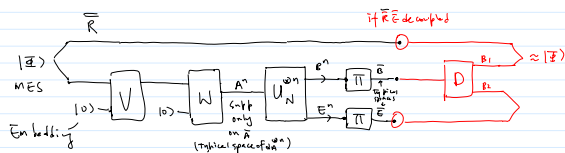
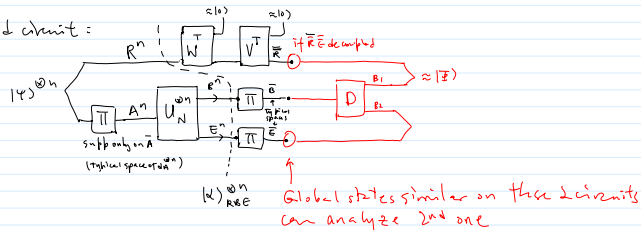


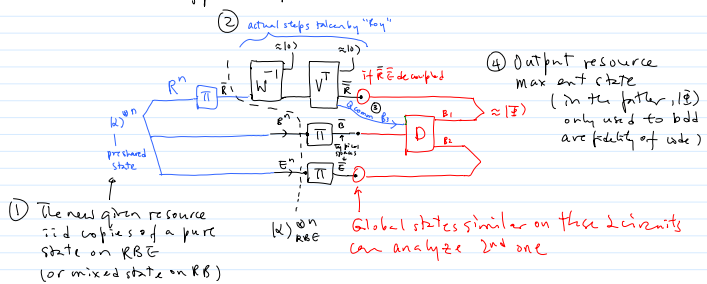
- Recall the coding protocol for the LSD theorem:



- A related circuit:

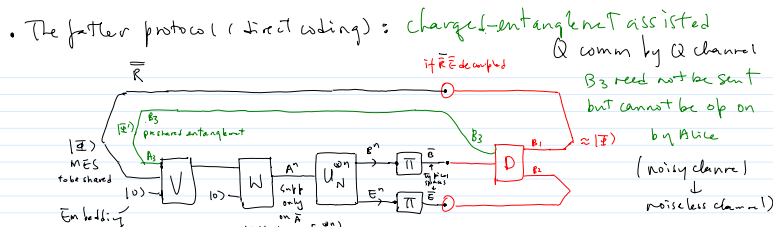


- A protocol "mother" suggested by the related circuit:

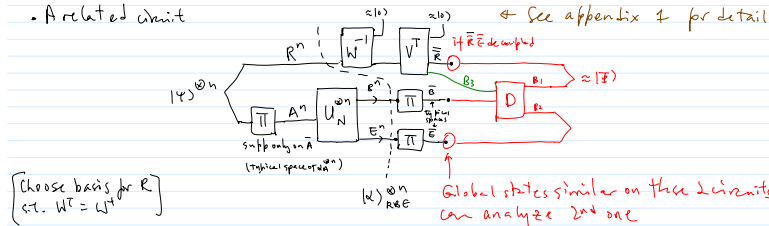


Mother: Mixed state entanglement purification (noisy ent → noiseless ent) using charged Q communication

Enough to analyze mother (father follows with $|\alpha\rangle = I \otimes U_N |\Phi\rangle$ same dim for \bar{R} & B_3).



- A related circuit



Goal: $\max \dim(\bar{R})$ (Q comm or distillable entanglement obtained)

$\min \dim(B_3)$ (Noiseless ent or Q communication spent)

While decoupling $\bar{R} \bar{E}$ (to guarantee the job done right)

Recall approx decoupling lemma: (Lecture 10 p 11)

$$\text{If } \left\| \rho_{\bar{R}\bar{E}} - \left(\frac{I}{2^{n_R}} \right)_{\bar{R}} \otimes \rho_{\bar{E}} \right\|_{\text{tr}} \leq \epsilon'$$

then $\exists |\Psi\rangle_{\bar{R}E, B_2}$ purifying $\rho_{\bar{R}\bar{E}}$

$$\text{s.t. } \left\| \text{tr}_{E, B_2} [|\Psi\rangle\langle\Psi|_{\bar{R}E, B_2}] - \left(\frac{I}{2^{n_R}} \right)_{\bar{R}} \otimes \rho_{\bar{E}} \right\|_{\text{tr}} \leq 2\sqrt{\epsilon'}$$

$$\text{So we bound } \left\| \rho_{\bar{R}\bar{E}} - \left(\frac{I}{2^{n_R}} \right)_{\bar{R}} \otimes \rho_{\bar{E}} \right\|_{\text{tr}}$$

