# Four Unsolved Problems in Congruence Permutable Varieties

Ross Willard

University of Waterloo, Canada

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# Congruence permutable varieties

## Definition

A variety V is **congruence permutable** (or **CP**) if for each  $A \in V$ , Con A is a lattice of *permuting* equivalence relations.

 $\theta, \varphi$  permute if  $\theta \lor \varphi = \theta \circ \varphi = \varphi \circ \theta$ .

Examples of CP varieties: Any variety of . . .

- groups
- expansions of groups (e.g., rings, modules, non-associative rings, near rings, boolean algebras, etc.)
- quasi-groups in the language  $\{\cdot,/,\setminus\}$

#### But not:

• lattices, semilattices, semigroups, unary algebras.

## Basic facts about CP varieties

#### Fact 1

 $CP \Rightarrow$  congruence modularity.

# Fact 2 (Mal'tsev, 1954)

For a variety V, TFAE:

- ullet  $\mathcal V$  is CP.
- $\mathcal V$  has a term m(x,y,z) satisfying, in all  $\mathbf A \in \mathcal V$ ,

$$m(x,x,z) = z$$
 and  $m(x,z,z) = x$  (\*)

#### **Definitions**

- Mal'tsev term: a term m(x, y, z) satisfying (\*).
- Mal'tsev algebra: an algebra having a Mal'tsev term.
- Mal'tsev variety: a variety having a common Mal'tsev term.

# Fact 2 (restated)

CP varieties = Mal'tsev varieties.

## Aim of lecture

Mal'tsev algebras and varieties are . . .

- not "far" removed from groups, rings, near-rings, quasi-groups, etc...
- "old-fashioned," "solved."

**Aim of this lecture**: to correct this perception, by stating some open problems that:

- are general
- are of current interest
- are open
- are ripe for study in Mal'tsev algebras and varieties.

# 1. Subpower membership problem

Fix a finite algebra A.

# Subpower membership problem for A

Input:  $X \subseteq A^n$  and  $f \in A^n$   $(n \ge 1)$ 

Question: is  $f \in \operatorname{Sg}_{\mathbf{A}^n}(X)$ ?

How hard can it be?

#### HARD:

- Naive algorithm is EXPTIME
- There is no better algorithm (Friedman 1982; Bergman et al 1999.
   ADDED IN PROOF: Kozik, announced 2007).

However, for groups and rings the problem is solvable in polynomial time.

# Subpower membership problem for groups

(adapted from Sims 1971; Furst, Hopcroft, Luks 1980)

Fix a finite group **G**. Suppose  $H \leq \mathbf{G}^n$ .

Consider

$$H = H^{(0)} \ge H^{(1)} \ge \cdots \ge H^{(n)} = \{e\}$$

where

$$H^{(i)} = \{g \in H : g = (\underbrace{e, \dots, e}_{i}, *, \dots, *)\}.$$

Let  $M_i$  be a transversal for the cosets of  $H^{(i)}$  in  $H^{(i-1)}$ , including  $\hat{e}$ . Concretely:

**②** Every such form witnessed in H is represented in  $M_i$  exactly once.

Put  $M = \bigcup_{i=1}^n M_i$ .

#### Facts:

- **1** *M* is small (|M| = O(n))
- $\langle M \rangle = H$ . In fact,
  - $H = M_1 M_2 \cdots M_n$
  - every element  $h \in H$  is uniquely expressible in the form  $h = g_1g_2 \cdots g_n$  with each  $g_i \in M_i$ . ("Canonical form")
- **3** Given  $h \in H$ , we can find  $g_i \in M_i$  recursively, efficiently (knowing M).
- **3** Same algorithm tests arbitrary  $f \in G^n$  for membership in H.
- **1** Thus the subpower membership problem for **G** is solvable in polynomial time **if**, given  $X \subseteq G^n$ , we can find such an M for  $H = \langle X \rangle$ .

