

FINITENESS PROPERTIES OF LOCALLY FINITE ABELIAN VARIETIES

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ABSTRACT. We show that any locally finite abelian variety is generated by a finite algebra. We solve a problem posed by D. Hobby and R. McKenzie by exhibiting a nonfinitely based finite abelian algebra.

1. INTRODUCTION

A *variety* of algebras is an equationally definable class of algebras in a fixed language. A variety \mathcal{V} is said to be *locally finite* if its finitely generated algebras are finite. An algebra \mathbf{A} is *abelian* if it satisfies all sentences of the form

$$\forall \mathbf{x}, \mathbf{y}, \mathbf{z}, \mathbf{u} (t(\mathbf{x}, \mathbf{u}) = t(\mathbf{x}, \mathbf{v}) \implies t(\mathbf{y}, \mathbf{u}) = t(\mathbf{y}, \mathbf{v}))$$

where $t(\mathbf{x}, \mathbf{y})$ is an $(m + n)$ -ary term in the language of \mathbf{A} . A variety of algebras is said to be abelian if all of its members are abelian.

For any ring \mathbf{R} the variety of left \mathbf{R} -modules is abelian. For any monoid \mathbf{M} the variety of left \mathbf{M} -sets is abelian. A variety of \mathbf{R} -modules is locally finite if and only if \mathbf{R} is finite while a variety of \mathbf{M} -sets is locally finite if and only if \mathbf{M} is finite. These two types of examples are fundamental, since it is shown in [1] that the polynomial structure of any finite abelian algebra is locally like that of an \mathbf{R} -module or an \mathbf{M} -set. Therefore it is natural to wonder if one can associate to each locally finite abelian variety \mathcal{V} a finite structure, of which a ring or a monoid is a special case, which acts on the members of \mathcal{V} and which ‘determines’ \mathcal{V} in some sense. If such were the case, then we would expect a locally finite abelian variety to share the finiteness properties of locally finite varieties of \mathbf{R} -modules and \mathbf{M} -sets. For example, since a subvariety of a variety of \mathbf{R} -modules (or \mathbf{M} -sets) may be identified with a variety of \mathbf{R}' -modules (\mathbf{M}' -sets) where \mathbf{R}' (\mathbf{M}') is a quotient of \mathbf{R} (\mathbf{M}), it follows that a variety of either type has finitely many subvarieties. This implies that either type of variety is finitely generated and, with a little extra argument, that either type of variety is finitely based. If to each locally finite abelian variety we could associate a finite structure which determined the variety, we would expect every locally finite abelian variety to be finitely generated, finitely based and to have a finite subvariety lattice.

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In the first part of this paper we prove that any locally finite abelian variety is finitely generated. In the second part we then give an example to show that a locally finite abelian variety need not be finitely based (and therefore need not have a finite subvariety lattice). By the first part of the paper, the variety of the second part is generated by a finite abelian algebra. Therefore the results contained in this paper answer in the negative Problem 3 from the book [1] of D. Hobby and R. McKenzie which asks:

If \mathbf{A} is a finite abelian algebra of finite type, is $\mathcal{V}(\mathbf{A})$ finitely axiomatizable?

The nonfinitely based abelian algebra which we construct generates an abelian variety. By a result of [2] any such algebra fails to be inherently nonfinitely based. It remains open whether there is a finite abelian algebra which is inherently nonfinitely based.

2. FINITE GENERATION

In this section we will prove that every locally finite abelian variety is generated by a finite algebra. We begin by making some general remarks about what it means for a variety to be generated by a finite algebra.

Let \mathcal{V} be a variety. Any equation in the language of \mathcal{V} may be written (after possibly renaming variables and adding ‘fictitious variables’) as $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ where s and t are $(n + 1)$ -ary term operations, $\mathbf{x} = (x_1, \dots, x_n)$ is a sequence of distinct variables and y is a distinguished variable not in the sequence \mathbf{x} . We shall consider only equations of this form. The *rank* of the equation $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ is defined to be the length of the sequence \mathbf{x} . (Observe that by adding a fictitious variable to this equation we change its rank, so we consider the modified equation to be different from the original equation.) An equation $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ is *falsifiable* in \mathcal{V} if there is some $\mathbf{A} \in \mathcal{V}$ and a tuple $(\mathbf{a}, b) \in A^{n+1}$ such that $s(\mathbf{a}, b) \neq t(\mathbf{a}, b)$. That is, an equation is falsifiable if it is not an equation of \mathcal{V} . A *minimal falsifiable equation* is a falsifiable equation $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ which is no longer falsifiable if two variables are set equal.