Useful lemmas:

- (L1) Cauchy-Schwarz inequality

$$\|M\|_{\text{tr}}^2 \leq \text{rank}(M) \text{tr}(M^\dagger M) = \text{rank}(M) \|M\|_2^2$$

- (L2) $\text{tr}(M^2) = \text{tr}(\text{SWAP } M \otimes M)$

the operator taking $|ij\rangle$ to $|ji\rangle$ e.g. $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$

$$\text{Pf: } \text{tr}(M^2) = \sum_j \langle j | M I M | j \rangle$$

Alternative proof in appendix 2

$$\begin{aligned} &= \sum_{i,j} \langle j | M | i \rangle \langle i | M | j \rangle \\ &= \sum_{i,j} \langle j | \langle i | M \otimes M | i \rangle | j \rangle \end{aligned}$$

$$= \text{tr} \left(\underbrace{\sum_{ij} |ij\rangle \langle ij|}_{\text{SWAP}} M \otimes M \right)$$

(L3) Let $S_1 = S_{11} S_{12} \dots S_{1s}$ be s 1-qubit systems
 $S_2 = S_{21} S_{22} \dots S_{2s} \dots$

$$\text{Then } \text{SWAP}_{S_1 S_2} = \frac{1}{2^s} \sum_{P \in P_s} P \otimes P$$

$P \in P_s = 4^s$ Pauli matrices on s qubits

$$P_f: \text{SWAP}_{S_{1i} S_{2i}} = \frac{1}{2} (1 + XX + YY + ZZ) = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}_{S_{1i} S_{2i}}$$

$$\text{SWAP}_{S_1 S_2} = \bigotimes_{i=1}^s \text{SWAP}_{S_{1i} S_{2i}} = \frac{1}{2^s} \sum_{P \in P_s} P \otimes P$$

state matrix extension of P_1 i.e. $\tau_{r_2} P_{12} = P_1$

(L4) $\text{Tr}(P_1 Q) = \text{Tr}(P_{12} Q \otimes I_2)$

(L5) $\mathbb{E}_{V \in C_{st}} (V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger) \left(\underbrace{I_{T_1 T_2}}_{\text{each with } t \text{ qubits}} \otimes \text{SWAP}_{S_1 S_2} \right) (V_{T_1 S_1} \otimes V_{T_2 S_2})$
 $V \in C_{st} = \text{Clifford group on } st \text{ qubits}$

$$= \alpha I_{(T_1 S_1)(T_2 S_2)} + \beta \text{SWAP}_{(T_1 S_1)(T_2 S_2)}$$

$$\text{for } \alpha = 2^s \left[\frac{4^t - 1}{4^{st} - 1} \right] \leq \frac{1}{2^s} = \frac{1}{1s.1}$$

$$\beta = 2^t \left[\frac{4^s - 1}{4^{st} - 1} \right] \leq \frac{1}{2^t} = \frac{1}{1t.1}$$

Pf: By (L3), LHS of (L5)

$$= \mathbb{E}_{V \in C_{st}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\underbrace{I_{T_1 T_2}}_{\sum_{P \in P_s} P_{S_1} \otimes P_{S_2}} \otimes \frac{1}{2^s} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$$

• If $P = I$, $\mathbb{E}_{V \in C_{st}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\underbrace{I_{T_1 T_2}}_{I_{S_1 S_2}} \otimes I_{S_1 S_2} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$
 $= I_{T_1 S_1 T_2 S_2}$

