

CONTINUOUS-TIME ←  
 "QUANTUM WALK + MIXING"  
 (ALGORITHMIC LENS)

OPEN PROBLEMS IN  
 ALGEBRAIC COMBINATORICS

X GRAPH,  $L(X) =$  Laplacian matrix associated to X

ex:  $K_2$

$$L = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} = I - \sigma_x$$

$$= 0 \times \frac{J}{2} + 2 \times (I - \frac{J}{2})$$

RANDOM WALK:

$$e^{-tL(X)} = \frac{J}{2} + e^{-2t} (I - \frac{J}{2})$$

converge  $\xrightarrow{t \rightarrow \infty}$   $\frac{J}{2} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{pmatrix}$  uniform mixing

QUANTUM WALK: [FARHI-GUTMANN]

$$e^{-itL(X)} = \frac{J}{2} + e^{-2it} (I - \frac{J}{2}) = \begin{pmatrix} \cos t & -i \sin t \\ -i \sin t & \cos t \end{pmatrix}$$

$t = \pi/4$   $\begin{pmatrix} 1/\sqrt{2} & -i/\sqrt{2} \\ -i/\sqrt{2} & 1/\sqrt{2} \end{pmatrix}$

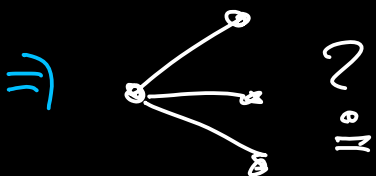
exact  $\downarrow$   $\rho_{ub} = \begin{pmatrix} 1/2 \\ 1/2 \end{pmatrix}$

X regular:

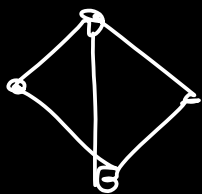
$$e^{-itL(X)} \cong e^{-itA(X)}$$

QUESTION: NON-REGULAR X (FAMILY)

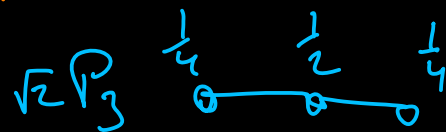
WITH LAPLACIAN UNIFORM MIXING



non?



$K_4 \setminus e$   
 ORPHAN



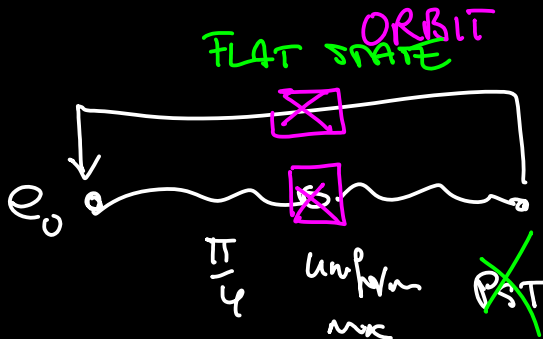
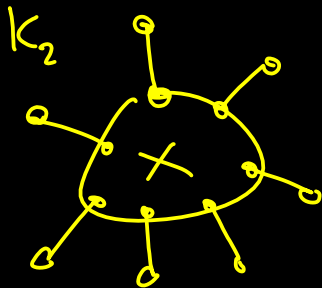
[ZHAN]

QUESTION:

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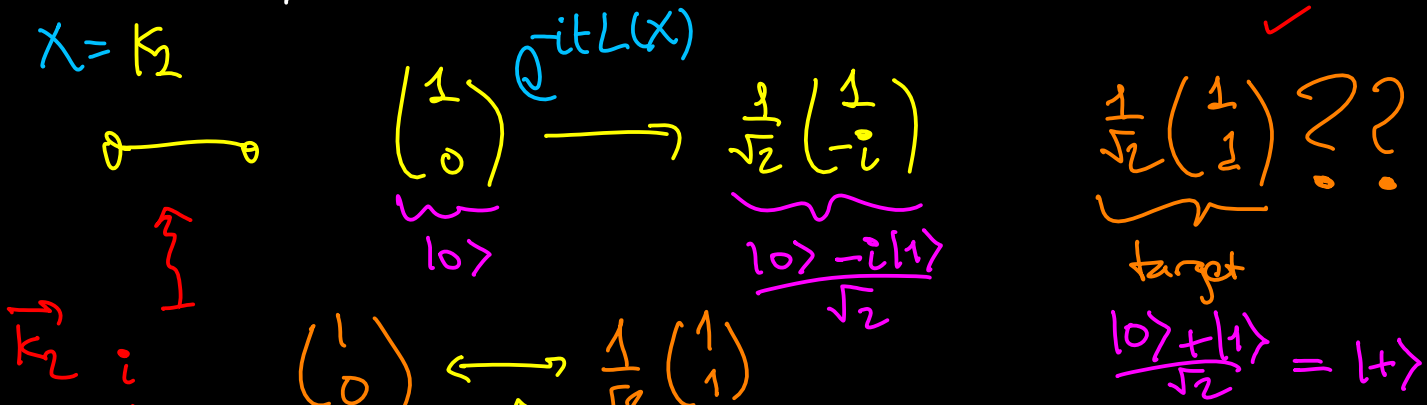
(UNARY) GRAPH OPERATORS