## Finding M.

## Rough idea. Given $X \subseteq G^n$ :

- Start with  $M_i = \{\widehat{e}\}$  for each i (so  $M = \{\widehat{e}\}$ ).
- For each  $g \in X$ , attempt to find the canonical form for g relative to M. (Will fail.)
- Each failure suggests an addition to some  $M_i$ .
  - The addition is always from  $\langle X \rangle$ .
  - **Action**: increment this  $M_i$  by the suggested addition.
- Repeat until each  $g \in X$  passes; i.e.,  $X \subseteq M_1 M_2 \cdots M_n$ .
- Next, for each  $g, h \in M$ , attempt to find the canonical form for gh.
  - Make additions to appropriate  $M_i$  upon each failure.
  - Loop until  $g, h \in M \Rightarrow gh$  passes.

## When to stop:

#### Lemma

 $M_1M_2\cdots M_n=\langle X\rangle$  as soon as  $g,h\in M\Rightarrow gh\in M_1M_2\cdots M_n$ .

## Corollary

The subpower membership problem is solvable in polynomial time for any finite group  $\mathbf{G}$ .

Remark. Similar technique works for any expansion of a group by multilinear operations (e.g., rings, modules, nonassociative rings).

## Corollary

The subpower membership problem is solvable in polynomial time for any finite ring or module.

## Partial generalization to Mal'tsev algebras

(Adapted from A. Bulatov and V. Dalmau, A simple algorithm for Mal'tsev constraints, *SIAM J. Comput.* **36** (2006), 16–27.)

Fix a finite algebra **A** with Mal'tsev term m(x, y, z).

## **Definition**

An *index* for  $A^n$  is a triple  $(i, a, b) \in \{1, 2, ..., n\} \times A \times A$ .

## **Definition**

A pair  $(g, h) \in A^n \times A^n$  witnesses (i, a, b) if

$$g = (x_1, \ldots, x_{i-1}, a, *, \ldots, *)$$

$$h = (x_1, \ldots, x_{i-1}, b, *, \ldots, *)$$

Consider  $\mathbf{B} \leq \mathbf{A}^n$ .

#### **Definition**

A **structured signature** for **B** is an *n*-tuple  $(M_1, \ldots, M_n)$  where

- (i = 1):
  - $M_1 \subseteq B$
  - Each form  $(a, *, ..., *) \in B$  is represented exactly once in  $M_1$ .
- **2**  $(2 \le i \le n)$ :
  - $M_i \subseteq B^2$
  - Each  $(g, h) \in M_i$  witnesses some index (i, a, b).
  - Each index (i, a, b) witnessed in B is represented exactly once in  $M_i$

Suppose  $(M_1, \ldots, M_n)$  is a structured signature for  $\mathbf{B} \leq \mathbf{A}^n$ . Let M be the set of all  $g \in A^n$  mentioned in  $(M_1, \ldots, M_n)$ .

#### Facts:

- $(M_1,\ldots,M_n)$  and M are small (|M|=O(n))
- **3** In fact, every element  $h \in B$  is expressible in the "canonical form"

$$h = m(m(\cdots m(m(f_1, g_2, h_2), g_3, h_3), \cdots), g_n, h_n)$$

with  $f_1 \in M_1$  and  $(g_i, h_i) \in M_i$  for  $2 \le i \le n$ .

• Note: can also require

$$g_2(2) = f_1(2)$$
  
 $g_3(3) = m(f_1, g_2, h_2)(3)$ , etc.

- **1**  $f_1, g_2, h_2, \ldots, g_n, h_n$  as above are unique for h and can be found recursively and efficiently.
- **3** Same algorithm tests arbitrary  $f \in A^n$  for membership in B.

This was enough for Bulatov and Dalmau to give a simple polynomial-time solution to the "CSP problem with Mal'tsev constraints."

Question: What about the subpower membership problem?

Suppose  $X \subseteq A^n$  and put  $\mathbf{B} = \operatorname{Sg}_{\mathbf{A}^n}(X)$ .

We can mimic the group algorithm by attempting to "grow" a structured signature for  ${\bf B}$ .

Sticking point: knowing when to stop.

### Problem 1

Using structured signatures or otherwise, is the Subpower Membership Problem for finite Mal'tsev algebras solvable in polynomial time?

# 2. The Pixley Problem

**Recall**: An algebra is *subdirectly irreducible* (or s.i.) if it cannot be embedding in a direct product of proper homomorphic images. (Equivalently, if its congruence lattice is monolithic.)

#### Definition

A variety V is a **Pixley variety** if:

- its language is finite
- every s.i. in V is finite (i.e., V is residually finite)
- ullet  $\mathcal V$  has arbitrarily large (finite) s.i.'s.