• If $P \neq I$, $\mathbb{E}_{V \in C_{st}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\left(\underbrace{I_{T_1}}_{P_{S_1}} \otimes \underbrace{I_{T_2}}_{P_{S_2}} \right) \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$

see Appendix 3
 $= \mathbb{E}_{V \in C_{st}} V_{T_1 S_1}^\dagger (I_{T_1} P_{S_1}) V_{T_1 S_1} \otimes V_{T_2 S_2}^\dagger (I_{T_2} P_{S_2}) V_{T_2 S_2}$

The Clifford group permutes Pauli matrices by conjugation
 Δ acts transitively on all the non-identity Pauli's.
 (for any $P_1 \neq I, P_2 \neq I$, $\exists V$ s.t. $VP_1 V^\dagger = P_2$)

$$= \frac{1}{\underbrace{4^{st} - 1}_{\# \text{ of possible } Q = V^\dagger \otimes P V \neq I, Q \in P_{st}}} \sum_{Q \in P_{st}, Q \neq I} Q \otimes Q$$

(L3) $= \left(\frac{1}{4^{st} - 1} \right) \left(2^{st} \text{SWAP}_{(T_1 S_1)(T_2 S_2)} - I \otimes I \right)$

So LHS of (L5)

$$= \mathbb{E}_{V \in C_{st}} V_{T_1 S_1}^\dagger \otimes V_{T_2 S_2}^\dagger \left[\underbrace{I_{T_1 T_2}}_{\sum_{P \in P_s} P_{S_1} \otimes P_{S_2}} \otimes \frac{1}{2^s} \right] V_{T_1 S_1} \otimes V_{T_2 S_2}$$

$$= \frac{1}{2^s} (P=I \text{ case}) + \frac{4^s - 1}{2^s} (P \neq I \text{ case})$$

$$= \frac{1}{2^s} I_{T_1 S_1 T_2 S_2} + \frac{4^s - 1}{2^s} \frac{1}{4^{st} - 1} \left(2^{st} \text{SWAP}_{(T_1 S_1)(T_2 S_2)} - I_{T_1 S_1} I_{T_2 S_2} \right)$$

$$= \frac{1}{2^s} \left(1 - \frac{4^s - 1}{4^{st} - 1} \right) I_{T_1 S_1 T_2 S_2} + 2^t \frac{4^s - 1}{4^{st} - 1} \text{SWAP}_{(T_1 S_1)(T_2 S_2)}$$

$$\frac{1}{2^s} \left(\frac{4^{st} - 4^s}{4^{st} - 1} \right) = 2^s \left(\frac{4^t - 1}{4^{st} - 1} \right) = \alpha \quad \parallel \quad \beta$$

NB The average over Clifford group in (L5) same as
 average over V drawn over the Haar meas.

The above allows stabilizer codes (not random codes)
 to be used.

A slightly modified proof (appendix 3)

shows that the Clifford group is a "2-design".

Back to $\| \rho_{\bar{R}\bar{E}} - (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \|_2^2$

$$\begin{aligned}
 &= \text{tr} \left[\left(\rho_{\bar{R}\bar{E}} - (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \right) \left(\rho_{\bar{R}\bar{E}} - (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \right) \right] \\
 &= \text{tr}(\rho_{\bar{R}\bar{E}}^2) - 2 \text{tr} \left(\rho_{\bar{R}\bar{E}} (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \right) + \text{tr} \left((\frac{I}{2^{nr}})_{\bar{R}}^2 \otimes \rho_{\bar{E}}^2 \right) \\
 &= \text{tr}(\rho_{\bar{R}\bar{E}}^2) - 2 \underbrace{\left(\frac{1}{2^{nr}} \text{tr}(\rho_{\bar{R}\bar{E}} (I_{\bar{R}} \otimes \rho_{\bar{E}})) \right)}_{\textcircled{14}} + \frac{1}{2^{nr}} \text{tr} \rho_{\bar{E}}^2 \\
 &\quad \text{tr}(\text{tr}_{\bar{R}}(\rho_{\bar{R}\bar{E}}) \cdot \rho_{\bar{E}}) = \text{tr}(\rho_{\bar{E}} \cdot \rho_{\bar{E}}) \\
 &= \text{tr}(\rho_{\bar{R}\bar{E}}^2) - \frac{1}{2^{nr}} \text{tr} \rho_{\bar{E}}^2
 \end{aligned}$$

