

# **CO350 Linear Programming**

## **Chapter 9: The Revised Simplex Method**

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## Example of unboundedness

Solve the LP using revised simplex method with smallest-subscript rules.

$$\begin{array}{ll}
 \max & 3x_1 - 2x_2 - 3x_3 \\
 \text{s.t.} & x_1 - x_2 - x_3 \leq 1 \\
 (P) & 7x_1 - 8x_2 - 11x_3 \leq 2 \\
 & 2x_1 - 2x_2 - 3x_3 \leq 1 \\
 & x_1, x_2, x_3 \geq 0
 \end{array}$$

Adding slack variables  $x_4, x_5, x_6$  gives

$$\begin{array}{ll}
 \max & c^T x \\
 (P') & \text{s.t.} \quad Ax = b \\
 & x \geq 0
 \end{array}$$

where

$$A = \begin{bmatrix} 1 & -1 & -1 & 1 & 0 & 0 \\ 7 & -8 & -11 & 0 & 1 & 0 \\ 2 & -2 & -3 & 0 & 0 & 1 \end{bmatrix}, \quad b = \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix}$$

and

$$c = [3, -2, -3, 0, 0, 0]^T$$

Start from feasible basis  $B = \{4, 5, 6\}$ .

Iteration 1:

$$B = \{4, 5, 6\}, x_B^* = \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix}, A_B = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}.$$

$$\text{Compute } \bar{c}_1 = c_1 - A_1^T y = 3 - [1 \ 7 \ 2]y = 3 > 0.$$

$x_1$  enters.

$$\text{Solve } A_B d = A_1 \text{ to get } d = \begin{bmatrix} 1 \\ 7 \\ 2 \end{bmatrix}.$$

$$t = \min\{\frac{1}{1}, \frac{2}{7}, \frac{1}{2}\} = \frac{2}{7}. x_5 \text{ leaves.}$$

$$\text{Update } x_1^* = t = \frac{2}{7}, x_4^* = 1 - (1)(\frac{2}{7}) = \frac{5}{7}, x_6^* = 1 - (2)(\frac{2}{7}) = \frac{3}{7}.$$

Iteration 2:

$$B = \{1, 4, 6\}, x_B^* = \begin{bmatrix} 2/7 \\ 5/7 \\ 3/7 \end{bmatrix}, A_B = \begin{bmatrix} 1 & 1 & 0 \\ 7 & 0 & 0 \\ 2 & 0 & 1 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 3 \\ 0 \\ 0 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 0 \\ 3/7 \\ 0 \end{bmatrix}.$$

Compute  $\bar{c}_2 = c_2 - A_2^T y = -2 - [-1 \quad -8 \quad -2]y = 10/7 > 0$ .  
 $x_2$  enters.

$$\text{Solve } A_B d = A_2 \text{ to get } d = \begin{bmatrix} -8/7 \\ 1/7 \\ 2/7 \end{bmatrix}.$$

$$t = \min\{-, \frac{5/7}{1/7}, \frac{3/7}{2/7}\} = \frac{3}{2}. \quad x_6 \text{ leaves.}$$

$$\text{Update } x_2^* = t = \frac{3}{2}, \quad x_1^* = \frac{2}{7} - (-\frac{8}{7})(\frac{3}{2}) = 2, \quad x_4^* = \frac{5}{7} - (\frac{1}{7})(\frac{3}{2}) = \frac{1}{2}.$$

Iteration 3:

$$B = \{1, 2, 4\}, \quad x_B^* = \begin{bmatrix} 2 \\ 3/2 \\ 1/2 \end{bmatrix}, \quad A_B = \begin{bmatrix} 1 & -1 & 1 \\ 7 & -8 & 0 \\ 2 & -2 & 0 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 3 \\ -2 \\ 0 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 0 \\ -1 \\ 5 \end{bmatrix}.$$

Compute  $\bar{c}_3 = c_3 - A_3^T y = -3 - [-1 \quad -11 \quad -3]y = 1 > 0$ .  
 $x_3$  enters.

