Series-parallel posets having a near-unanimity polymorphism

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AMS Fall Western Sectional Meeting Denver, October 8, 2016 All posets are finite.

If P, Q are posets, then P + Q is their **ordinal sum**:

 $\mathbf{P} \cup \mathbf{Q}$ is their disjoint union.

$$1 \cup 1 = \bullet \bullet = 2$$

$$1+1 =$$

Definition. Let **P** be a poset.

A function $f: P^n \to P$ is a **near unanimity (NU) polymorphism** of **P** if

- n > 3.
- $\forall 1 \leq i \leq n, \forall a, b \in P$,

$$f(a, a, \ldots, a, b, a, \ldots, a) = a$$
 \uparrow
 i

• f is monotone in each variable.

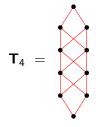
Clone theorists (last century) and CSPers (this century) care about which posets have an NU polymorphism.

Key examples

$$\mathsf{T}_2 = \bigvee_{=}^{\bullet}$$

$$= 1 + 2 + 2 + 1$$

$$T_3 =$$



Facts

Every lattice-ordered poset has an NU polymorphism of arity 3.

T₂: Has an NU polymorphism of arity 5 (Demetrovics et al, 1984).

T₃: Does not have an NU polymorphism. (Demetrovics et al, 1984) Does have "weaker" (Taylor) polymorphisms (McKenzie, 1990).

T₄: Does not even have "weaker" polymorphisms (Dem. & Rónyai, 1989).

 T_2, T_3, T_4, \dots are examples of **series-parallel posets**.

Definition

A poset is **series-parallel** if it can be constructed from (copies of) 1 by finitely many applications of + and \cup .

Equivalently (Valdes, Tarjan, Lawler 1982), a poset is series-parallel iff does not embed into it.

Dalmau, Krokhin, Larose (2008) characterized those series-parallel posets which have "weaker" (Taylor) polymorphisms:

- \bullet By "forbidden retracts" (list of 5, including $T_4,\ 2+2,$ and 2+2+2).
- By an internal characterization, easily checkable in polynomial time.

Our main result: We can do something similar for NU polymorphisms.

New operations: $P \boxtimes Q$, $P \triangle Q$, $P \nabla Q$, and $P \lozenge Q$

 $\mathbf{P} \mathbf{X} \mathbf{Q}$: defined when \mathbf{P} has 1 and \mathbf{Q} has 0.

 $\mathbf{P} \triangle \mathbf{Q}$: defined when both \mathbf{P} and \mathbf{Q} have 1.

 $\mathbf{P} \nabla \mathbf{Q}$: defined when both \mathbf{P} and \mathbf{Q} have 0.

 $\mathbf{P} \lozenge \mathbf{Q}$: defined when both \mathbf{P} and \mathbf{Q} have $1 \neq 0$.

$$\bigvee$$
 \times \bigvee $=$ \bigvee

$$\triangle \triangle \triangle = \triangle$$

$$\bigvee \bigvee \bigvee = \bigvee$$

Here is our result.

Theorem

Let **P** be a series-parallel poset. TFAE:

- P has an NU polymorphism.
- ② P does not retract onto 2+2, 2+2+1 or its dual, or T_3 .
- **3** Each connected component of **P** having more than one element is in the closure of $\{1+1\}$ under +, X, A, A, A, A.

About the proof:

- Hardest part is showing that **P** being in the closure of $\{1+1\}$ under $+, \Sigma, \triangle, \nabla, \diamondsuit$ implies **P** has an NU polymorphism.
- Use a complicated induction on the construction of P.
- On not actually construct an NU; instead, use criterion for existence due to Kun and Szabó (2001).

Thus we have no control over (and know nothing about) the NU's arity.

Problem

For fixed $k \ge 3$, characterize the series-parallel posets which have a k-ary NU polymorphism.

Added post-lecture: the above problem is solved by Corollary 3.3 of L. Zádori, Series parallel posets with nonfinitely generated clones, *Order* **10** (1993), 305–316.