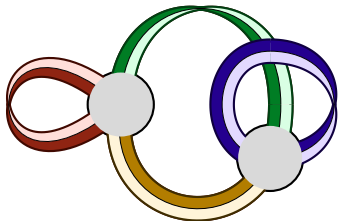
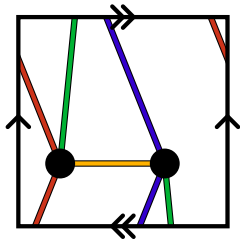


# Jack Symmetric Functions and the Non-Orientability of Rooted Maps

Michael La Croix

University of Waterloo

January 4, 2012



# Graphs, Surfaces, and Maps

## Definition

A **surface** is a compact 2-manifold without boundary.

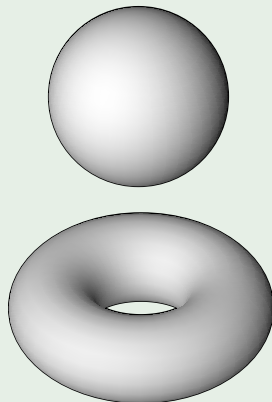
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A **graph** is a finite set of *vertices* together with a finite set of *edges*, such that each edge is associated with either one or two vertices.

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A **map** is a 2-cell embedding of a graph in a surface.

## Example



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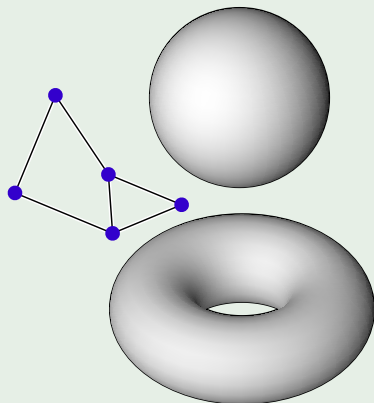
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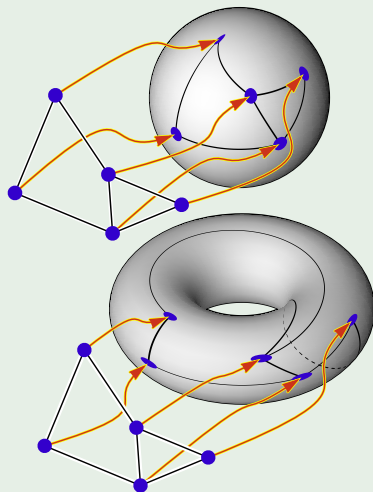
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# Ribbon Graphs and Flags

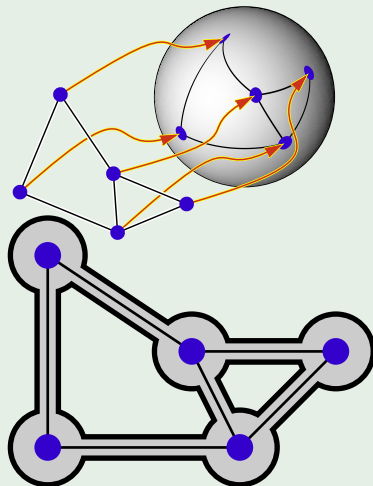
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The neighbourhood of the graph determines a **ribbon graph**, and the boundaries of ribbons determine **flags**.

## Definition

Automorphisms permute flags, and a **rooted map** is a map together with a distinguished orbit of flags.

## Example



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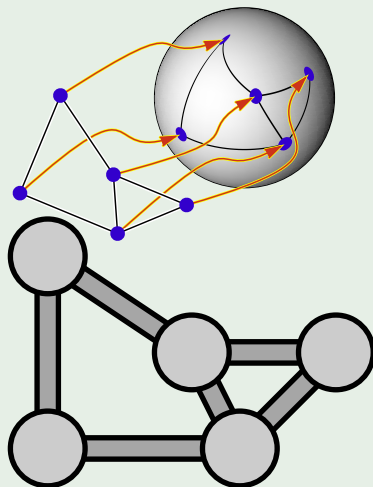
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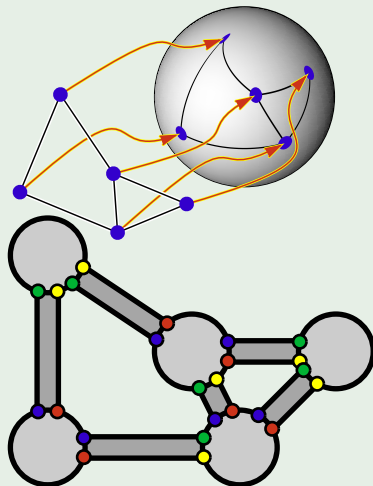
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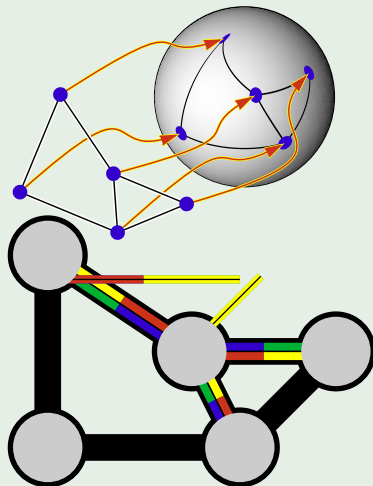
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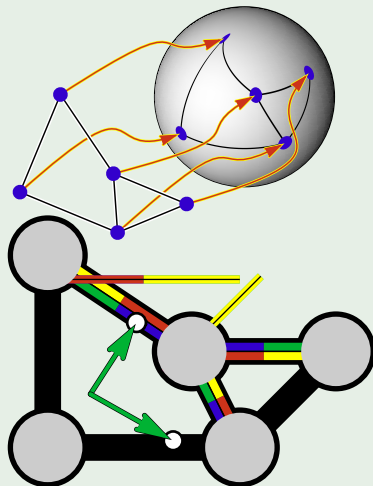
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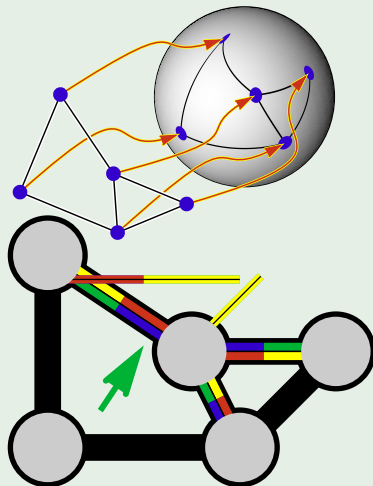
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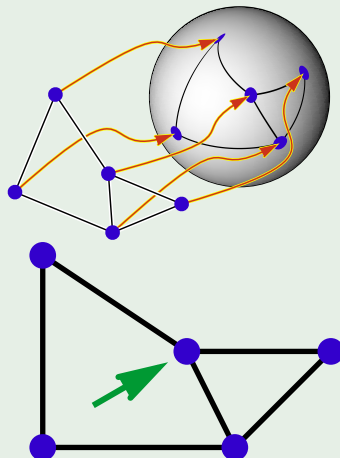
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## Note