THAT PRESERVE UNIFORM MIXING?



# [AHARONOV, TA-SHMA] STATE GENERATION

$X = K_2$



PROB DISTR  $P \in \mathcal{C}_{\mathcal{N}, \mathcal{N}}$

$|ψ\rangle = \sum_{a \in V(X)} \sqrt{P_a} |a\rangle$   
 $\sum_{a \in V(X)} P_a = 1$

$Q^{-it} \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$   
 $\begin{pmatrix} \cos^2 t & \\ & \sin^2 t \end{pmatrix} \leftarrow \begin{pmatrix} \cos t & \\ & -i \sin t \end{pmatrix}$

GENERALIZE  $K_2$ : 3 WAYS

①  $\mathcal{G}_N$

②  $K_N$

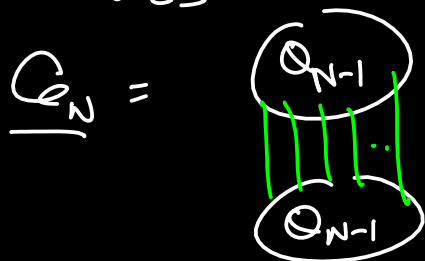
$\mathcal{P}_N$

$\mathcal{C}_N$

③  $K_{1,N}$

[MOORE, RUSSELL]

① N-CUBES



$\begin{pmatrix} A(Q_{N-1}) & I \\ I & A(Q_{N-1}) \end{pmatrix}$

oriented N-cube

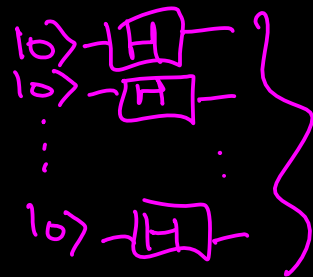
Commuting

$\sum_{\gamma} I \otimes \sigma_{\gamma} \otimes \dots \otimes I$

$Q^{-it} A(Q_N) = \left( Q^{-it} A(K_2) \right)^{\otimes N}$

$t = \pi/4$

$|ψ\rangle = \frac{1}{\sqrt{2^N}} \sum_{k=0}^{2^N-1} |k\rangle$



$\begin{pmatrix} 1/\sqrt{2} \\ -i/\sqrt{2} \end{pmatrix}^{\otimes N}$

MIXING SPEEDUP

lazy stochastic  

$$P = \frac{1}{2}I + \frac{1}{2} \frac{1}{N} A(G_N)$$

RANDOM WALK:

$\rightarrow \Omega(N \log_2 N)$

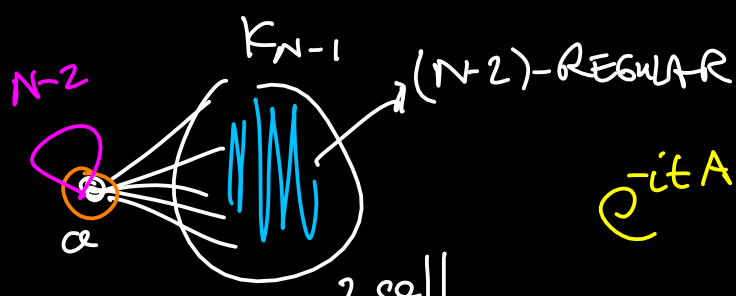
QUANTUM WALK:

QUESTION: FASTER SPEEDUP? EXPONENTIAL  $\sim$ ?  

