# Reconstruction of polytopes and Kalai＇s conjecture on reconstruction of spheres 

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## Convex polytopes

(Convex) polytope $P$ :
convex hull of finitely many points in Euclidian space.
The graph $G(P)$ : the graph consisting of the vertices and edges of $P$.


P

$G(P)$

Simple polytope $P$ :
number of edges incident to each vertex equals the dimension of $P$.

## Reconstruction of polytopes

## Theorem (Blind-Mani, 1987)

If $P$ is a simple polytope, then the graph $G(P)$ determines the entire combinatorial structure of $P$.


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Kalai, 1988: A simple constructive proof.

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This holds for arbitrary polytopes (not only simple) in dimension 3, but not in higher dimensions.

## Example

Let $\Delta_{m}$ be a $m$-dimensional simplex. The following are two non isomorphic 6-dimensional polytopes with the same graph (complete graph on 7 vertices)

$$
\left(\Delta_{2} \times \Delta_{4}\right)^{*} \not \not\left(\left(\Delta_{3} \times \Delta_{3}\right)^{*}\right.
$$

## Duality of polytopes

Every nonempty $d$-polytope $P$ in $\mathbb{R}^{d}$ admits a dual polytope in $\mathbb{R}^{d}$ :

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P^{*}=\left\{y \in \mathbb{R}^{d}: x^{T} y \leq 1 \text { for all } x \in P\right\}
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where $P$ is assumed to contain the origin in its interior.


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Under this duality:


## Simple vs simplicial

## Simplicial polytope $P$ :

all faces are simplices.
The facet-ridge graph $G_{F R}(P)$ :
the graph whose vertices are facets of $P$
two facets are connected by an edge if they intersect in a ridge.

$$
\begin{aligned}
P \text { is simple } & \longleftrightarrow P^{*} \text { is simplicial } \\
G(P) & =G_{F R}\left(P^{*}\right)
\end{aligned}
$$



## Reconstruction of polytopes and spheres

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Theorem (Blind-Mani, 1987)
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Conjecture (Blind-Mani, 1987; Kalai, 2009)
Simplicial spheres are completely determined by their facet-ridge graphs.
A simplicial sphere is a simplicial complex which is homeomorphic to a sphere.


## Most spheres are not polytopal

For $d \geq 3$, most $d$-spheres are not polytopal.

- Goodman-Pollack, 1986
- Kalai, 1988
- Pfeifle-Ziegler, 2004

Deciding polytopality of spheres is a difficult problem
Mnëv's Universality theorem: Realization spaces of polytopes can take arbitrary (semi-algebraic) shapes and thus can exhibit all kinds of pathologies.

The realizability problem for 4-polytopes is NP-hard.


## Goal

## Our initial goal was:

Look for a counterexample to Kalai's Conjecture among a special family of simplicial spheres which are conjectured to be polytopal.
(kill two conjectures at once)
Instead:
We proved the conjecture for this family. (spherical subword complexes)

Rest of the talk:
Introduce subword complexes and state our main result.

## Subword complexes preliminaries

Symmetric group $\mathbb{S}_{n+1}$ :
group of permutations of $\{1, \ldots, n+1\}$

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length of w : smallest $r$ such that $w=s_{i_{1}} \ldots s_{i_{r}}$ longest element: permutation $[n+1, \ldots, 1]$
reduced expression for $w$ : expression for $w$ of minimal length

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In this talk: finite Coxeter groups
(very similar to the symmetric group)

## Subword complexes

$W$ finite Coxeter group with generating set $S$
$Q=\left(q_{1}, \ldots, q_{m}\right)$ a word in $S$
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$\pi \in W$

Definition (Knutson-Miller, 2004)
The subword complex $\Delta(Q, \pi)$ is the simplicial complex whose
faces $\longleftrightarrow$ subwords $P$ of $Q$ such that $Q \backslash P$ contains a reduced expression of $\pi$

Knutson-Miller. Gröbner geometry of Schubert polynomials. Ann. Math., 161(3), '05 Knutson-Miller. Subword complexes in Coxeter groups. Adv. Math., 184(1), '04

## Subword complexes - Example 1

In type $A_{2}$ :
$W=\mathbb{S}_{3}, S=\left\{s_{1}, s_{2}\right\}=\left\{\left(\begin{array}{ll}1 & 2),(23)\}\end{array}\right.\right.$

## Subword complexes - Example 1

In type $A_{2}$ :
$W=\mathbb{S}_{3}, S=\left\{s_{1}, s_{2}\right\}=\left\{\left(\begin{array}{ll}1 & 2\end{array}\right),\left(\begin{array}{ll}2 & 3\end{array}\right)\right\}$
$Q=\begin{gathered}\left(s_{1}, s_{2}, s_{1}, s_{2}, s_{1}\right. \\ q_{1}, q_{2}, q_{3}, q_{4}, q_{5}\end{gathered}$ and $\pi=\left[\begin{array}{ll}3 & 2\end{array} 1\right]$

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$$
\begin{gathered}
q_{2} \\
0
\end{gathered}
$$

$q_{3} \circ$
$\Delta(Q, \pi)$ is isomorphic to

- $q_{1}$
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## Subword complexes - Example 2

In type $A_{3}$ :
$W=\mathbb{S}_{4}, S=\left\{s_{1}, s_{2}, s_{3}\right\}=\{(12),(23),(34)\}$
$Q=\begin{gathered}\left(s_{1}, s_{2}, s_{1}, s_{2}, s_{1}, s_{3}\right) \\ q_{1}, q_{2}, q_{3}, q_{4}, q_{5}, q_{6}\end{gathered}$ and $\pi=\left[\begin{array}{ll}3 & 2\end{array} 1\right]=s_{1} s_{2} s_{1}=s_{2} s_{1} s_{2}$

## Subword complexes - Example 2

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## Subword complexes

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Subword complexes are vertex decomposable spheres or balls.

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Spherical subword complexes are polytopal.

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Spherical subword complexes are polytopal.

Special cases include:

- Cyclic polytopes
- Duals of associahedra
- Cluster complexes of cluster algebras of finite type
- Duals of pointed-pseudotriangulation polytopes
- Simplicial multi-associahedra (conjectured)

Woo, Pilaud-Pocchiola, Serrano-Stump, Stump, C.-Labbé-Stump, Rote-Santos-Streinu, Jonsson, ...

## Our main theorem

## Theorem (C.-Doolittle)

Spherical subword complexes of finite type are completely determined by their facet-ridge graph. In other words, they satisfy Kalai's Conjecture.

Our current proof is not constructive.
It is based on the topological tools developed by Blind and Mani.

Thank you!

