

# Introduction to Quantum Information Processing

QIC 710 / CS 768 / PH 767 / CO 681 / AM 871

## Lecture 11 (2017)

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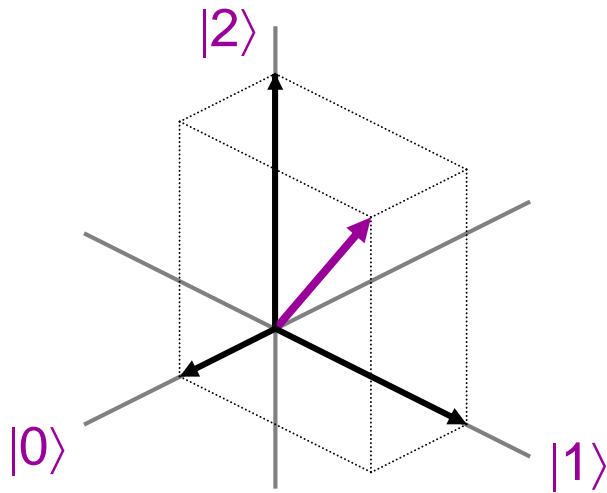
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# General measurements and POVMS

(POVM = Positive Operator Valued Measure)

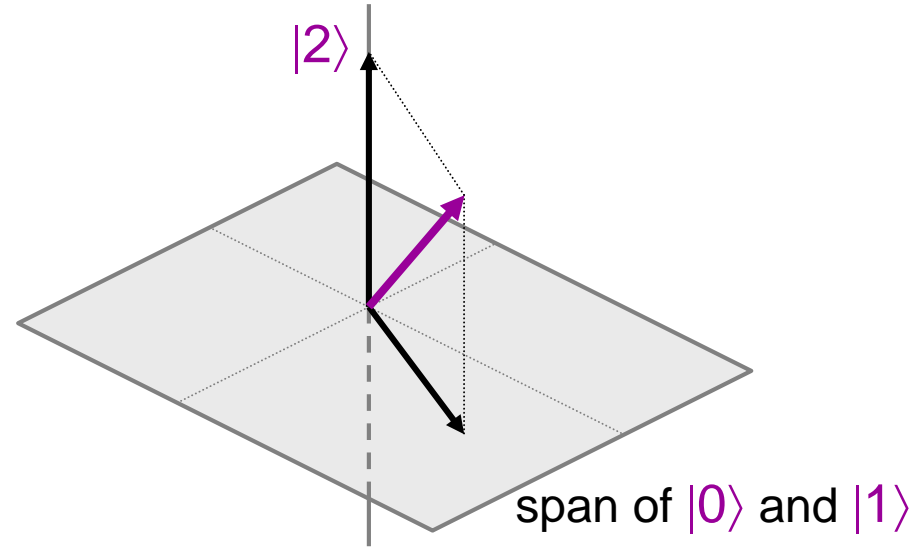
# Prelude: projective measurements



$$P_0 = |0\rangle\langle 0|$$

$$P_1 = |1\rangle\langle 1|$$

$$P_2 = |2\rangle\langle 2|$$



$$P_1 = |0\rangle\langle 0| + |1\rangle\langle 1|$$

$$P_2 = |2\rangle\langle 2|$$

In both cases, there is a complete set of mutually orthogonal projectors:

$$\sum_j P_j = I \quad \text{and} \quad P_i P_j = 0$$

The probability of outcome  $j$  is  $\langle \psi | P_j^\dagger P_j | \psi \rangle = \text{Tr}(|\psi\rangle\langle \psi | P_j^\dagger P_j)$  using  $\text{Tr}(AB) = \text{Tr}(BA)$

The collapsed state is the projected vector, but normalized

# General measurements (1)

Let  $A_1, A_2, \dots, A_m$  be *any* matrices satisfying  $\sum_{j=1}^m A_j^\dagger A_j = I$

Corresponding **measurement** is a stochastic operation on  $\rho$

that, with probability  $\text{Tr}(A_j \rho A_j^\dagger)$ , produces outcome:

$$\left\{ \begin{array}{l} \mathbf{j} \text{ (classical information)} \\ \frac{A_j \rho A_j^\dagger}{\text{Tr}(A_j \rho A_j^\dagger)} \text{ (the collapsed quantum state)} \end{array} \right.$$