$$e^{-it \frac{1}{N} A(G_N)}$$
  

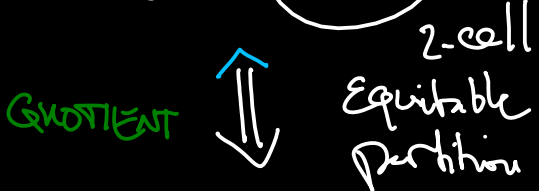
$$N \frac{\pi}{4} = \underline{O(N)}$$

2) CLIQUE:  $K_N$

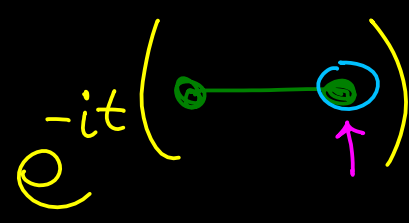
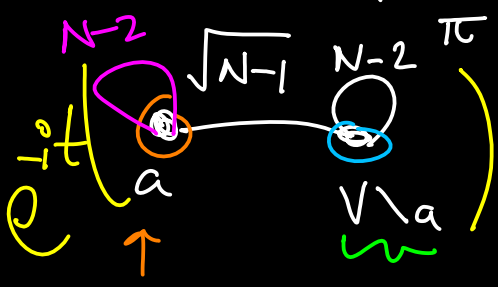


$$\frac{1}{\sqrt{N-1}} \begin{pmatrix} 0 \\ 1 \end{pmatrix} \approx \frac{1}{\sqrt{N}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$e^{-it A(K_N)} \approx I$$



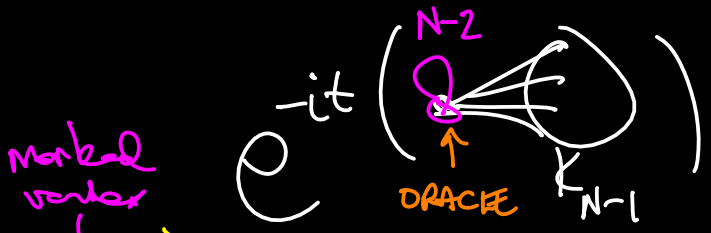
[GOODILL] SEDENTARY GRAPHS  
 QUESTION: ARE ALL EXPANDERS SEDENTARY?



$t = \sqrt{N-1} \frac{\pi}{2}$

How DO WE ESCAPE a?

PERFECT STATE TRANSFER  $t = \pi/2$



[FARHI-GUTMANN]

marked vertex  $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$

GROVER SEARCH  $\rightarrow \frac{1}{\sqrt{N}} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$

REVERSE MIXING

# 3) CLAWS

## MIXING GAME

Alice

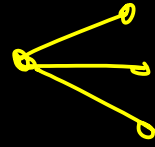
GRAPH

Bob

$K_2$ :



$K_{1,3}$



STATE GENERATION

1) CHOOSE

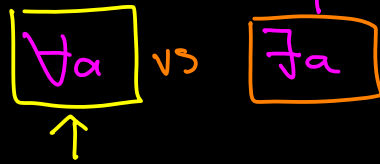
$P \in (0,1]^{V(X)}$   
(FULL SUPPORT)

2) CHOOSE WEIGHT

$W \in \mathbb{R}_+^{E(X)}$

BOB WINS IFF  $\exists t$

$$|\langle e_b, e^{-itA(X,W)} e_a \rangle| = \sqrt{p_b} \quad \forall b \in V(X)$$



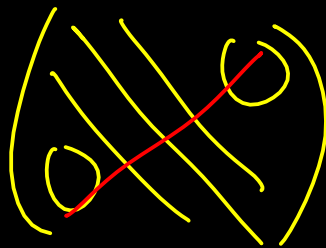
" $K_2$ "

CLAIM [CARLSON ET AL] BOB WINS ON "TREES"

$K_{1,N}$

$\underbrace{P}_N$ ?  $\rightarrow$  [KAY]

$t = \pi$



QUESTION: WHAT OTHER GRAPHS WHERE BOB WINS?

$\rightarrow 0$   
 $\rightarrow 2\mathbb{Z}+1$

$\lambda_1 > \lambda_2 > \dots > \lambda_n$

$\Downarrow$

Jacobi matrix unique

