### MATH 249, WINTER 2017, ASSIGNMENT 5

This assignment is due Friday February 17 at 3pm. It should be submitted using crowd-mark. The crowdmark instructions are the same as before.

## Questions.

- (1) (3 points BONUS) Draw the Petersen graph on a torus (with no edge crossings).
- (2) (5 points) Let G be a planar graph with n vertices and m edges.
  - (a) Let T be a spanning tree of G. Prove that the edges of  $G^*$  which do not correspond to edges of T form a spanning tree of  $G^*$ .
  - (b) Let  $[I_{n-1}|A]$  be an  $n-1 \times m$  matrix. Suppose that the submatrix formed by the columns indexed by  $J \subset \{1, 2, ..., m\}$  with |J| = n-1 has nonzero determinant. Prove that the columns not indexed by J in  $[-A^T|I_{m-n+1}]$  also have nonzero determinant.

In fact you just proved the same thing twice because if we take the signed incidence matrix of G and remove a row and choose the order on the edges so that the first n-1 columns form a spanning tree then we can row reduce this matrix to the form  $[I_{n-1}|A]$ . In this case the signed incidence matrix of  $G^*$  with one row removed is row-equivalent to  $[-A^T|I_{m-n+1}]$ .

- (3) (5 points) The girth of a graph is the length of a shortest cycle.
  - (a) Prove that every planar graph with girth at least 4 is 4-colourable (don't use the 4-colour theorem).
  - (b) What girth is required to guarantee that a planar graph with that girth is 3-colourable using analogous techniques to the previous part and what we did in class?

# (4) **(6 points)**

- (a) Let G be a connected planar graph which is isomorphic to its dual. Prove that |E(G)| = 2|V(G)| 2.
- (b) Find a connected planar graph on 6 vertices which is isomorphic to its dual.
- (c) How many non-isomorphic connected planar multigraphs on 4 vertices which are isomorphic to their duals are there?

### (5) (4 points)

- (a) Show that if two opposite corner squares of a chessboard (chessboards are 8 × 8) are removed then the resulting board cannot be covered with 31 dominoes each of which covers two adjacent squares.
- (b) What did the previous part have to do with things we discussed in class?

#### (6) **(5 points)**

- (a) Let G be a bipartite graph, and let  $\Delta$  be the maximum degree of any vertex in G. Show that G has a matching with at least  $|E(G)|/\Delta$  edges.
- (b) Give an example of a non-bipartite graph for which this lower bound on the size of a matching is not true.