Now bounding $\mathbb{E} \text{tr}(\rho_{\bar{R}\bar{E}})^2$

$$\begin{aligned}
 &\textcircled{12} \quad \mathbb{E} \text{tr} \left[\left(\rho_{\bar{R}_1 \bar{E}_1} \otimes \rho_{\bar{R}_2 \bar{E}_2} \right) \text{SWAP}_{(\bar{R}_1 \bar{E}_1)(\bar{R}_2 \bar{E}_2)} \right] \\
 &\textcircled{14} \quad \mathbb{E} \text{tr} \left[\underbrace{\left(\rho_{B_{31} \bar{R}_1 \bar{E}_1} \otimes \rho_{B_{32} \bar{R}_2 \bar{E}_2} \right)}_{V_{B_{31} \bar{R}_1} \otimes I_{\bar{E}_1} (\alpha_{\bar{R}_1 \bar{E}_1})^\dagger \otimes I_{\bar{E}_2} V_{B_{31} \bar{R}_1}^\dagger} \left(I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right) \right] \\
 &\quad \text{Similarly } \boxed{\bar{R}_1 = \bar{R}_2, B_{31}} \\
 &\text{qdict} \quad = \mathbb{E} \text{tr} \left[\left(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \right) \times \left(V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger \otimes I_{\bar{E}_1 \bar{E}_2} (I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2}) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes I_{\bar{E}_1 \bar{E}_2} \right) \right]
 \end{aligned}$$

$$\begin{aligned}
 &= \text{tr} \left[\left(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \right) \times \left(\mathbb{E} V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger \otimes I_{\bar{E}_1 \bar{E}_2} (I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2}) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes I_{\bar{E}_1 \bar{E}_2} \right) \right] \\
 &= \text{tr} \left[\left(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \right) \times \left(\mathbb{E} V_{B_{31} \bar{R}_1}^\dagger \otimes V_{B_{32} \bar{R}_2}^\dagger \otimes (I_{B_{31} B_{32}} \otimes \text{SWAP}_{\bar{R}_1 \bar{R}_2}) V_{B_{31} \bar{R}_1} \otimes V_{B_{32} \bar{R}_2} \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right) \right] \\
 &\textcircled{15} \quad \leq \text{tr} \left[\left(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \right) \times \left[\left(\frac{1}{|\bar{R}|} I_{B_{31} \bar{R}_1, B_{32} \bar{R}_2} + \frac{1}{|B_3|} \text{SWAP}_{(B_{31} \bar{R}_1)(B_{32} \bar{R}_2)} \right) \otimes \text{SWAP}_{\bar{E}_1 \bar{E}_2} \right] \right] \\
 &\quad \downarrow \textcircled{14} \\
 &= \frac{1}{|\bar{R}|} \text{tr}(\alpha_{\bar{E}_1} \otimes \alpha_{\bar{E}_2} \text{SWAP}_{\bar{E}_1 \bar{E}_2}) + \frac{1}{|B_3|} \text{tr}(\alpha_{\bar{R}_1 \bar{E}_1} \otimes \alpha_{\bar{R}_2 \bar{E}_2} \text{SWAP}_{\bar{R}_1 \bar{E}_1, \bar{R}_2 \bar{E}_2})
 \end{aligned}$$

$$= \frac{1}{|\bar{R}|} \text{tr}(\alpha_{\bar{E}}^2) + \frac{1}{|B_3|} \text{tr}(\alpha_{\bar{R}\bar{E}}^2)$$

$$\begin{aligned}
 \left\| \rho_{\bar{R}\bar{E}} - (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \right\|_2^2 &\leq \underbrace{\left(\frac{1}{|\bar{R}|} \text{tr}(\alpha_{\bar{E}}^2) + \frac{1}{|B_3|} \text{tr}(\alpha_{\bar{R}\bar{E}}^2) \right)}_{\text{q}} - \frac{1}{2^{nr}} \text{tr} \rho_{\bar{E}}^2 \\
 &= \frac{1}{|B_3|} \text{tr}(\alpha_{\bar{R}\bar{E}}^2)
 \end{aligned}$$

$$\text{By } \textcircled{1}, \left\| \rho_{\bar{R}\bar{E}} - (\frac{I}{2^{nr}})_{\bar{R}} \otimes \rho_{\bar{E}} \right\|_{\text{tr}}^2 \leq \frac{|\bar{R}\bar{E}|}{|B_3|} \text{tr}(\alpha_{\bar{R}\bar{E}}^2)$$