Lemma 2.1. *A variety is finitely generated if and only if it is locally finite and there is a finite bound on the rank of any minimal falsifiable equation.*

Proof. Assume that $\mathcal{V} = \mathcal{V}(\mathbf{A})$ where \mathbf{A} is finite. It is well known that the variety generated by a finite algebra is locally finite. Choose a minimal falsifiable equation $s = t$. Since \mathbf{A} generates \mathcal{V} and $s = t$ is a minimal falsifiable equation, we can falsify $s = t$ by a substitution of distinct elements of A into the distinct variables of the equation. This implies that the number of variables in the equation $s = t$ does not exceed $|A|$, and therefore the rank of $s = t$ is $< |A|$. Hence $|A|$ is a finite bound on the rank of any minimal falsifiable equation.

Conversely, assume that \mathcal{V} is a locally finite variety and that N is a finite bound on the rank of any minimal falsifiable equation. Then $\mathbf{F}_{\mathcal{V}}(N)$ is a finite algebra which satisfies all the equations of \mathcal{V} and fails all the equations that can be falsified in \mathcal{V} . Thus, $\mathcal{V} = \mathcal{V}(\mathbf{F}_{\mathcal{V}}(N))$. \square

From now on \mathcal{A} will denote a fixed but arbitrarily chosen locally finite abelian variety. It is our goal to prove that there is a finite bound on the rank of any minimal falsifiable equation in the language of \mathcal{A} . First we explain a reduction. Let δ be a unary term in the language of \mathcal{A} and let \mathcal{E}_{δ} denote the set of equations $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ in the language of \mathcal{A} for which

$$\mathcal{A} \models s(y, y, \dots, y) = \delta(y) = t(y, y, \dots, y).$$

Clearly, if $\mathcal{A} \models \delta(y) = \delta'(y)$, then $\mathcal{E}_{\delta} = \mathcal{E}_{\delta'}$. Therefore, since \mathcal{A} has only finitely many unary terms up to equivalence, there are only finitely many different sets of the form \mathcal{E}_{δ} . Furthermore, if $s(\mathbf{x}, y) = t(\mathbf{x}, y)$ is any minimal falsifiable equation of rank greater than zero, then $\mathcal{A} \models s(y, \dots, y) = t(y, \dots, y)$; thus $s = t$ belongs to \mathcal{E}_{δ} for $\delta(y) := s(y, y, \dots, y)$. It follows that there is a finite bound on the rank of all minimal falsifiable equations if and only if there is a finite bound on the rank of the minimal falsifiable equations in \mathcal{E}_{δ} for each δ . Henceforth it will be our goal to show that there is a finite bound on the rank of any minimal falsifiable equation in a fixed but arbitrarily chosen \mathcal{E}_{δ} .

To accomplish our goal we need to understand the structure the minimal falsifiable equations in \mathcal{E}_{δ} . Our analysis depends in an essential way on the theorem of E. W. Kiss and M. A. Valeriote which connects abelian varieties to Hamiltonian varieties.

Definition 2.2. A variety \mathcal{V} is said to be *Hamiltonian* provided that whenever $\mathbf{A} \in \mathcal{V}$ and S is a subuniverse of \mathbf{A} , then S is a congruence block of \mathbf{A} .

The theorem of Kiss and Valeriote which is crucial for us is the following.

