Math 249, Winter 2013 Assignment 8

Not to be handed in.

- 1. Let G be a graph, and let M be a matching.
 - (a) Let M' be another matching, and let H be the spanning subgraph of G with $E(H) = M \cup M'$. Prove that every component of H is one of the following: (i) an isolated vertex; (ii) a single edge that lies in $M \cap M'$; (iii) an even cycle with edges alternately in M and M'; or (iv) a path with edges alternately in M and M'.
 - (b) If M is not a maximum matching, prove that an augmenting path exists.
- 2. The $a \times b$ Knight's Move graph, denoted KM(a,b) has vertex set $[a] \times [b]$ and $\{(s,t),(u,v)\}$ is an edge iff $(s-u)^2 + (t-v)^2 = 5$. Note that KM(a,b) is a subgraph of KM(a',b') if $a \le a'$ and $b \le b'$.
 - (a) Use the Bipartite Matching Algorithm to find a perfect matching in the following graphs:
 - (i) KM(3,4)
 - (ii) KM(3,6)
 - (iii) KM(5,6)

For (i), start with a perfect matching of KM(2,4). For (ii) and (iii), use the previous solution as a starting point.

- (b) Prove that KM(a, 2b) has a perfect matching, for $a \geq 3$, $b \geq 2$. [Hint: Hall's Theorem probably won't help. Use your solutions from (a) as the building blocks for a solution to (b).]
- $3.\ \,$ Find a 3-regular graph with no perfect matching. Prove that your answer is correct.

[*Hint*: If a graph has a bridge, then that bridge must either be in every perfect matching, or not be in any perfect matching. Why? If a graph has multiple bridges in awkward places, this can prevent the graph from having perfect matching.]

4. Let G be a graph, and let A be its adjacency matrix. If det(A) is odd, prove that G has a perfect matching.

[Hint: This is quite hard, but not impossible. You'll need the Leibniz formula for det(A):

$$\det(A) = \sum_{\sigma \in S_n} \operatorname{sgn}(\sigma) A_{1\sigma(1)} A_{2\sigma(2)} \cdots A_{n\sigma(n)}.$$

Here A is an $n \times n$ matrix, S_n is the set of all permutations of [n], and $\operatorname{sgn}(\sigma) = \pm 1$ is the sign of σ . Some of the permutations in S_n can be thought of as perfect matchings in K_n . Which ones correspond to matchings in G? How does this relate to the Leibniz formula? Once you understand the answers to these questions, let m(G) be the number of perfect matchings of G. Prove that $\det(A) \equiv m(G) \pmod{2}$.