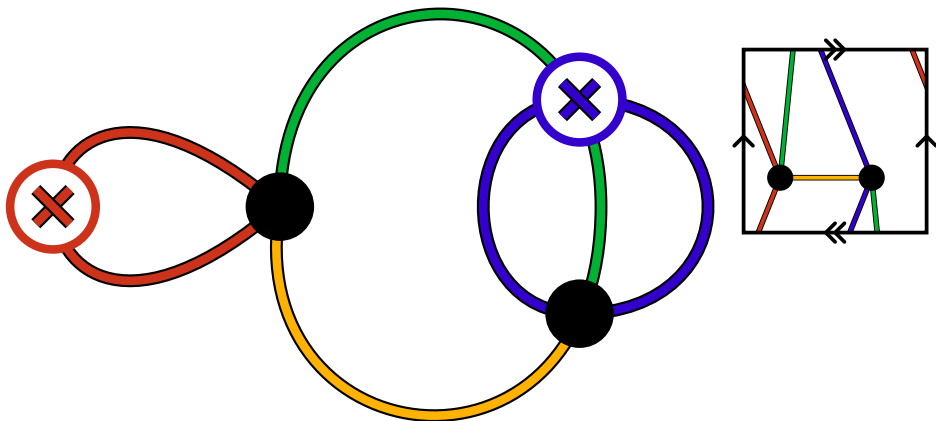
The map with no edges, , has a rooting.

## Example



# A Combinatorial Encoding of Maps

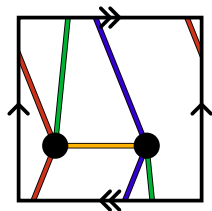
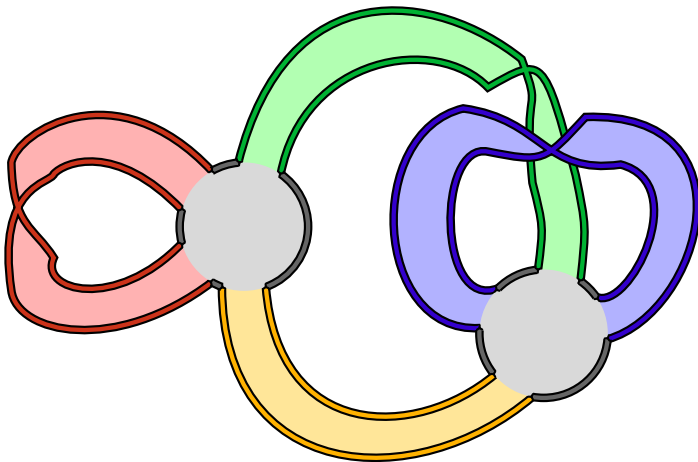
Equivalence classes can be encoded by perfect matchings of flags.



Start with a ribbon graph.

