) and for which $\theta^{-1}\theta$ is a congruence for (S, \mathcal{F}) . Then define an algebra (S_1, \mathcal{F}_1) by requiring that $g \in \mathcal{F}_1$ iff there is an $f \in \mathcal{F}$ such that

$$\theta f(x_1, \dots, x_n) = g[\theta(x_1), \dots, \theta(x_n)]$$

for all $x_1, \dots, x_n \in S$. With this the theorem follows easily.

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