Theorem 2.3. [3] *A locally finite abelian variety is Hamiltonian.*

A characterization of Hamiltonian varieties is given by L. Klukovits in [4]. Klukovits showed that a variety \mathcal{V} is Hamiltonian if and only if for each term $t(\mathbf{x})$ and each choice of a variable of this term, say the i -th variable, there is a ternary term $k(u, y, z)$ such that

$$\mathcal{V} \models k(t(x_1, \dots, x_{i-1}, y, x_{i+1}, \dots, x_n), y, z) = t(x_1, \dots, x_{i-1}, z, x_{i+1}, \dots, x_n).$$

We call the term k a *Klukovits term for t in its i -th variable*. For example, if \mathcal{V} is a variety of left \mathbf{R} -modules and $t(\mathbf{x}) = r_1x_1 + \dots + r_nx_n$, then a Klukovits term for t in its i -th variable is $k(u, y, z) = u - r_iy + r_iz$. If \mathcal{V} is a variety of \mathbf{M} -sets and $t(\mathbf{x}) = mx_i$, then a Klukovits term for t in its i -th variable is $k(u, y, z) = mz$; a Klukovits term for t in its j -th variable, $j \neq i$, is $k(u, y, z) = u$. The phrase “Klukovits term” will

mean any ternary term k which is a Klukovits term for t in its i -th variable for some t and i .

Let K be a complete set of \mathcal{A} -inequivalent Klukovits terms. Note that K is finite since Klukovits terms are ternary and \mathcal{A} is locally finite. Let ω denote the set of natural numbers with the usual ordering, and let $\omega^{K \times K}$ denote the set of functions from $K \times K$ to ω ordered pointwise. Define a binary relation \triangleright from \mathcal{E}_δ to $\omega^{K \times K}$ by the rule that

$$(s = t) \triangleright f$$

if and only if $f : K \times K \rightarrow \omega$ has the property that there exists a sequence

$$\Sigma = \langle (j_1, k_1), (j_2, k_2), \dots, (j_n, k_n) \rangle,$$

where n is the rank of the equation $s = t$, such that

- (1) j_i is a Klukovits term for s in its i -th variable and k_i is a Klukovits term for t in its i -th variable, and
- (2) for any $(j, k) \in K \times K$, $f(j, k)$ = the number of times (j, k) occurs in the sequence Σ .

Lemma 2.4. *The following statements hold for the equations in \mathcal{E}_δ .*

- (1) *If $(s = t) \triangleright f$, then the height of f in $\omega^{K \times K}$ equals the rank of the equation $s = t$.*
- (2) *If $(s = t) \triangleright f$, and $g < f$ in $\omega^{K \times K}$, then $s = t$ has a specialization $(s' = t') \in \mathcal{E}_\delta$ such that $(s' = t') \triangleright g$.*
- (3) *If $p = q$ and $s = t$ are both \triangleright -related to the element $f \in \omega^{K \times K}$, then $p = q$ is equivalent to $s = t$ modulo the equations of \mathcal{A} . (I.e., $\text{Eq}(\mathcal{A}) \models (p = q) \Leftrightarrow (s = t)$.)*

Proof. The height of a function f in $\omega^{K \times K}$ is the sum of the values of f . Now if $(s = t) \triangleright f$ where $s = t$ is an equation of rank n , then recall that $f(j, k)$ is defined to be the number of times (j, k) occurs in some sequence $\langle (j_1, k_1), (j_2, k_2), \dots, (j_n, k_n) \rangle$ where this sequence is a sequence of pairs of Klukovits terms for the first n -variables of the equation $s = t$. It follows that $\sum_{(j,k) \in K \times K} f(j, k) = n$, which shows that the height of f is n whenever the rank of $s = t$ is n and $(s = t) \triangleright f$. This proves (1).

For (2), assume that $s = t$ is the equation $s(x_1, \dots, x_n, y) = t(x_1, \dots, x_n, y)$ where (by part (1)) the height of f is n . From the definitions, there exists a sequence $\langle (j_1, k_1), \dots, (j_n, k_n) \rangle$ of pairs of Klukovits terms for the first n variables of $s = t$ such that, for any $(j, k) \in K \times K$, the number of times (j, k) occurs in this sequence is $f(j, k) \geq g(j, k)$. Therefore it is possible to select a subset $\{i_1, \dots, i_m\} \subseteq \{1, \dots, n\}$ such that $\langle (j_{i_1}, k_{i_1}), \dots, (j_{i_m}, k_{i_m}) \rangle$ has precisely $g(j, k)$ occurrences of (j, k) for each $(j, k) \in K \times K$. Then define $s'(x_1, \dots, x_m, y) = t'(x_1, \dots, x_m, y)$ to be the equation obtained from $s(x_1, \dots, x_n, y) = t(x_1, \dots, x_n, y)$ by substituting y for x_h whenever $h \notin \{i_1, \dots, i_m\}$ and substituting x_g for x_{i_g} $i_g \in \{i_1, \dots, i_m\}$. The equation $s' = t'$ is in \mathcal{E}_δ since it is obtained from $s = t$ by substituting new variables for old.