# A Combinatorial Encoding of Maps

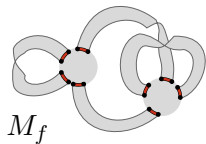
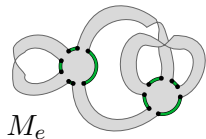
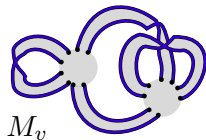
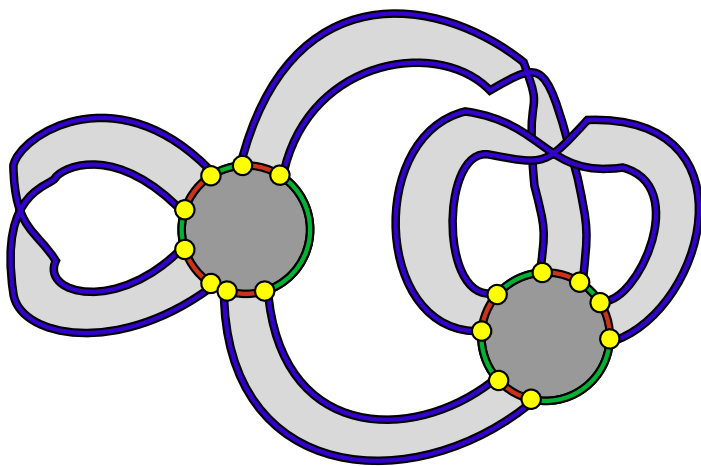
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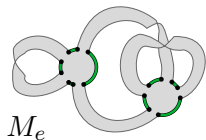
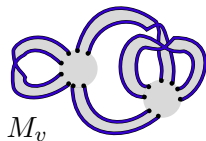
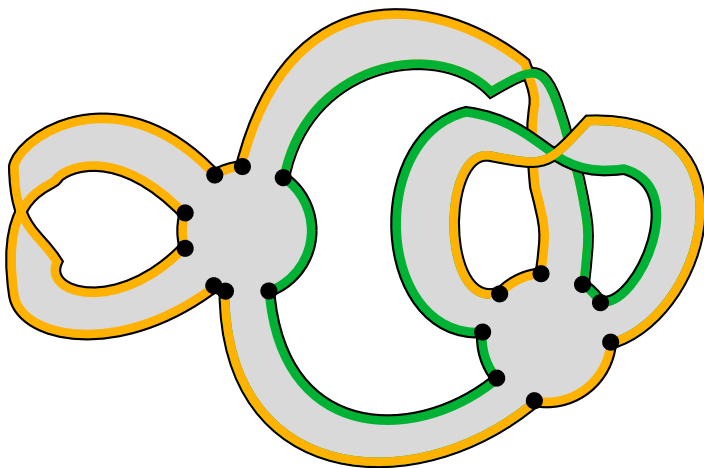
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Ribbon boundaries determine 3 perfect matchings of flags.

# A Combinatorial Encoding of Maps

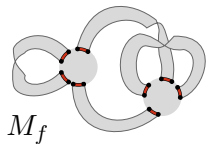
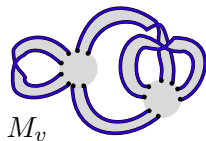
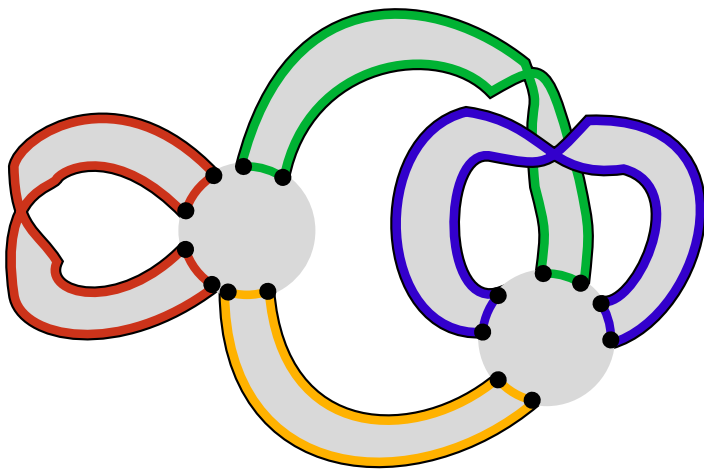
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# A Combinatorial Encoding of Maps

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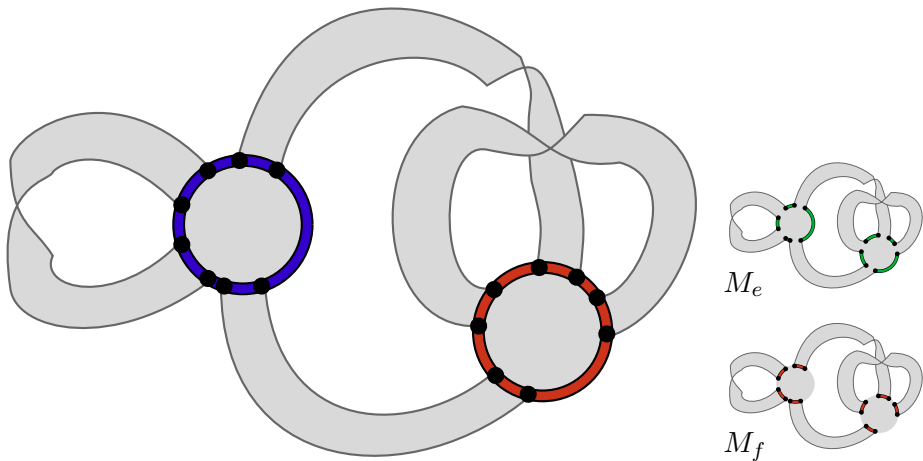