**Example 1:** ( $A_j = |\phi_j\rangle\langle\phi_j|$ ) (rank-1 orthogonal projectors)

Consistent with our first definition of measurements

**Question:** what if we do the above but don't look at  $\mathbf{j}$ ?

# General measurements (2)

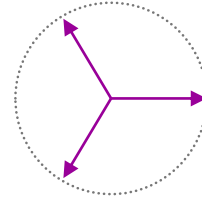
When  $A_j = |\phi_j\rangle\langle\phi_j|$  are orthogonal projectors and  $\rho = |\psi\rangle\langle\psi|$ ,

$$\begin{aligned}\text{Tr}(A_j\rho A_j^\dagger) &= \text{Tr}|\phi_j\rangle\langle\phi_j|\psi\rangle\langle\psi|\phi_j\rangle\langle\phi_j| \\ &= \langle\phi_j|\psi\rangle\langle\psi|\phi_j\rangle\langle\phi_j|\phi_j\rangle \\ &= |\langle\phi_j|\psi\rangle|^2\end{aligned}$$

Moreover, 
$$\frac{A_j\rho A_j^\dagger}{\text{Tr}(A_j\rho A_j^\dagger)} = \frac{|\phi_j\rangle\langle\phi_j|\psi\rangle\langle\psi|\phi_j\rangle\langle\phi_j|}{|\langle\phi_j|\psi\rangle|^2} = |\phi_j\rangle\langle\phi_j|$$

# General measurements (3)

Example 3 (trine state measurement):



$$\text{Let } |\phi_0\rangle = |0\rangle, \quad |\phi_1\rangle = -\frac{1}{2}|0\rangle + \frac{\sqrt{3}}{2}|1\rangle, \quad |\phi_2\rangle = -\frac{1}{2}|0\rangle - \frac{\sqrt{3}}{2}|1\rangle$$

$$\text{Define } A_0 = \sqrt{\frac{2}{3}}|\phi_0\rangle\langle\phi_0| = \sqrt{\frac{2}{3}}\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$A_1 = \sqrt{\frac{2}{3}}|\phi_1\rangle\langle\phi_1| = \frac{1}{4}\begin{pmatrix} \sqrt{2/3} & -\sqrt{2} \\ -\sqrt{2} & \sqrt{6} \end{pmatrix}$$

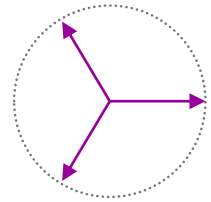
$$A_2 = \sqrt{\frac{2}{3}}|\phi_2\rangle\langle\phi_2| = \frac{1}{4}\begin{pmatrix} \sqrt{2/3} & \sqrt{2} \\ \sqrt{2} & \sqrt{6} \end{pmatrix}$$

$$\text{Then } A_0^\dagger A_0 + A_1^\dagger A_1 + A_2^\dagger A_2 = I.$$

If the input itself is an unknown trine state  $|\phi_k\rangle\langle\phi_k|$ , then the probability that classical outcome is  $k$  is  $2/3 = 0.6666\dots$

# General measurements (3)

**Question:** Are there states the trine measurement can't distinguish?



$$|\phi_0\rangle = |0\rangle, \quad |\phi_1\rangle = -\frac{1}{2}|0\rangle + \frac{\sqrt{3}}{2}|1\rangle, \quad |\phi_2\rangle = -\frac{1}{2}|0\rangle - \frac{\sqrt{3}}{2}|1\rangle$$

$$A_0 = \sqrt{\frac{2}{3}}|\phi_0\rangle\langle\phi_0| = \sqrt{\frac{2}{3}}\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$$

$$A_1 = \sqrt{\frac{2}{3}}|\phi_1\rangle\langle\phi_1| = \frac{1}{4}\begin{pmatrix} \sqrt{2/3} & -\sqrt{2} \\ -\sqrt{2} & \sqrt{6} \end{pmatrix}$$

$$A_2 = \sqrt{\frac{2}{3}}|\phi_2\rangle\langle\phi_2| = \frac{1}{4}\begin{pmatrix} \sqrt{2/3} & \sqrt{2} \\ \sqrt{2} & \sqrt{6} \end{pmatrix}$$

**Answer:**  $|i\rangle = \frac{|0\rangle + i|1\rangle}{\sqrt{2}}$ ,  $|-i\rangle = \frac{|0\rangle - i|1\rangle}{\sqrt{2}}$  ( $Y$ -eigenstates)

The trine measurement is **not informationally complete!**

# General measurements (4)

Often measurements arise in contexts where we only care about the classical part of the outcome (not the residual quantum state)

The probability of outcome  $j$  is  $\text{Tr}(A_j \rho A_j^\dagger) = \text{Tr}(\rho A_j^\dagger A_j)$

***Simplified definition of such measurements***

Let  $E_1, E_2, \dots, E_m$  be positive semidefinite and with  $\sum_{j=1}^m E_j = I$

The probability of outcome  $j$  is  $\text{Tr}(\rho E_j)$ .