Question (Pixley, 1984): Is there a congruence distributive Pixley variety?

Answer (Kearnes, W., 1999): No.

Problem: Generalize.

## What is the situation for groups, rings, etc.?

- Commutative rings with 1.
  - No Pixley varieties here, since principal ideals are first-order definable.
- @ Groups.
  - Ol'shanskii (1969) described all residually finite varieties of groups.
  - None are Pixley varieties.
- 3 Rings (with or without 1).
  - McKenzie (1982) analyzed all residually small varieties of rings.
  - None are Pixley varieties.
- Modules.
  - Goodearl (priv. comm.): if R is an infinite, f.g. prime ring for which all nonzero ideals have finite index, then all nonzero injective left R-modules are infinite.
  - Kearnes (unpubl.): hence no variety of modules is Pixley.

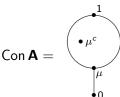
# Commutator Theory

Mal'tsev varieties (and congruence modular varieties) have a well-behaved theory of abelianness, solvability, centralizers and nilpotency.

#### Fundamental notions:

- " $\theta$  centralizes  $\varphi$ " ( $\theta, \varphi \in \mathsf{Con}\, \mathbf{A}$ ), i.e.,  $[\theta, \varphi] = 0$ .
- $\varphi^c$  = largest  $\theta$  which centralizes  $\varphi$ .

Frequently important: if A is s.i.:



**Fact**: if V is a CM Pixley variety, then (by the Freese-McKenzie theorem) for every s.i. in V,  $\mu^c$  is abelian.

# An argument

Suppose  ${\cal V}$  is a congruence modular variety in a finite language and having arbitrarily large finite s.i.'s.

**Case 1**: There exist arbitrarily large finite s.i.'s  $\mathbf{A} \in \mathcal{V}$  with  $|A/\mu^c|$  bounded.

• Use the module result to get an infinite s.i.  $\mathbf{A} \in \mathcal{V}$  with  $|A/\mu^c|$  bounded.

Case 2: Else.

- Define  $C(x, y, z, w) \leftrightarrow \text{``Cg}(x, y)$  centralizes Cg(z, w)."
- Assume C(x,y,z,w) is first-order definable in  $\mathcal{V}$ . Then use compactness to get an s.i.  $\mathbf{A} \in \mathcal{V}$  with  $|A/\mu^c|$  infinite.

Hence:

# Theorem (Kearnes, W., unpubl.)

If V is congruence modular and C(x, y, z, w) is definable in V, then V is **not** a Pixley variety.

#### Notes:

- Previous theorem handles all varieties of groups, rings and modules.
- Doesn't handle varieties of non-associative rings.

#### Problem 2

Does there exist a congruence permutable Pixley variety?

• What about varieties of non-associative rings?

# 3. McNulty's Problem

#### Definition

A variety  $\mathcal{V}$  is **strange** if

- its language is finite.
- V is locally finite.
- V is not finitely based.
- There exists a finitely based variety  $\mathcal{W}$  having exactly the same finite members as  $\mathcal{V}$ .

## Definition

A finite algebra is strange if the variety it generates is.

Question (Eilenberg, Schützenberger, 1976): Does there exist a strange finite algebra?

McNulty has asked the same question for varieties.

# Lemma (Cacioppo, 1993)

If **A** is strange, then it is inherently nonfinitely based (INFB).

# Theorem (McNulty, Székely, W., 2007?)

If  ${\bf A}$  can be shown to be INFB by the "shift automorphism method," then  ${\bf A}$  is **not** strange.

Examples of algebras known to be INFB but not by the shift automorphism method:

- (ADDED IN PROOF thank you, George): INFB Semigroups.
   Characterized by Sapir; George has checked that none are strange.
- 2 Isaev's non-associative ring (1989).

That's it!

## Problem 3

- 1 Is Isaev's algebra strange?
- ② Find more INFB algebras that are expansions of groups. Are any of them strange?

# 4. Dualizability

#### Definition

A finite algebra  $\mathbf{X}$   $\mathbf{\underline{M}}$  is dualizable if

- there exists an "alter ego" M ...
- ... partial operations ... relations ... discrete topology ...
- $\bullet$  ... **ISP** and **IS**<sub>c</sub>**P**<sup>+</sup> ...
- ... contravariant hom-functors ...
- ... dual adjunction  $(D, E, e, \varepsilon)$  ...
- AARRRGGHH!!! STOP THE INSANITY!!