Pairs of matchings determine, faces, **edges**,



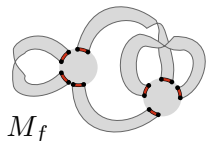
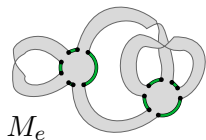
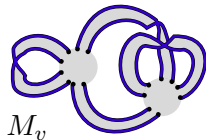
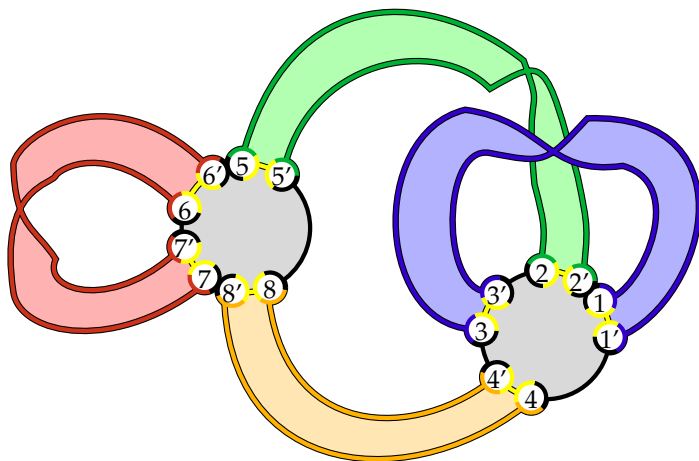
# A Combinatorial Encoding of Maps

Equivalence classes can be encoded by perfect matchings of flags.



Pairs of matchings determine, faces, edges, and **vertices**.

# A Combinatorial Encoding of Maps



$$M_v = \{\{1, 3\}, \{1', 3'\}, \{2, 5\}, \{2', 5'\}, \{4, 8'\}, \{4', 8\}, \{6, 7\}, \{6', 7'\}\}$$

$$M_e = \{\{1, 2'\}, \{1', 4\}, \{2, 3'\}, \{3, 4'\}, \{5, 6'\}, \{5', 8\}, \{6, 7'\}, \{7, 8'\}\}$$

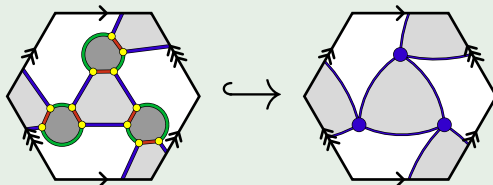
$$M_f = \{\{1, 1'\}, \{2, 2'\}, \{3, 3'\}, \{4, 4'\}, \{5, 5'\}, \{6, 6'\}, \{7, 7'\}, \{8, 8'\}\}$$

# Hypermaps

Generalizing the combinatorial encoding, an arbitrary triple of perfect matchings determines a **hypermap** when the triple induces a connected graph, with cycles of  $M_e \cup M_f$ ,  $M_e \cup M_v$ , and  $M_v \cup M_f$  determining vertices, hyperfaces, and hyperedges. [▶ Example](#)

Hypermaps both **specialize** and generalize maps.

## Example



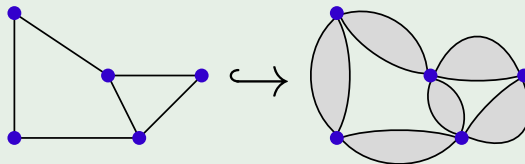
Hypermaps can be represented as face-bipartite maps.

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Hypermaps both specialize and **generalize** maps.

## Example



Maps can be represented as hypermaps with  $\epsilon = [2^n]$ .

# The Hypermap Series

## Definition

The **hypermap series** for a set  $\mathcal{H}$  of hypermaps is the combinatorial sum

$$H(\mathbf{x}, \mathbf{y}, \mathbf{z}) := \sum_{\mathfrak{h} \in \mathcal{H}} \mathbf{x}^{\nu(\mathfrak{h})} \mathbf{y}^{\phi(\mathfrak{h})} \mathbf{z}^{\epsilon(\mathfrak{h})}$$

where  $\nu(\mathfrak{h})$ ,  $\phi(\mathfrak{h})$ , and  $\epsilon(\mathfrak{h})$  are the vertex-, hyperface-, and hyperedge-degree partitions of  $\mathfrak{h}$ . [▶ Example](#)

## Example

Rootings of  contribute  $12 \left( x_2^3 x_3^2 \right) \left( y_3 y_4 y_5 \right) z_2^6$  to the sum.

# The Hypermap Series

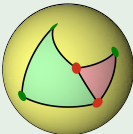
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# Explicit Formulae

The hypermap series can be computed explicitly when  $\mathcal{H}$  consists of orientable hypermaps or all hypermaps. [▶ sketch](#)

## Theorem (Jackson and Visentin - 1990)

When  $\mathcal{H}$  is the set of orientable hypermaps, [▶ encoding details](#)

$$H_{\mathcal{O}}(p(\mathbf{x}), p(\mathbf{y}), p(\mathbf{z}); 0) = t \frac{\partial}{\partial t} \ln \left( \sum_{\theta \in \mathcal{P}} H_{\theta} s_{\theta}(\mathbf{x}) s_{\theta}(\mathbf{y}) s_{\theta}(\mathbf{z}) \right) \Big|_{t=0}.$$

## Theorem (Goulden and Jackson - 1996)

When  $\mathcal{H}$  is the set of all hypermaps (orientable and non-orientable),

$$H_{\mathcal{A}}(p(\mathbf{x}), p(\mathbf{y}), p(\mathbf{z}); 1) = 2t \frac{\partial}{\partial t} \ln \left( \sum_{\theta \in \mathcal{P}} \frac{1}{H_{2\theta}} Z_{\theta}(\mathbf{x}) Z_{\theta}(\mathbf{y}) Z_{\theta}(\mathbf{z}) \right) \Big|_{t=0}.$$

# Jack Symmetric Functions

Jack symmetric functions, ▸ Definition, are a one-parameter family, denoted by  $\{J_\theta(\alpha)\}_\theta$ , that generalizes both Schur functions and zonal polynomials.

