

eg Let $|Y_1\rangle = |0\rangle$

$$|Y_2\rangle = \cos\theta |0\rangle + \sin\theta |1\rangle$$

$$|Y_3\rangle = \cos\theta |0\rangle - \sin\theta |1\rangle$$



s.t.

$$0.4 |Y_1\rangle\langle Y_1| + 0.3 |Y_2\rangle\langle Y_2| + 0.3 |Y_3\rangle\langle Y_3| = \frac{I}{2}$$

Let $P_x = 0.9 |Y_x\rangle\langle Y_x| + 0.1 \frac{I}{2}$.

$|Y_2\rangle, |Y_3\rangle$ are closer to one another than to $|Y_1\rangle$

Consider Q-box emitting P_x when input = x , $x=1,2,3$.

$$\text{Then } C(Q) = \max_{p_1, p_2, p_3} S\left(\sum_{x=1}^3 p_x P_x\right) - \sum_x p_x S(P_x)$$

Claim: optimal distribution is $p_1 = 0.4$, $p_2 = 0.3$, $p_3 = 0.3$.

Pf: Here $S(P_x)$ independent for x so we can focus on

maximizing $S\left(\sum_x p_x P_x\right)$, which is $S\left(\frac{I}{2}\right) = 1$

attained for $p_1 = 0.4$, $p_2 = p_3 = 0.3$.

NB Each P_x has eigenvalues 0.95, 0.05, so $S(P_x) \approx 0.28640$

$$\therefore C(Q) = 1 - 0.28640 \approx 0.71360$$

What is the code?

Consider all strings of 1,2,3 of length n with

1,2,3 occurring $\approx 40\%$, 30% , 30% of the time.

$H(0.4, 0.3, 0.3) \approx 1.571$ so there are $\approx 2^{1.571n}$ such strings

Strings are drawn with replacement, "almost" equiprobably,

$\approx 2^{(0.71-1)n}$ times, from this set of $2^{1.571n}$.

For $n = 100$, $\frac{2^{0.71n}}{2^{1.57n}} \approx 2^{-0.86n} \approx 2^{-86}$ so a very

sparse set has been chosen to be code words even for small n .

Say, $C_1 = 211232 \dots \dots \dots 1233 \leftarrow$ highly correlated

$C_2 = 312211 \dots \dots \dots 3123$

\vdots

$C_{2^{n(0.71-d)}} = 322111 \dots \dots \dots 3211$

So Bob has to distinguish

$\gamma_1 = p_2 \otimes p_1 \otimes p_1 \otimes p_2 \otimes p_3 \otimes p_2 \dots \otimes p_1 \otimes p_2 \otimes p_3 \otimes p_3$

$\gamma_2 = p_3 \otimes p_1 \otimes p_2 \otimes p_2 \otimes p_1 \otimes p_1 \dots \otimes p_3 \otimes p_1 \otimes p_2 \otimes p_3$

\vdots

$\gamma_{2^{n(0.71-d)}} = p_3 \otimes p_2 \otimes p_2 \otimes p_1 \otimes p_1 \otimes p_1 \dots \otimes p_3 \otimes p_2 \otimes p_1 \otimes p_1$
 π_2 (bracketed over the last four terms)
 π_1 on p_i (bracketed under the last four terms)

$\Pi_{C_{2^{n(0.71-d)}}} = \pi_1 \otimes \pi_2 \otimes \pi_3$
 on the appropriate sys. defined on $C_{2^{n(0.71-d)}}$

$\Gamma_{C_{2^{n(0.71-d)}}} = \Pi \Pi_{C_{2^{n(0.71-d)}}} \Pi$

$\Gamma = \sum_{i=1}^{2^{n(0.71-d)}} \Gamma_{C_i}$

$M_i = \Gamma^{-\frac{1}{2}} \Gamma_{C_i} \Gamma^{-\frac{1}{2}}$ \leftarrow note heavily entangling

Finally $C(\mathcal{Q}) \approx 0.71 \gg I_{acc}(\{p_x, p_y\})$ (cf. 0.585 for the bases)