Moreover, a sequence of Klukovits terms for the first m variables of $s'(x_1, \dots, x_m, y) = t'(x_1, \dots, x_m, y)$ is $\langle (j_{i_1}, k_{i_1}), \dots, (j_{i_m}, k_{i_m}) \rangle$. It follows that $(s' = t') \triangleright g$.

Finally we prove (3). To do this, we first define an action of pairs of Klukovits terms on equations. If $u = v$ is an equation of rank ℓ , then the pair (j, k) of Klukovits terms acts on $u = v$ (on the right) to produce a new equation of rank $\ell + 1$ as follows:

$$(u = v) \circ (j, k) := j(u(x_1, \dots, x_\ell, y), y, x_{\ell+1}) = k(v(x_1, \dots, x_\ell, y), y, x_{\ell+1}).$$

To start the proof of (3) assume that $p = q$ and $s = t$ are both \triangleright -related to f . This implies that there are sequences of pairs $\langle (j_1, k_1), \dots, (j_n, k_n) \rangle$ and $\langle (J_1, K_1), \dots, (J_n, K_n) \rangle$ where

- (a) the first sequence is a sequence of pairs of Klukovits terms for $p = q$,
- (b) the second sequence is a sequence of pairs of Klukovits terms for $s = t$ and
- (c) the second sequence is a permutation of the first sequence.

We must use this information to prove that $p = q$ and $s = t$ are equivalent modulo the equations of \mathcal{A} .

Claim. Let π be a permutation of $\{1, \dots, n\}$. Modulo the equations of \mathcal{A} , the equation $p = q$ is equivalent to

$$[\dots[(\delta(y) = \delta(y)) \circ (j_{\pi(1)}, k_{\pi(1)})] \circ \dots] \circ (j_{\pi(n)}, k_{\pi(n)}).$$

The proof of this claim establishes part (3) of this lemma. To see this, note that we can apply the claim once with π chosen so that $(j_{\pi(i)}, k_{\pi(i)}) = (J_i, K_i)$ to get an expression equivalent to $p = q$ which, by a second application of the claim to the equation $s = t$ and the permutation $\pi = \text{id}$, is equivalent to $s = t$ modulo the equations of \mathcal{A} . Thus we get that $p = q$ is equivalent to $s = t$ modulo the equations of \mathcal{A} .

To prove the claim, first note that since $(p = q) \in \mathcal{E}_\delta$ we have that $\mathcal{A} \models p(y, y, \dots, y) = \delta(y) = q(y, y, \dots, y)$. Therefore we are trying to show that $p = q$ is equivalent to

$$j_{\pi(n)}(\dots j_{\pi(1)}(p(y, \dots, y), y, x_1) \dots, y, x_n) = k_{\pi(n)}(\dots k_{\pi(1)}(q(y, \dots, y), y, x_1) \dots, y, x_n).$$

Using the Klukovits equations, which are equations of \mathcal{A} , this equation can be greatly simplified. We simplify it in n steps, working our way through this nested composition from the innermost part outwards. At the first step we have $j_{\pi(1)}(p(y, y, \dots, y), y, x_1)$ on the lefthand side, and the Klukovits equations reduce this to $p(y, y, \dots, x_1, \dots, y)$ with x_1 in the $\pi(1)$ -rst position. At the innermost part on the righthand side we have $k_{\pi(1)}(q(y, y, \dots, y), y, x_1)$ which simplifies to $q(y, y, \dots, x_1, \dots, y)$ with x_1 in the $\pi(1)$ -rst position. Similarly, as we work our way through each step of the composition we simply replace the y in position $\pi(i)$ on both sides of the equation with the variable x_i during the i -th step. The result is that the previously displayed equation is equivalent modulo the equations of \mathcal{A} to

$$p(x_{\pi^{-1}(1)}, \dots, x_{\pi^{-1}(n)}, y) = q(x_{\pi^{-1}(1)}, \dots, x_{\pi^{-1}(n)}, y),$$

which differs from the equation $p = q$ only by a permutation of variables. Thus the claim is proved. \square