$\alpha_{\bar{R}\bar{E}} = (\alpha_{\bar{R}\bar{E}})^{\otimes n}$ projected onto typical space of. Each eigenvalue $\approx 2^{nS(\bar{R}\bar{E})}$ $\therefore \text{tr}(\alpha_{\bar{R}\bar{E}}^2) \approx 2^{nS(\bar{R}\bar{E})}$

If we choose $|\bar{R}| = 2^{\frac{n}{2}[S(R:B) - \delta]}$
 $|B_3| = 2^{\frac{n}{2}[S(R:E) + \delta]}$

then $\| \cdot \|_{\text{tr}} \leq \frac{2^{\frac{n}{2}[S(R:B) - \delta]}}{2^{\frac{n}{2}[S(R:E) + \delta]}} \times \frac{2^{n(S(\bar{E}) + \epsilon)}}{2^{n(S(\bar{R}\bar{E}) - \epsilon)}}$

$$\begin{aligned}
 &= 2^{\frac{n}{2}[(S(R+B) - S(R\bar{E})) - (S(R) + S(\bar{E}) - S(\bar{R}\bar{E})) - 2\delta]} \cdot 2^{n(S(\bar{E}) - S(\bar{R}\bar{E}) + 2\epsilon)} \\
 &= 2^{n(S(\bar{R}) - S(\bar{R}\bar{E}) - \delta)} \cdot 2^{n(S(\bar{E}) - S(\bar{R}\bar{E}) + 2\epsilon)} \\
 &= 2^{-n(\delta - 2\epsilon)} \quad \text{choose } \delta = 3\epsilon \quad \text{fixed by how good the typical spaces are} \\
 &= 2^{-n\epsilon} \rightarrow 0 \text{ as } n \rightarrow \infty
 \end{aligned}$$

NB the unassisted case (the LSD thm) has $|B_3| = 1$

We choose $|\bar{R}| = 2^{\frac{n}{2}[S(R:B) - S(R:E) - 2\delta]}$

$$= 2^{n[I_c(R>B) - \delta]}$$

and the same decoupling condition for $\bar{R}\bar{E}$ holds.

The mapping from the direct coding scheme to the decoupling condition based on $\alpha^{\otimes n}$ has a glitch: $\sum_{(b)} \text{tr}(\alpha^{\otimes n}) = \text{tr}(\alpha^{\otimes n})$

for other means out comes, decoupling still works but the proof is involved (see 0702005 p9-10 proof Thm IV), & omitted

Alt: just use the fatter to get LSD.

① matter:

$$n \{ \{ \rightarrow \} \} + \frac{n}{2} (S(R:E) + d) \{ \{ \rightarrow \} \} \geq \frac{n}{2} (S(R:B) - d) \{ \{ \rightarrow \} \}$$

noisy static 2-party quantum resource (2RB here) evaluated on d qbit ebits

noiseless dynamical

the more Erhas the more assistance it takes to decompose

Bob's full potential

② father:

$$n \{ \{ \rightarrow \} \} + \frac{n}{2} (S(R:E) + d) \{ \{ \rightarrow \} \} \geq \frac{n}{2} (S(R:B) - d) \{ \{ \rightarrow \} \}$$

noisy dynamical resource (NRC here) evaluation IOWN | 0 > = | 2 > ebits qbit

Note that we have asymptotic approximate resource inequalities here. Say, XXX >= YYY. We demand the output resource YYY (lesser side) to be close to the ideal resource in trace distance or diamond norm. This ensures the protocol underlying the resource inequality can be used as a subroutine in any other protocol to produce YYY (using XXX) and when YYY is consumed, it is basically as good as ideal.