$$\text{Solve } A_B d = A_3 \text{ to get } d = \begin{bmatrix} -1 \\ 1/2 \\ 1/2 \end{bmatrix}.$$

$$t = \min\{-, \frac{3/2}{1/2}, \frac{1/2}{1/2}\} = 1. \quad x_4 \text{ leaves.}$$

$$\text{Update } x_3^* = t = 1, \quad x_1^* = 2 - (-1)(1) = 3, \quad x_2^* = \frac{3}{2} - (\frac{1}{2})(1) = 1.$$

Iteration 4:

$$B = \{1, 2, 3\}, \quad x_B^* = \begin{bmatrix} 3 \\ 1 \\ 1 \end{bmatrix}, \quad A_B = \begin{bmatrix} 1 & -1 & -1 \\ 7 & -8 & -11 \\ 2 & -2 & -3 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 3 \\ -2 \\ -3 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 2 \\ -1 \\ 4 \end{bmatrix}.$$

$$\text{Compute } \bar{c}_4 = c_4 - A_4^T y = 0 - [1 \ 0 \ 0]y = -2 \leq 0.$$

$$\text{Compute } \bar{c}_5 = c_5 - A_5^T y = 0 - [0 \ 1 \ 0]y = 1 > 0.$$

$x_5$  enters.

$$\text{Solve } A_B d = A_5 \text{ to get } d = \begin{bmatrix} -1 \\ -1 \\ 0 \end{bmatrix}.$$

$$d \leq 0 \implies (P') \text{ is unbounded.}$$

Hence  $(P)$  is unbounded.

Proof of unboundedness

Let  $x_5(t) = t$ ,  $x_4(t) = x_6(t) = 0$ , and  $x_B(t) = x_B^* - td$

$$\text{i.e., } [x_1(t), x_2(t), x_3(t)]^T = [3 + t, 1 + t, 1]^T$$

Check:  $x(t)$  is feasible for  $(P')$  when  $t \geq 0$ .

$$\text{Objective value: } c^T x(t) = 4 + t \rightarrow \infty \text{ as } t \rightarrow \infty$$

### Example of revised two-phase method (see also §9.3)

Solve the LP using revised two-phase method with smallest-subscript rules.

$$\begin{array}{ll}
 \max & (z =) \quad x_1 \quad \quad \quad + 3x_3 \\
 (P) \quad \text{s.t.} & x_1 + 3x_2 - x_3 + 2x_4 = 5 \\
 & x_1 - 3x_2 + 5x_3 - 4x_4 = -1 \\
 & x_1, \quad x_2, \quad x_3, \quad x_4 \geq 0
 \end{array}$$

Multiply the second equation by  $-1$  and add artificial variables  $x_5, x_6$  to get the auxiliary problem

$$\begin{array}{ll}
 \max & g^T x \\
 (A) \quad \text{s.t.} & Dx = f \\
 & x \geq 0
 \end{array}$$

where

$$D = \begin{bmatrix} 1 & 3 & -1 & 2 & 1 & 0 \\ -1 & 3 & -5 & 4 & 0 & 1 \end{bmatrix}, \quad f = \begin{bmatrix} 5 \\ 1 \end{bmatrix}$$

and

$$g = [0, 0, 0, 0 - 1, -1]^T$$

Iteration 1:

$$B = \{5, 6\}, x_B^* = \begin{bmatrix} 5 \\ 1 \end{bmatrix}, D_B = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

$$\text{Solve } D_B^T y = g_B = \begin{bmatrix} -1 \\ -1 \end{bmatrix} \text{ to get } y = \begin{bmatrix} -1 \\ -1 \end{bmatrix}.$$

$$\text{Compute } \bar{g}_1 = g_1 - D_1^T y = 0 - [1 \quad -1]y = 0 \leq 0.$$

$$\text{Compute } \bar{g}_2 = g_2 - D_2^T y = 0 - [3 \quad 3]y = 6 > 0. \quad x_2 \text{ enters.}$$

$$\text{Solve } D_B d = D_2 \text{ to get } d = \begin{bmatrix} 3 \\ 3 \end{bmatrix}.$$

$$t = \min\{\frac{5}{3}, \frac{1}{3}\} = \frac{1}{3}. \quad x_6 \text{ leaves. Drop artificial } x_6.$$