Theorem 2.5. *\mathcal{A} is finitely generated.*

Proof. Define an order filter in $\omega^{K \times K}$ as follows:

$$F = \{x \in \omega^{K \times K} \mid (\exists g) x \geq g \text{ where } (p = q) \triangleright g \text{ for some falsifiable equation } p = q\}.$$

We claim that if $(s = t) \in \mathcal{E}_\delta$ is a minimal falsifiable equation and $(s = t) \triangleright f$, then f is a minimal member of F . To see this, assume that $s = t$ is falsifiable and that f is not minimal in F ; we will prove that $s = t$ is not a minimal falsifiable equation. Since f is not minimal in F there is a falsifiable equation $p = q$ such that $(p = q) \triangleright g$ and $g < f$. From Lemma 2.4 we deduce that the equation $s = t$ has a specialization $(s' = t') \in \mathcal{E}_\delta$ such that $(s' = t') \triangleright g$ and $s' = t'$ is equivalent to $p = q$ modulo the equations of \mathcal{A} . The equivalence of $s' = t'$ with the falsifiable equation $p = q$ implies that $s' = t'$ is falsifiable, which proves that $s = t$ has a falsifiable specialization of smaller rank. This establishes our claim.

Since K is finite, the ordered set $\omega^{K \times K}$ has the property that its order filters are finitely generated (see [5]). So, there is a natural number N such that every minimal element of F has height $\leq N$. It follows from the previous paragraph and Lemma 2.4 (1) that N is a bound on the rank of any minimal falsifiable equation in \mathcal{E}_δ . As we observed earlier, the fact that this is true for an arbitrarily chosen δ implies the existence of a finite bound on the rank of all minimal falsifiable equations. Lemma 2.1 can now be invoked to deduce that \mathcal{A} is finitely generated. \square

3. A NONFINITELY BASED ABELIAN ALGEBRA

In this section we describe a finite algebra which generates a nonfinitely based abelian variety. This provides a strong negative answer to Problem 3 of [1], and complements the result in [2] that states that no finite algebra can generate an inherently nonfinitely based abelian variety.

The idea behind our example is extremely simple, so we give a rough description now before facing the details. We plan to construct a variety \mathcal{P} of algebras whose models are (essentially) pairs of isomorphic Boolean groups¹ \mathbf{B} and \mathbf{C} glued together at a subset containing the common identity element, 0. We try to show that the subvariety of algebras where $B \cap C$ is a subgroup is not finitely based relative to \mathcal{P} . That is, we try to show that it is impossible to express the idea that $B \cap C$ is closed under sums without looking at a large number of elements of $B \cap C$ simultaneously.

Our idea does not work in the form just explained, because it is *not* hard to express the fact that $B \cap C$ is closed under sums: one can check elements of $B \cap C$ two at a time to see if their sum is in $B \cap C$. Therefore, to make this idea work, we need

¹A *Boolean group* is a group of exponent 2.

a subset Q disjoint from $B \cup C$ and an operation $s : Q \rightarrow B \cap C$ whose duty is to ‘select’ a subset of $B \cap C$. What we actually show is that it is hard to express the fact that the subgroup generated by $s(Q)$ lies in $B \cap C$. Here it may be that all sums of few elements of $s(Q)$ lie in $B \cap C$, but some sum of many elements lies outside $B \cap C$.