Say, XXX + ZZZ >= YYY + ZZZ >= KKK

The first inequality only holds if XXX >= YYY is given by a protocol producing sufficiently

Appendix 1:

Let $\sum_{i=1}^d |i\rangle\langle i| = \sqrt{d} \times \text{max ent state}$

The well-known transpose trick says the following:
 $\forall d \times d$ matrices M

$$\begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \text{---} \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} = \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \text{---} \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \quad (\text{this holds also for MES})$$

This can be generalized:

$\forall d \times d \times d$ M

$$\begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \text{---} \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} = \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \text{---} \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \text{---} \begin{array}{c} d \\ \hline \text{---} \\ d \end{array} \quad (\text{remember to re-insert the normalization when applying this to max ent states})$$

Pf The output of the LHS

$$= I \otimes V \left(\sum_{i,j} |i\rangle\langle j| \right) |0\rangle \quad \left[\text{let } V|i,j\rangle = \sum_{rs} V_{ijrs} |rs\rangle \right]$$

$$= \sum_{i,j} |i\rangle\langle j| \sum_{rs} V_{ijrs} |rs\rangle$$

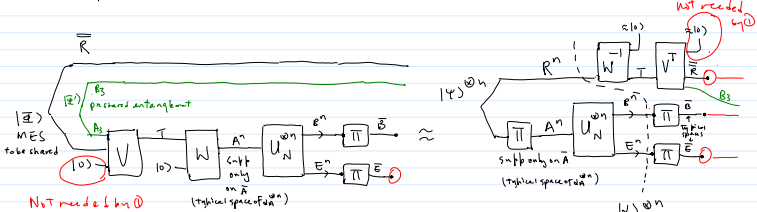
$$= \sum_{rs} \left(\sum_{i,j} V_{ijrs} |i\rangle\langle j| \right) |rs\rangle$$

$$\left[\sum_{i,j} V_{ijrs} |i\rangle\langle j| \right]$$

$$= \sum_{rs} (I \otimes |0\rangle\langle 0|) (V^T |rs\rangle) \otimes |rs\rangle$$

$$= (I \otimes |0\rangle\langle 0| \otimes I) (V^T \otimes I) \sum_{rs} |rs\rangle\langle rs| = \text{output of RHS}$$

for our purpose, ie to show



we don't need the above generalization of the transpose trick.

① In the end, $\dim(\bar{R}) = 2^{\frac{n}{2} [S(R:B) - d]}$

$$\dim(B_3) = 2^{\frac{n}{2} [S(R:E) + d]}$$

$$\text{So } \dim(\bar{R} B_3) = \dim(T) = 2^{n(S(R:E) + d)} \quad \text{and } \begin{array}{c} \text{---} \\ 10 \end{array} \text{---} = \begin{array}{c} \text{---} \\ 10 \end{array} \text{---}$$

(if necessary add 2 terms to dim(B₃))

(For the unassisted case the LSD with (B₃) = 0)

Moving V to the upper register requires an actual proj on the upper register. We can have a meas affecting one out of many possible projections.

Then IV of 0702005 is needed to ensure any outcome gives the same decompiling condition.)

② To move W: note initial state on RHS \approx MES on $2^{n S(R:B)}$ dims not 2^{ndR} dims. So moving W to W^T only requires the original transpose trick.

Appendix 2:

Alt proof for L2: $\text{tr}(M^2) = \text{tr}(\text{SWAP } M \otimes M)$

Pf: $\text{tr}(\text{SWAP } M \otimes M)$

$$= \text{tr}(\text{SWAP } M \otimes I \cdot I \otimes M)$$

$$= \text{tr}(\underbrace{\text{SWAP } M \otimes I}_{I \otimes M} \text{SWAP } I \otimes M)$$

$$= \text{tr}(I \otimes M \text{SWAP } I \otimes M)$$

$$= \text{tr}(\text{SWAP } I \otimes M^2) \stackrel{④}{=} \text{tr}(\underbrace{(\text{tr SWAP})}_{\mathbb{I}} M^2)$$

$$= \text{tr}(M^2)$$

$$\text{NB } \text{tr}_c(\text{SWAP}) = I$$

— Kronecker δ -fun.