## Proposition (Stanley - 1989)

*Jack symmetric functions are related to Schur functions and zonal polynomials by:*

$$J_\lambda(1) = H_\lambda s_\lambda,$$

$$J_\lambda(2) = Z_\lambda,$$

*and*

$$\langle J_\lambda, J_\lambda \rangle_1 = H_\lambda^2,$$

$$\langle J_\lambda, J_\lambda \rangle_2 = H_{2\lambda},$$

*where  $2\lambda$  is the partition obtained from  $\lambda$  by multiplying each part by two.*



# A Generalized Series

## $b$ -Conjecture (Goulden and Jackson - 1996)

*The generalized series,*

$$\begin{aligned} H(p(\mathbf{x}), p(\mathbf{y}), p(\mathbf{z}); b) \\ &:= (1+b)t \frac{\partial}{\partial t} \ln \left( \sum_{\theta \in \mathcal{P}} \frac{J_\theta(\mathbf{x}; 1+b) J_\theta(\mathbf{y}; 1+b) J_\theta(\mathbf{z}; 1+b)}{\langle J_\theta, J_\theta \rangle_{1+b}} \right) \Big|_{t=0} \\ &= \sum_{n \geq 0} \sum_{\nu, \phi, \epsilon \vdash n} c_{\nu, \phi, \epsilon}(b) p_\nu(\mathbf{x}) p_\phi(\mathbf{y}) p_\epsilon(\mathbf{z}), \end{aligned}$$

*has an combinatorial interpretation involving hypermaps. In particular*

$$c_{\nu, \phi, \epsilon}(b) = \sum_{\mathfrak{h} \in \mathcal{H}_{\nu, \phi, \epsilon}} b^{\beta(\mathfrak{h})} \text{ for some invariant } \beta \text{ of rooted hypermaps.}$$

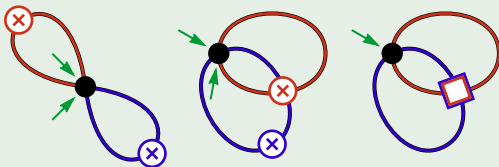
# A $b$ -Invariant

The  $b$ -Conjecture assumes that  $c_{\nu,\phi,\epsilon}(b)$  is a polynomial, and numerical evidence suggests that its degree is the genus of the hypermaps it enumerates. A  $b$ -invariant must:

- 1 be **zero** for orientable hypermaps,
- 2 be positive for non-orientable hypermaps, and
- 3 depend on rooting.

## Example

Rootings of precisely three maps are enumerated by  $c_{[4],[4],[2^2]}(b) = 1 + b + 3b^2$ .



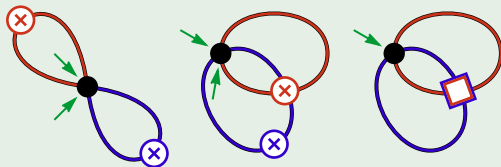
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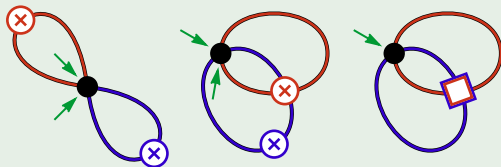
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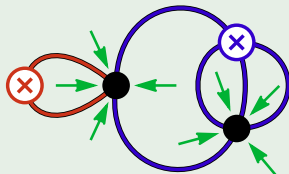
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## Example

There are precisely eight rooted maps enumerated by  $c_{[4,4],[3,5],[2^4]}(b) = 8b^2$ .



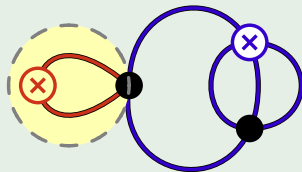
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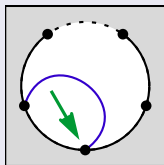
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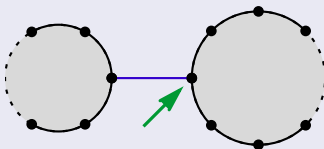
# A root-edge classification

There are four possible types of root edges in a map.