We will get a properly decreasing sequence of varieties $\mathcal{P} = \mathcal{V}_1 \supset \mathcal{V}_2 \supset \mathcal{V}_3 \supset \dots$, where \mathcal{V}_n is the collection of algebras where all sums of $\leq n$ elements of $s(Q)$ are in $B \cap C$. The intersection $\mathcal{V}_\infty = \bigcap_{n < \omega} \mathcal{V}_n$ is the nonfinitely based variety of algebras in \mathcal{P} where $s(Q)$ generates a subgroup of $B \cap C$. Since \mathcal{V}_∞ is locally finite and abelian, the result of the last section proves that \mathcal{V}_∞ is generated by a (nonfinitely based) finite algebra. We produce a concrete 6-element generating algebra for \mathcal{V}_∞ at the end of this section.

Our nonfinitely based algebra is of type $\langle 0, 1, 1, 1, 2, 2 \rangle$ and the corresponding operation symbols are $\langle 0, e, f, s, +, \oplus \rangle$. Our algebra will be a member of the (abelian) variety \mathcal{P} whose defining equations assert that in each $\mathbf{A} \in \mathcal{P}$:

- (A) $\{0\}$ is a subuniverse.
- (B) $ee(x) = e(x), ff(x) = f(x), ss(x) = 0,$
 $ef(x) = e(x), fe(x) = f(x),$
 $es(x) = fs(x) = s(x),$
 $se(x) = sf(x) = 0.$
- (C) $x + y = e(x) + e(y) = e(x + y) = e(x \oplus y),$
 $x \oplus y = f(x) \oplus f(y) = f(x \oplus y) = f(x + y).$
- (D) $\langle e(A); +, 0 \rangle$ and $\langle f(A); \oplus, 0 \rangle$ are Boolean groups.

We will soon see that \mathcal{P} is a locally finite abelian variety which contains a non-finitely based algebra. First we describe how to construct models of these equations.

We refine our earlier discussion of the models of \mathcal{P} by discussing a class of three-sorted structures of the form

$$\langle \mathbf{B}, \mathbf{C}, Q; \iota; e_Q, s_Q \rangle.$$

Here $\mathbf{B} = \langle B; *, 0 \rangle$ and $\mathbf{C} = \langle C; \circ, 0 \rangle$ are Boolean groups which have a common identity element. Q is a set which is disjoint from $B \cup C$. The unary function ι is an isomorphism $\iota : \mathbf{C} \rightarrow \mathbf{B}$ for which $\iota(x) = x$ for all $x \in B \cap C$. Both $e_Q : Q \rightarrow B$ and $s_Q : Q \rightarrow (B \cap C)$ are functions. There is no restriction on them other than that they have the correct domain and range.

From such a three-sorted structure $\langle \mathbf{B}, \mathbf{C}, Q; \iota; e_Q, s_Q \rangle$ we can construct a member of \mathcal{P} . Our algebra will have universe $A = B \cup C \cup Q$. We interpret the operations $\langle 0, e, f, s, +, \oplus \rangle$ as follows. We interpret 0 as the element already named $0 \in A$. We

define e and f by

$$e(x) = \begin{cases} x & \text{if } x \in B, \\ \iota(x) & \text{if } x \in C, \\ e_Q(x) & \text{if } x \in Q, \end{cases} \quad f(x) = \begin{cases} x & \text{if } x \in C, \\ \iota^{-1}(x) & \text{if } x \in B, \\ \iota^{-1}e_Q(x) & \text{if } x \in Q. \end{cases}$$

We define s so that $s(B \cup C) = \{0\}$ while $s|_Q = s_Q$. Next we define $x + y$ to be $e(x) * e(y)$, where $*$ is the group operation of \mathbf{B} . Similarly, $x \oplus y = f(x) \circ f(y)$ where \circ is the group operation of \mathbf{C} .

Lemma 3.1. *The algebra \mathbf{A} constructed as in the previous paragraph belongs to \mathcal{P} . Conversely, any member of \mathcal{P} is isomorphic to such an algebra.*

Sketch of proof. The first statement requires only the straightforward verification that the equations defining \mathcal{P} hold in \mathbf{A} .

