Name: ID:

## PMath 330 Assignment 5

Mark

## A Resolution Derivation

Given the collection of 8 clauses

1. $\{Q, S\}$
2. $\{R, S\}$
3. $\{\neg P, Q\}$
4. $\{P, R\}$
5. $\{P, \neg Q\}$
6. $\{\neg P, \neg R\}$
7. $\{\neg Q, \neg S\}$
8. $\{\neg R, \neg S\}$
fill in the reasons for the following resolution derivation:


Is it possible to find an assignment of truth values for the propositional variables $P, Q, R, S$ that will satisfy the original eight clauses? (if yes, give one)

Given the collection $\mathcal{S}$ of 6 clauses

1. $\{P, \neg Q\}$
2. $\{Q, S\}$
3. $\{P, R\}$
4. $\{\neg P, \neg R\}$
5. $\{\neg Q, \neg S\}$
6. $\{\neg R, \neg S\}$
fill in the reasons for the following resolution steps:

| 7. $\{Q, \neg R\}$ | 14. $\{P, S\}$ |
| :---: | :---: |
| 8. $\{P, \neg S\}$ | 15. $\{P, \neg R\}$ |
| 9. $\{S, \neg S\}$ | 16. $\{P\}$ |
| 10. $\{Q, \neg Q\}$ | 17. $\{\neg R, S\}$ |
| 11. $\{\neg Q, \neg R\}$ | 18. $\{P, Q\}$ |
| 12. $\{P, \neg P\}$ | 19. $\{\neg R\}$ |
| 13. $\{R, \neg R\}$ |  |

Can you obtain any other clauses by resolution?
What does this say about the satisfiability of $\mathcal{S}$ ?
$\qquad$

Apply the Davis-Putnam Procedure to the First Problem, showing just the $\mathcal{S}_{i}{ }^{\prime}$ and $\mathcal{U}_{i}$ steps (as done for the resolution on $R$ below).

## Resolution on $R$ :

(1)
(2)
(3)
(4)
$\mathcal{S}_{1}{ }^{\prime}:\{Q, S\} \quad\{R, S\} \quad\{\neg P, Q\} \quad\{P, R\} \quad\{P, \neg Q\} \quad\{\neg P, \neg R\} \quad\{\neg Q, \neg S\} \quad\{\neg R, \neg S\}$
$(1,3) \quad(1,4) \quad(2,3) \quad(2,4)$
$\mathcal{U}_{1}: \quad\{\neg P, S\} \quad\{S, \neg S\} \quad\{P, \neg P\} \quad\{P, \neg S\}$
Resolution on $P$ :
$\mathcal{S}_{2}{ }^{\prime}:$
$\mathcal{U}_{2}$ :

## Resolution on $Q$ :

$\mathcal{S}_{3}{ }^{\prime}:$
$\mathcal{U}_{3}:$

## Resolution on $S$ :

$\mathcal{S}_{4}{ }^{\prime}:$
$\mathcal{U}_{4}:$

Given the collection of five Horn clauses

1. $\{\neg P, Q\}$
2. $\{P\}$
3. $\{\neg R\}$
4. $\{S, \neg Q\}$
5. $\{\neg S, R\}$
find all clauses that can be derived using unit resolution:


Given the clauses $\{\neg P, Q, R\}$ and $\{\neg Q, R, S\}$ prove, using just the definitions, that if $\vec{e}$ is a truth evaluation of $P, Q, R, S$ that makes the two clauses true then it also makes the clause $\{\neg P, R, S\}$, obtained by resolving the two over $Q$, true.

Consider the propositional argument:

$$
\begin{array}{ll}
\mathrm{F}_{1}: & (P \rightarrow \neg Q) \rightarrow R \\
\mathrm{~F}_{2}: & P \vee \neg(Q \rightarrow R) \\
\mathrm{F}: & P \leftrightarrow(Q \leftrightarrow R)
\end{array}
$$

Give the conjunctive normal form for each of the following formulas:
$F_{1}$ : $\qquad$
$\mathrm{F}_{2}$ : $\qquad$
$\neg F$ : $\qquad$
From this derive a set $\mathcal{S}$ of clauses such that $\mathrm{F}_{1}, \mathrm{~F}_{2} \therefore \mathrm{~F}$ is valid iff $\neg \operatorname{Sat}(\mathcal{S}) . \mathcal{S}$ has the clauses:

1. $\qquad$ 4. $\qquad$
2. $\qquad$ 5. $\qquad$
3. $\qquad$ 6. $\qquad$
Is $\mathcal{S}$ satisfiable? (Reasons)

Is the original argument valid? (Reasons)