Called a ***POVM (Positive Operator-Valued Measure)***

**It is a measure valued in positive (-semidefinite) operators.**



# Informationally-complete POVMs

A POVM  $E_1, E_2, \dots, E_m$  is **informationally complete** if  $\text{span}_{\mathbb{R}}(E_1, E_2, \dots, E_m) = \text{all Hermitian } d \times d \text{ matrices.}$

Such POVMs can distinguish **any** states.

**Example:** Informationally complete POVMs such that  $\text{rank}(E_j) = 1$  for each  $j$ , i.e. such that

$E_1 = \alpha_1 |\phi_1\rangle\langle\phi_1|, E_2 = \alpha_2 |\phi_2\rangle\langle\phi_2|, \dots, E_m = \alpha_m |\phi_m\rangle\langle\phi_m|,$   
are sometimes called **tight frames**.

**(Very hard) question:** Do informationally-complete POVMs exist with  $\text{rank}(E_j) = 1$  for every  $j$  and  $\text{Tr}(E_i E_j) = \text{constant}$  for  $i \neq j$ ?

**Answer:** Apparently yes (no proof yet), known as SIC-POVMs (Symmetric Informationally Complete POVMs) 9

# “Mother of all operations”

Let  $A_{1,1}, A_{1,2}, \dots, A_{1,k_1}$   
 $A_{2,1}, A_{2,2}, \dots, A_{2,k_2}$   
 $A_{m,1}, A_{m,2}, \dots, A_{m,k_m}$  satisfy  $\sum_{j=1}^m \sum_{i=1}^{k_m} A_{j,i}^\dagger A_{j,i} = I$

Then there is a quantum operation that, on input  $\rho$ , produces with probability  $\sum_{i=1}^{k_m} A_{j,i} \rho A_{j,i}^\dagger$  the state:

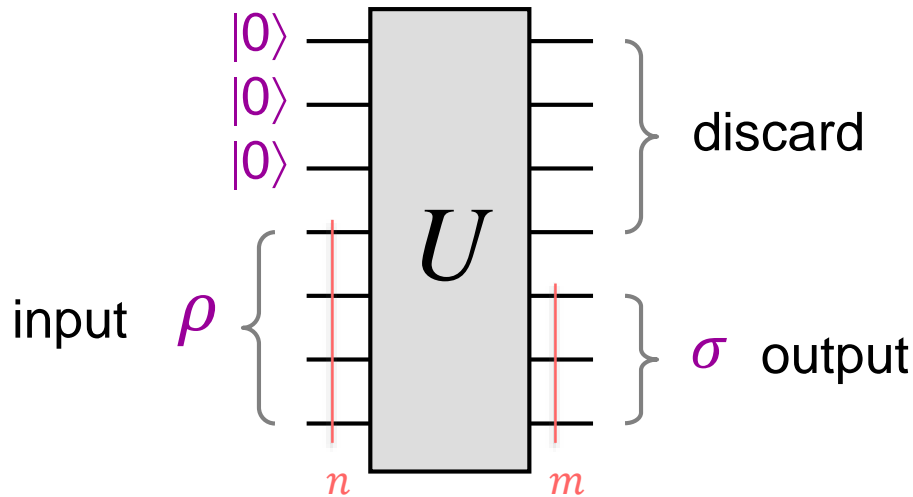
$$\left\{ \begin{array}{l} \mathbf{j} \text{ (classical information)} \\ \frac{\sum_{i=1}^{k_m} A_{j,i} \rho A_{j,i}^\dagger}{\sum_{i=1}^{k_m} \text{Tr} (A_{j,i} \rho A_{j,i}^\dagger)} \text{ (the collapsed quantum state)} \end{array} \right.$$