For the second statement, choose any $\mathbf{D} \in \mathcal{P}$. Let $B = e(D)$, $C = f(D)$ and $Q = D - (B \cup C)$. The equations of \mathcal{P} of types (A), (B) and (D) ensure that B is closed under $+$ and 0 , C is closed under \oplus and 0 , and that $\mathbf{B} := \langle B; +', 0 \rangle$ and $\mathbf{C} := \langle C; \oplus', 0 \rangle$ are Boolean groups. Here the prime on $+$ and \oplus indicates that we are using the restrictions of the corresponding operations of \mathbf{D} . If we let $\iota = e|_C$, then the equations of type (C) guarantee that $\iota : \mathbf{C} \rightarrow \mathbf{B}$ is an isomorphism of groups which is the identity on $B \cap C$. Let $e_Q = e|_Q$ and $s_Q = s|_Q$. Of course, $e_Q(Q) \subseteq e(D) = B$. The equations of type (B) involving s ensure that $s_Q(Q) \subseteq e(D) \cap f(D) = B \cap C$. Thus, \mathbf{D} yields a three-sorted structure $\langle \mathbf{B}, \mathbf{C}, Q; \iota; f_Q, s_Q \rangle$.

We can apply the procedure outlined before the proof of this lemma to the three-sorted structure derived from \mathbf{D} . By doing so we reconstruct an algebra in \mathcal{P} which has the same universe as \mathbf{D} . One can check that each operation of the constructed algebra coincides with the corresponding operation in \mathbf{D} . Thus, \mathbf{D} is reconstructible from $\langle \mathbf{B}, \mathbf{C}, Q; \iota; e_Q, s_Q \rangle$. \square

Now we analyze the term operations of \mathcal{P} . In the next lemma we let $E(x) = e(x) + s(x)$ and $F(x) = f(x) \oplus s(x)$.

Lemma 3.2. *Any term of \mathcal{P} is \mathcal{P} -equivalent either to a variable or to a term of the form*

$$u_1(x_1) + u_2(x_2) + \cdots + u_n(x_n)$$

where each $u_i \in \{0, e, s, E\}$ or to

$$v_1(x_1) \oplus v_2(x_2) \oplus \cdots \oplus v_n(x_n)$$

where each $v_i \in \{0, f, s, F\}$.

Sketch of proof. The proof is a straightforward induction argument using the equations for \mathcal{P} . \square

Corollary 3.3. *\mathcal{P} is a locally finite abelian variety.*

Proof. To see that \mathcal{P} is abelian, choose $\mathbf{A} \in \mathcal{P}$ and a term $t(x, \mathbf{y})$. Without loss of generality we may assume that $t(x, \mathbf{y}) = u_0(x) + u_1(y_1) + \cdots + u_n(y_n)$. To check that the term condition holds for t we must show that for all $a, b \in A$ and $\mathbf{c}, \mathbf{d} \in A^n$

$$u_0(a) + u_1(c_1) + \cdots + u_n(c_n) = u_0(a) + u_1(d_1) + \cdots + u_n(d_n)$$

implies

$$u_0(b) + u_1(c_1) + \cdots + u_n(c_n) = u_0(b) + u_1(d_1) + \cdots + u_n(d_n).$$

The second equality follows from the first by adding $u_0(a) + u_0(b)$ to both sides of the first equality. This shows that \mathcal{P} is abelian.

It follows immediately from the previous lemma that \mathcal{P} has only finitely many n -ary terms up to equivalence for any finite n . Thus \mathcal{P} is locally finite. \square

Next we consider equations of the form

$$(E_n): \quad s(x_1) + s(x_2) + \cdots + s(x_n) = s(x_1) \oplus s(x_2) \oplus \cdots \oplus s(x_n).$$

Notice that by substituting 0 in for x_n in E_n we obtain E_{n-1} . Thus $E_n \Rightarrow E_{n-1}$. The next lemma proves that the reverse implication does not hold.

Lemma 3.4. *For each $n > 1$ there is an algebra in \mathcal{P} which satisfies E_{n-1} , but which fails E_n .*

Proof. Let $\mathbf{B} = \mathbf{Z}_2^n$ where $\mathbf{Z}_2 = \langle \{0, 1\}; *, 0 \rangle$ is the two-element group. Let $b_i \in B$ denote the element which has a one in the i -th position and zeros elsewhere. Let $\mathbf{0} \in B$ denote the element which has zeros in every position and let $\mathbf{1} \in B$ be the element with ones in every position. Let \mathbf{C} be a Boolean group obtained from \mathbf{B} by replacing the element $\mathbf{1}$ with a new element $\mathbf{1}'$ and naming the resulting group operation \circ . Observe that $B \cap C = B - \{\mathbf{1}\}$.