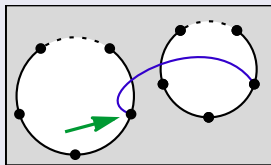
Borders



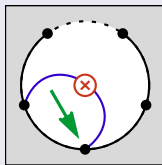
Bridges



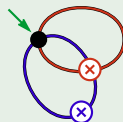
Handles



Cross-Borders

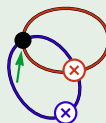


Example



A handle

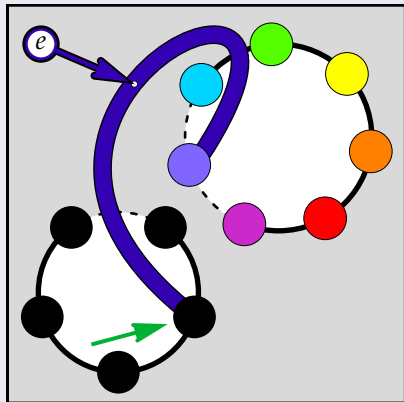
Example



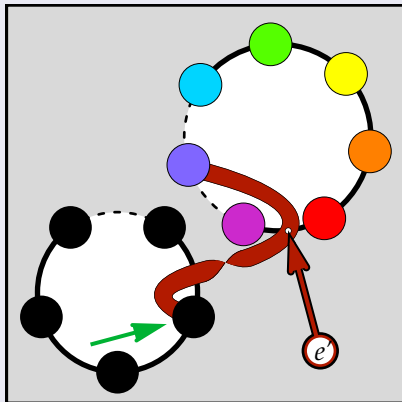
A cross-border

# A root-edge classification

Handles occur in pairs



Untwisted



Twisted

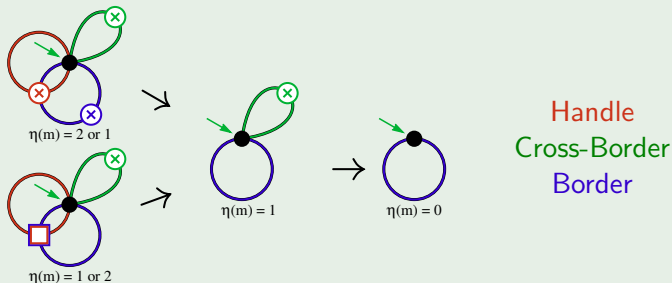


# A family of invariants

## The invariant $\eta$

- Iteratively deleting the root edge assigns a type to each edge in a map.
- An invariant,  $\eta$ , is given by
$$\eta(\mathfrak{m}) := (\# \text{ of cross-borders}) + (\# \text{ of twisted handles}).$$
- Different handle twisting determines a different invariant.

### Example



# Main result (marginal $b$ -invariants exist)

## Theorem (La Croix)

*If  $\phi$  partitions  $2n$  and  $\eta$  is a member of the family of invariants then,*

$$d_{v,\phi}(b) := \sum_{\ell(\nu)=v} c_{\nu,\phi,[2^n]}(b) = \sum_{\mathbf{m} \in \mathcal{M}_{v,\phi}} b^{\eta(\mathbf{m})}.$$

## Proof (sketch).

- A [▶ generating series](#) for maps with respect  $\eta$  satisfies a [▶ PDE](#) with a unique solution.
- The corresponding specialization of  $H$  has an analytic presentation.  
[▶ Details](#)
- An algebraic refinement to distinguish between root and non-root faces in the generating series satisfies the same PDE. □

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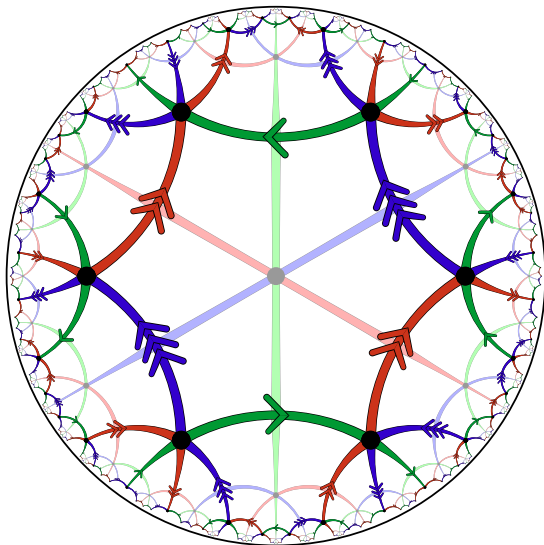
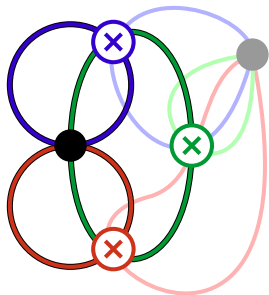
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## Implications of the proof

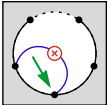
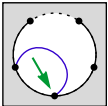
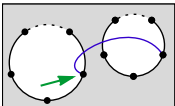
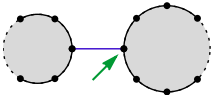
- $d_{v,\phi}(b)$  is of the form  $\sum_{0 \leq i \leq g/2} h_{v,\phi,i} b^{g-2i} (1+b)^i$ .
- The degree of  $d_{v,\phi}(b)$  is the genus of the maps it enumerates.
- The top coefficient,  $h_{v,\phi,0}$ , enumerates **unhandled** maps.
- $\eta$  and root-face degree are independent among maps with given  $\phi$ .