$$\text{Pf: } \text{SWAP} = \sum_{i,j} |j\rangle\langle i| \otimes |i\rangle\langle j|, \text{tr}(|j\rangle\langle i|) = \delta_{ij}$$

$$\therefore \text{tr}_c(\text{SWAP}) = \sum_{i,j} \delta_{ij} (|i\rangle\langle j|) = I \text{ (on } 2^{\text{nd}} \text{ sys)}$$

Appendix 3:

$$\text{Let } \tau(M) = \mathbb{E}_{V \in C_{\text{st}}} V \otimes V M V^\dagger \otimes V^\dagger$$

$V \in C_{\text{st}}$ (support group on st qubits)

① If $P \neq I$, $P \notin P_{\text{st}}$ (Pauli group on st qubits)

$$\text{then } \tau(P \otimes P) = \frac{1}{4^{\text{st}} - 1} \sum_{\substack{Q \in P_{\text{st}} \\ Q \neq I}} Q \otimes Q$$

Pf: Note that C_{st} is transitive on $P_{\text{st}} - \{I\}$

ie for any $Q_1, Q_2 \in P_{\text{st}} - \{I\}$

$$\exists W \in C_{\text{st}} \text{ s.t. } W Q_1 W^\dagger = Q_2$$

Also, $\forall V \in C_{\text{st}}, V(P_{\text{st}} - \{I\}) V^\dagger$ only permutes elements of $P_{\text{st}} - \{I\}$.

$$\text{So } \tau(P \otimes P) = \sum_{Q \in P_{\text{st}} - \{I\}} \mu(Q) Q \otimes Q \text{ for some distribution } \mu(Q)$$

If $\mu(Q)$ not uniform, then $\exists Q_1, Q_2$ s.t. $\mu(Q_1) \neq \mu(Q_2)$

$$\text{Let } W' Q_1 W'^\dagger = Q_2$$

$$\text{then } \tau(P \otimes P) = W' \tau(P \otimes P) W'^\dagger \left(\begin{array}{l} \text{this merely changes} \\ \mathbb{E}_{V \in C_{\text{st}}} \text{ to } \mathbb{E}_{W'V \in C_{\text{st}}} \end{array} \right)$$

$Q_1 \otimes Q_2$ has weight $\mu(Q_1)$ on LHS
 $Q_2 \otimes Q_2$ — — — $\mu(Q_1)$ on RHS } \otimes
 (∵ W' is a permutation on $P_{\text{st}} - \{I\}$, the $Q_2 \otimes Q_2$ term in the

2nd line can only come from the $Q_1 \otimes Q_1$ term before conjugation by W' .)

But $\{Q \otimes Q\}$ is trace orthonormal.

∴ \otimes is a contradiction ∵ $\mu(Q)$ has to be uniform

$$\text{② } \tau(P \otimes Q) = 0 \quad \forall P \neq Q, P, Q \in P_{\text{st}}$$

Pf. WLOG $P \neq I$ (at least one of $P, Q \neq I$)

Then are $4^{\text{st}} - 2$ Pauli's anti commuting with P
 & commuting with Q

Let R be one of them

$$\tau(P \otimes Q) = \mathbb{E}_{V \in C_{\text{st}}} V \otimes V P \otimes Q V^\dagger \otimes V^\dagger$$

$$= \frac{1}{2} \left[\mathbb{E}_{V \in C_{\text{st}}} V \otimes V P \otimes Q V^\dagger \otimes V^\dagger + \mathbb{E}_{V \in C_{\text{st}}} V R \otimes V R P \otimes Q (V R)^\dagger \otimes (V R)^\dagger \right]$$

$$= 0.$$

So $\forall M, \tau(M)$ = linear combination of II & SWAP .

It easy to show that the coeffs are same as that of

$$\int_{\text{Haar}} U \otimes U M U^\dagger \otimes U^\dagger \quad (\text{average } U \text{ over Haar meas.})$$

(see Q. state hiding paper DiVincenzo, L. Terhal for detail.)