Also known as **quantum instrument** or **heralded quantum channel**

# Simulations among operations

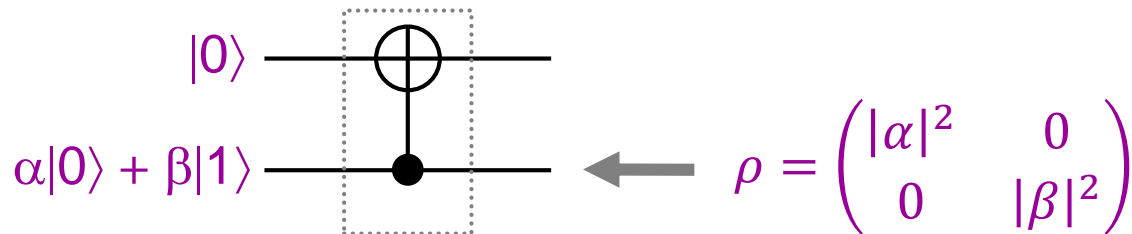
# Simulations among operations (1)

**Theorem 1:** any *quantum operation* can be simulated by applying a unitary operation on a larger quantum system:



This specification of a quantum operation is called the **Stinespring** form, or **isometric extension**

**Example:** decoherence



# Simulations among operations (2)

## Proof of Theorem 1:

Let  $A_1, A_2, \dots, A_{2^k}$  be any  $2^m \times 2^n$  matrices such that

$$\sum_{j=1}^{2^k} A_j^\dagger A_j = I$$

This defines a mapping from  $m$  qubits to  $n$  qubits:

$$\rho \mapsto \sum_{j=1}^{2^k} A_j \rho A_j^\dagger$$

This specification of the quantum operation is called the **Kraus** form

# Simulations among operations (3)

Set  $V = \begin{bmatrix} A_1 \\ A_2 \\ \vdots \\ A_{2^k} \end{bmatrix}$

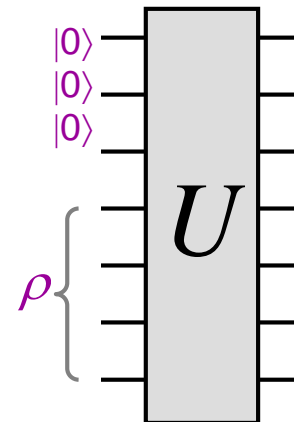
Since  $V^\dagger V = \begin{bmatrix} A_1^\dagger & A_2^\dagger & \cdots & A_{2^k}^\dagger \end{bmatrix} \begin{bmatrix} A_1 \\ A_2 \\ \vdots \\ A_{2^k} \end{bmatrix} = I$   
 the columns of  $V$  are **orthonormal**

Let  $U$  be any unitary matrix with first  $2^n$  columns from  $V$

$$U = [V \mid W]$$

$U$  is a  $2^{m+k} \times 2^{m+k}$  matrix (and its columns partition into  $2^{m-n+k}$  blocks of size  $2^n$ )

Now, consider the circuit:



# Simulations among operations (4)

The output state of the circuit is  $U(|00 \dots 0\rangle\langle 00 \dots 0| \otimes \rho)U^\dagger$

$$= \begin{bmatrix} A_1 & & & \\ & \vdots & & \\ & & \ddots & \\ & & & A_{2^k} \end{bmatrix} W \begin{bmatrix} \rho & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix} \begin{bmatrix} A_1^\dagger & A_1^\dagger & \dots & A_{2^k}^\dagger \\ \hline & & W^\dagger & \end{bmatrix}$$

$$= \begin{bmatrix} A_1 \rho & 0 & \dots & 0 \\ A_2 \rho & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ A_{2^k} \rho & 0 & \dots & 0 \end{bmatrix} \begin{bmatrix} A_1^\dagger & A_1^\dagger & \dots & A_{2^k}^\dagger \\ \hline & & W^\dagger & \end{bmatrix}$$