The End

Thank You

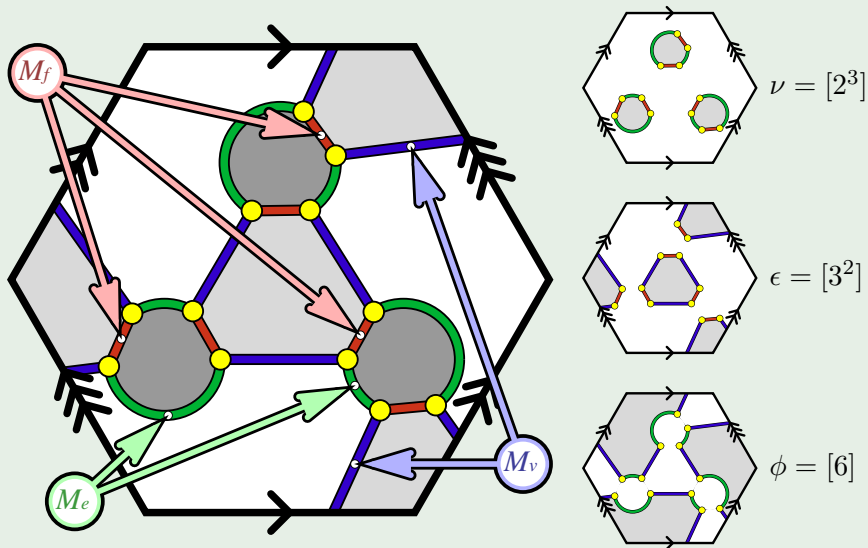


# Finding a partial differential equation

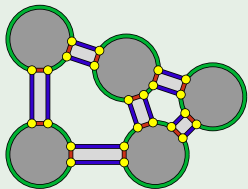
Root-edge type	Schematic	Contribution to $M$
Cross-border		$z \sum_{i \geq 0} (i+1) b r_{i+2} \frac{\partial}{\partial r_i} M$
Border		$z \sum_{i \geq 0} \sum_{j=1}^{i+1} r_j y_{i-j+2} \frac{\partial}{\partial r_i} M$
Handle		$z \sum_{i,j \geq 0} (1+b) j r_{i+j+2} \frac{\partial^2}{\partial r_i \partial y_j} M$
Bridge		$z \sum_{i,j \geq 0} r_{i+j+2} \left( \frac{\partial}{\partial r_i} M \right) \left( \frac{\partial}{\partial r_j} M \right)$

Return

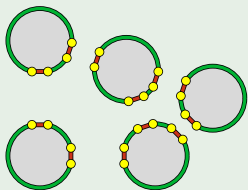
# Example



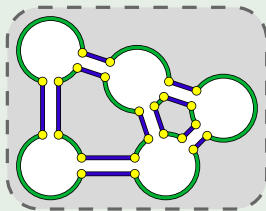
## Example



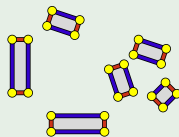
is enumerated by  $(x_2^3 x_3^2) (y_3 y_4 y_5) (z_2^6)$ .



$$\nu = [2^3, 3^2]$$



$$\phi = [3, 4, 5]$$



$$\epsilon = [2^6]$$

◀ Return

# Jack Symmetric Functions

With respect to the inner product defined by

$$\langle p_\lambda(\mathbf{x}), p_\mu(\mathbf{x}) \rangle_\alpha = \delta_{\lambda,\mu} \frac{|\lambda|!}{|C_\lambda|} \alpha^{\ell(\lambda)},$$

Jack symmetric functions are the unique family satisfying:

(P1) (Orthogonality) If  $\lambda \neq \mu$ , then  $\langle J_\lambda, J_\mu \rangle_\alpha = 0$ .

(P2) (Triangularity)  $J_\lambda = \sum_{\mu \preccurlyeq \lambda} v_{\lambda\mu}(\alpha) m_\mu$ , where  $v_{\lambda\mu}(\alpha)$  is a rational function in  $\alpha$ , and ' $\preccurlyeq$ ' denotes the natural order on partitions.

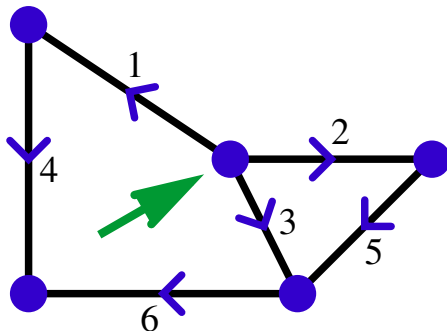
(P3) (Normalization) If  $|\lambda| = n$ , then  $v_{\lambda, [1^n]}(\alpha) = n!$ .

◀ Return



# Encoding Orientable Maps

- 1 Orient and label the edges.
- 2 This induces labels on flags.
- 3 Clockwise circulations at each vertex determine  $\nu$ .
- 4 Face circulations are the cycles of  $\epsilon\nu$ .



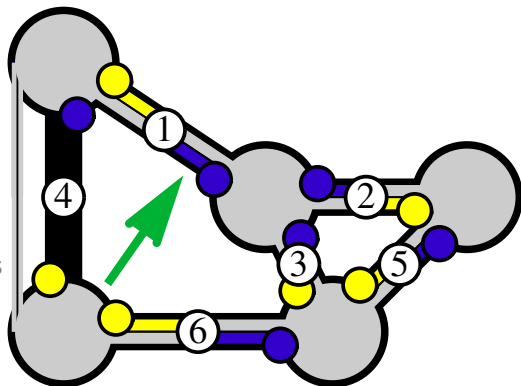
$$\epsilon = (1 \ 1')(2 \ 2')(3 \ 3')(4 \ 4')(5 \ 5')(6 \ 6')$$