# Dualizability

All that you need to know about dualizability (but were afraid to ask):

- "Dualizability" is a property that a finite algebra may, or may not, have.
- In practice, "dualizability" coincides with an apparently stronger property, called "finite dualizability."
- By a theorem of Zádori and myself, "finite dualizability" can be characterized in purely clone-theoretic terms.

# Classical clone theory

Fix a finite algebra A.

#### Recall that:

- $\bigcirc$  Inv(A) determines Clo(A), in the sense that

$$\forall f: A^n \to A, f \in Clo(\mathbf{A}) \text{ iff } f \text{ preserves every } r \in Inv(\mathbf{A}).$$

- Can speak of
  - a subset  $\mathcal{R} \subseteq Inv(\mathbf{A})$  determining  $Clo(\mathbf{A})$
  - Clo(A) being finitely determined.

### Old Theorem

The following are equivalent:

- R determines Clo(A)
- Every  $r \in Inv(\mathbf{A})$  can be defined from  $\Re$  by a  $\exists \& atomic$  formula.

# Partial operations with c.a.d. domains

Fix A.

A subset  $D \subseteq A^n$  is **c.a.d.** (conjunction-atomic-definable) if it is definable in **A** by a &atomic formula.

#### **Definition**

 $\mathrm{Clo}|_{\mathit{cad}}(\mathbf{A}) := \{ \mathrm{all} \ \mathrm{restrictions} \ \mathrm{of} \ \mathrm{term} \ \mathrm{operations} \ \mathrm{of} \ \mathbf{A} \ \mathrm{to} \ \mathrm{c.a.d.} \ \mathrm{domains} \}.$ 

#### Then:

- $Inv(\mathbf{A})$  determines  $Clo|_{cad}(\mathbf{A})$ , in the same sense:  $\forall f: D \to A \text{ with c.a.d. domain, } f \in Clo|_{cad}(\mathbf{A}) \text{ iff } f \text{ preserves every } r \in Inv(\mathbf{A}).$
- Can speak of
  - a subset  $\mathcal{R} \subseteq Inv(\mathbf{A})$  determining  $Clo|_{cad}(\mathbf{A})$
  - Clo<sub>cad</sub>(A) being finitely determined.

# Lemma/Definition

The following are equivalent:

- A is "finitely dualizable" ( ⇒ dualizable)
- **3** There is a finite set  $\mathcal{R} \subseteq Inv(\mathbf{A})$  such that every "hom-transparent"  $r \in Inv(\mathbf{A})$  is &atomic definable from  $\mathcal{R}$ .

**Def**.  $r \in Inv(A)$  is hom-transparent (or balanced) if

- Every homomorphism  $h: \mathbf{r} \to \mathbf{A}$  is a coordinate projection, and
- No two coordinate projections are the same.

# **Dualizability problem**: which finite **A** are (finitely) dualizable?

- O CD case:
  - (finitely) dualizable ⇔ A has a near-unanimity term
    - ← by Baker-Pixley, 
       ⇒ by (Davey, Heindorf, McKenzie, 1995)
- 2 Commutative rings with 1:
  - - (Clark, Idziak, Sabourin, Szabó, W., 2001)
- Groups:
  - (finitely) dualizable ⇔ G generates a residually small variety.
    - ⇒ by (Quackenbush, Szabó, 2002), ← by (Nickodemus, 2007?)
- Rings (with or without 1):
  - (finitely) dualizable  $\stackrel{?}{\Leftrightarrow}$  **R** generates a residually small variety.
    - ⇒ by (Szabó, 1999), ← by recent work of Kearnes, Szendrei?
- But:
  - if  $G = S_3$ , then  $G_G$  is *not* dualizable, yet generates a residually small variety (Idziak, unpubl., 1994)
  - ∃ expansion of (Z<sub>4</sub>, +) that is (finitely) dualizable, yet generates a residually large variety (Davey, Pitkethly, W., 2007?)

## Problem 4

- Which finite Mal'tsev algebras are (finitely) dualizable?
  - Can we at least answer this for expansions of groups?
- 2 Is the answer to (1) decidable?