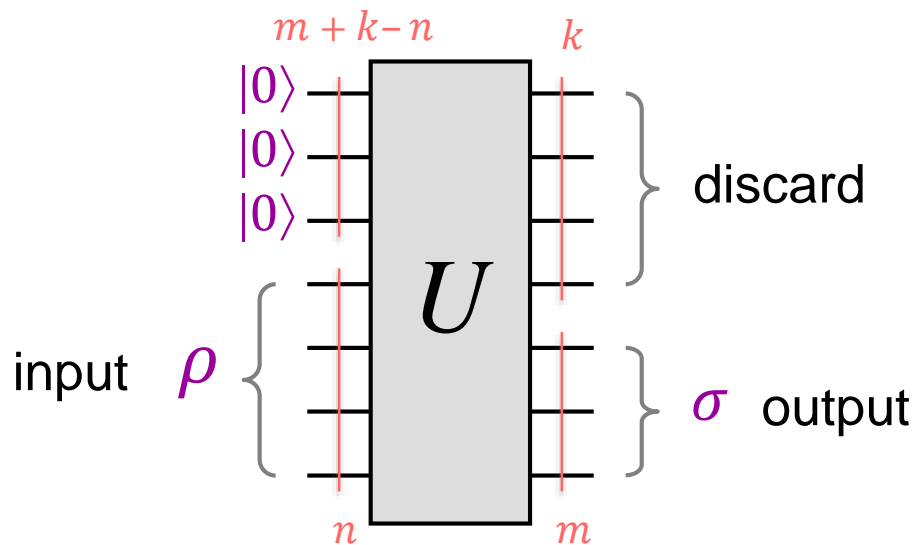
$$= \begin{bmatrix} A_1 \rho A_1^\dagger & A_1 \rho A_2^\dagger & \dots & A_1 \rho A_{2^k}^\dagger \\ A_2 \rho A_1^\dagger & A_2 \rho A_2^\dagger & \dots & A_2 \rho A_{2^k}^\dagger \\ \vdots & \vdots & \ddots & \vdots \\ A_{2^k} \rho A_1^\dagger & A_{2^k} \rho A_2^\dagger & \dots & A_{2^k} \rho A_{2^k}^\dagger \end{bmatrix}$$

# Simulations among operations (5)

Tracing out the high-order  $k$  qubits of this state yields

$$A_1 \rho A_1^\dagger + A_2 \rho A_2^\dagger + \cdots + A_{2^k} \rho A_{2^k}^\dagger$$

exactly the output of mapping that we want to simulate

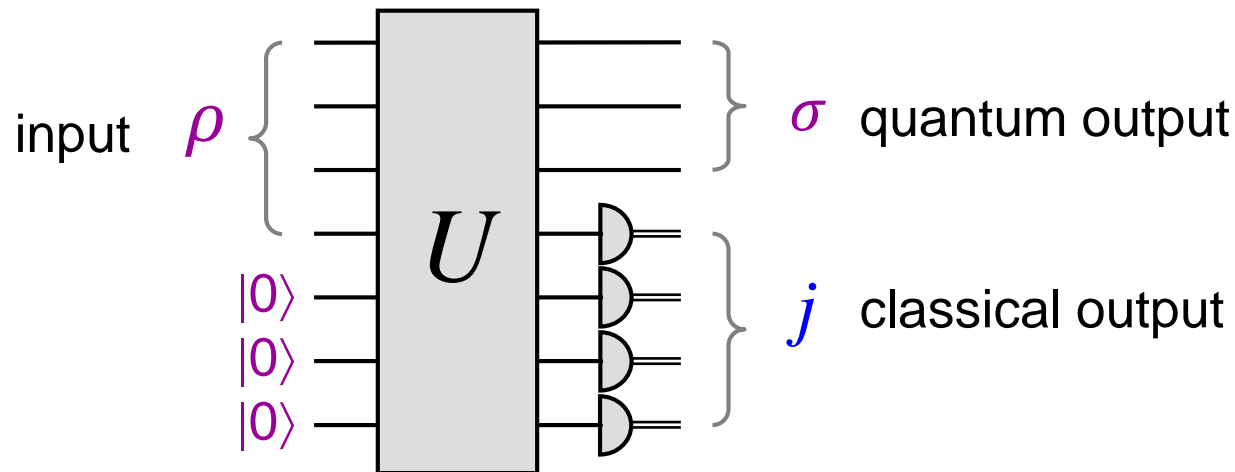


**Note:** this approach is *not always optimal* in the number of ancillary qubits used—there are more efficient methods



# Simulations among operations (6)

**Theorem 2:** *any measurement* can also be simulated by applying a unitary operation on a larger quantum system and then measuring:



This is the same diagram as for Theorem 1 (drawn with the extra qubits at the bottom) but where the “discarded” qubits are measured and part of the output