$$\nu = (1 \ 2 \ 3)(1' \ 4')(2' \ 5')(3' \ 5' \ 6')(4' \ 6')$$

$$\epsilon\nu = \phi = (1 \ 4 \ 6' \ 3')(1' \ 2 \ 5 \ 6 \ 4')(2' \ 3 \ 5')$$

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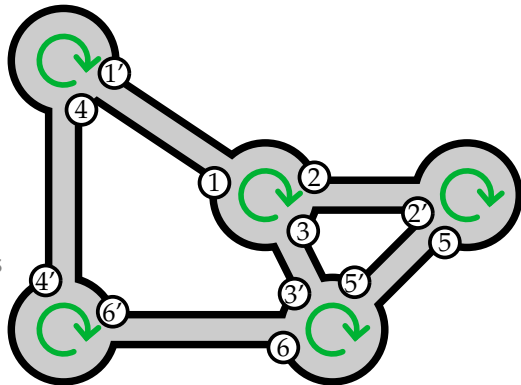
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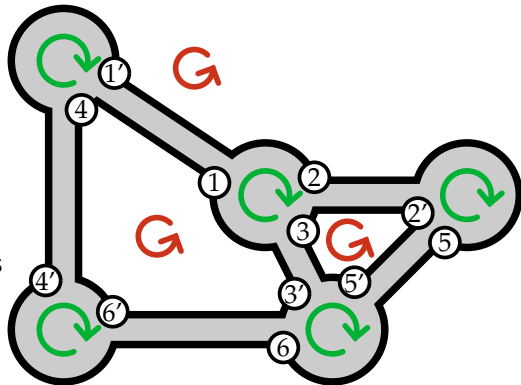
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# The Map Series

An enumerative problem associated with maps is to determine the number of rooted maps with specified vertex- and face- degree partitions.

## Definition

The **map series** for a set  $\mathcal{M}$  of rooted maps is the combinatorial sum

$$M = M(x, \mathbf{y}, z, \mathbf{r}; b) := \sum_{\mathfrak{m} \in \mathcal{M}} x^{|V(\mathfrak{m})|} \mathbf{y}^{\phi(\mathfrak{m}) \setminus \rho(\mathfrak{m})} z^{|E(\mathfrak{m})|} r_{\rho(\mathfrak{m})} b^{\eta(\mathfrak{m})},$$

where the sum is taken over all rooted maps, including the map with no edges,  $V(\mathfrak{m})$  is the vertex set of  $\mathfrak{m}$ ,  $\phi(\mathfrak{m})$  is the face-degree partition of  $\mathfrak{m}$ ,  $\rho(\mathfrak{m})$  is the degree of the root face of  $\mathfrak{m}$ , and  $E(\mathfrak{m})$  is the edge set of  $\mathfrak{m}$ .

◀ Return

# How to enumerate maps with symmetric functions

- Instead of counting rooted maps, we can count labelled hypermaps. This adds easily computable multiplicities.
- Labelled counting problems are turned into problems involving counting factorizations.
- These can be answered via character theory.
- Appropriate characters appear as coefficients of symmetric functions.
- The logarithms restrict to connected maps, and the differential operators remove the decoration.

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# A Specialization

## Definition

For a function  $f: \mathbb{R}^N \rightarrow \mathbb{R}$ , define an expectation operator  $\langle \cdot \rangle$  by

$$\langle f \rangle_{1+b} := c_{1+b} \int_{\mathbb{R}^N} |V(\boldsymbol{\lambda})|^{\frac{2}{1+b}} f(\boldsymbol{\lambda}) e^{-\frac{1}{2(1+b)} p_2(\boldsymbol{\lambda})} d\boldsymbol{\lambda},$$

with  $c_{1+b}$  chosen such that  $\langle 1 \rangle_{1+b} = 1$ .

## Theorem (Okounkov - 1997)

*If  $N$  is a positive integer,  $1 + b$  is a positive real number, and  $\theta$  is an integer partition of  $2n$ , then*

$$\langle J_\theta(\boldsymbol{\lambda}, 1+b) \rangle_{1+b} = J_\theta(\mathbf{1}_N, 1+b) [p_{[2^n]}] J_\theta,$$

*where  $\mathbf{1}_N = (1, \dots, 1, 0, 0, \dots)$  consists of  $N$  leading 1's followed by 0's.*

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with  $c_{1+b}$  chosen such that  $\langle 1 \rangle_{1+b} = 1$ .

$$M(N, \mathbf{y}, z, \mathbf{r}; b) = r_0 N + (1+b) \sum_{j \geq 1} r_j \frac{\partial}{\partial y_j} \ln \left\langle e^{\frac{1}{1+b} \sum_{k \geq 1} \frac{1}{k} y_k p_k(\boldsymbol{\lambda}) \sqrt{z^k}} \right\rangle_{1+b}$$

◀ Return

# $b$ is ubiquitous

## The many lives of $b$

	$b = 0$		$b = 1$
Hypermaps	Orientable	?	Locally Orientable
Symmetric Functions	$s_\theta$	$J_\theta(b)$	$Z_\theta$
Matrix Integrals	Hermitian	?	Real Symmetric
Moduli Spaces	over $\mathbb{C}$	?	over $\mathbb{R}$
Matching Systems	Bipartite	?	All