$$\text{Update } x_2^* = t = \frac{1}{3}, x_5^* = 5 - (3)(\frac{1}{3}) = 4.$$

Iteration 2:

$$B = \{2, 5\}, x_B^* = \begin{bmatrix} 1/3 \\ 4 \end{bmatrix}, D_B = \begin{bmatrix} 3 & 1 \\ 3 & 0 \end{bmatrix}.$$

$$\text{Solve } D_B^T y = g_B = \begin{bmatrix} 0 \\ -1 \end{bmatrix} \text{ to get } y = \begin{bmatrix} -1 \\ 1 \end{bmatrix}.$$

$$\text{Compute } \bar{g}_1 = g_1 - D_1^T y = 0 - [1 \quad -1]y = 2 > 0. \quad x_1 \text{ enters.}$$

$$\text{Solve } D_B d = D_1 \text{ to get } d = \begin{bmatrix} -1/3 \\ 2 \end{bmatrix}.$$

$$t = \min\{-, \frac{4}{2}\} = \frac{4}{2}. \quad x_5 \text{ leaves. Drop artificial } x_5.$$

$$\text{Update } x_1^* = t = 2, x_2^* = \frac{1}{3} - (-\frac{1}{3})(2) = 1.$$

All artificial variables are dropped, hence  $B = \{1, 2\}$  is feasible for  $(P)$ .

We start Phase 2 on  $(P)$  with  $B = \{1, 2\}$ .

Iteration 1:

$$B = \{1, 2\}, x_B^* = \begin{bmatrix} 2 \\ 1 \end{bmatrix}, A_B = \begin{bmatrix} 1 & 3 \\ 1 & -3 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 1/2 \\ 1/2 \end{bmatrix}.$$

Compute  $\bar{c}_3 = c_3 - A_3^T y = 3 - [-1 \ 5]y = 1 > 0$ .  $x_3$  enters.

$$\text{Solve } A_B d = A_3 \text{ to get } d = \begin{bmatrix} 2 \\ -1 \end{bmatrix}.$$

$t = \min\{\frac{2}{2}, -\} = \frac{2}{2}$ .  $x_1$  leaves.

Update  $x_3^* = t = 1$ ,  $x_2^* = 1 - (-1)(1) = 2$ .



Iteration 2:

$$B = \{2, 3\}, x_B^* = \begin{bmatrix} 2 \\ 1 \end{bmatrix}, A_B = \begin{bmatrix} 3 & -1 \\ -3 & 5 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 0 \\ 3 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 3/4 \\ 3/4 \end{bmatrix}.$$

$$\text{Compute } \bar{c}_1 = c_1 - A_1^T y = 1 - [1 \ 1]y = -1/2 \leq 0.$$

$$\text{Compute } \bar{c}_4 = c_4 - A_4^T y = 0 - [2 \ -4]y = 3/2 > 0. \ x_4 \text{ enters.}$$

$$\text{Solve } A_B d = A_4 \text{ to get } d = \begin{bmatrix} 1/2 \\ -1/2 \end{bmatrix}.$$

$$t = \min\{\frac{2}{1/2}, -\} = \frac{2}{1/2}. \ x_2 \text{ leaves.}$$

$$\text{Update } x_4^* = t = 4, x_3^* = 1 - (-\frac{1}{2})(4) = 3.$$

Iteration 3:

$$B = \{3, 4\}, x_B^* = \begin{bmatrix} 3 \\ 4 \end{bmatrix}, A_B = \begin{bmatrix} -1 & 2 \\ 5 & -4 \end{bmatrix}.$$

$$\text{Solve } A_B^T y = c_B = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \text{ to get } y = \begin{bmatrix} 2 \\ 1 \end{bmatrix}.$$

$$\text{Compute } \bar{c}_1 = c_1 - A_1^T y = 1 - [1 \ 1]y = -2 \leq 0.$$

$$\text{Compute } \bar{c}_2 = c_2 - A_2^T y = 0 - [3 \ -3]y = -3 \leq 0.$$

$$\text{An optimal solution is } x^* = [0, 0, 3, 4]^T.$$