Let $Q = \{1, 2, \dots, n\}$. Let $\iota : \mathbf{C} \rightarrow \mathbf{B}$ be the isomorphism which fixes $B \cap C$ and maps $\mathbf{1}'$ to $\mathbf{1}$. Define e_Q arbitrarily and for each $i \in Q$ let $s_Q(i) = b_i$. We now have a three-sorted structure $\langle \mathbf{B}, \mathbf{C}, Q; \iota, e_Q, s_Q \rangle$. Satisfaction of E_k in the associated algebra $\mathbf{A} \in \mathcal{P}$ is equivalent to the satisfaction of

$$s_Q(x_1) * \cdots * s_Q(x_k) = s_Q(x_1) \circ \cdots \circ s_Q(x_k)$$

in $\langle \mathbf{B}, \mathbf{C}, Q; \iota, e_Q, s_Q \rangle$. The only way for this equation to fail is for the left hand side to equal $\mathbf{1}$ and (therefore) for the right hand side to equal $\mathbf{1}'$. If $k < n$ there are too few summands for this to happen, but when $k = n$ we may take $x_i = i$ and we get a failure of this equation. Hence the algebra \mathbf{A} fails E_n but satisfies all E_k for $k < n$. \square

Let \mathcal{V}_n denote the subvariety of \mathcal{P} axiomatized by E_n and the equations of \mathcal{P} . Let $\mathcal{V}_\infty = \bigcap_{n < \omega} \mathcal{V}_n$. The previous lemma shows that \mathcal{V}_∞ is not finitely based. Since \mathcal{V}_∞ is a locally finite abelian variety it is generated by a finite algebra. We shall produce a finite generating algebra shortly, but first we describe the subvariety lattice of \mathcal{P} .

Theorem 3.5. *The subvarieties of \mathcal{P} are: $\mathcal{P} = \mathcal{V}_1, \mathcal{V}_2, \dots, \mathcal{V}_\infty$ and the six proper subvarieties of \mathcal{V}_∞ .*

Proof. We begin by considering the situation where \mathcal{U} and \mathcal{W} are subvarieties of \mathcal{P} , $\mathcal{U} \subset \mathcal{W} \subseteq \mathcal{P}$ and \mathcal{U} and \mathcal{W} satisfy the same one-variable equations. Let $p = q$ be an equation that holds in \mathcal{U} but fails in \mathcal{W} . We shall argue that $p = q$ is \mathcal{P} -equivalent to some E_n . The purpose of this is to show that every subvariety of \mathcal{P} is axiomatizable relative to \mathcal{P} by a set of one-variable equations together with some of the E_n 's.

If both p and q are \mathcal{P} -equivalent to variables, then $p = q$ is \mathcal{P} -equivalent to $x = x$ or to $x = y$. Since $p = q$ fails in \mathcal{W} it cannot be equivalent to $x = x$. The equation $x = y$ is \mathcal{P} -equivalent to $x = 0$, which is a one-variable equation. Since $p = q$ holds in \mathcal{U} but not in \mathcal{W} , and \mathcal{U} and \mathcal{W} satisfy the same one-variable equations, we conclude that $p = q$ is not equivalent to $x = y$, either. Henceforth we assume that p is not \mathcal{P} -equivalent to a variable.

Now, according to Lemma 3.2, $p(\mathbf{x})$ is \mathcal{P} -equivalent to $p_1(x_1) + \dots + p_n(x_n)$ or the same with $+$ replaced by \oplus . According to which case we are in ($+$ or \oplus), this implies that either $ep = p$ or $fp = p$ is an equation of \mathcal{P} . If q is \mathcal{P} -equivalent to a variable then $p = q$ has the consequence $e(x) = x$, since $eq = ep = p = q$, or else the consequence $f(x) = x$. But $e(x) = x$ and $f(x) = x$ are equivalent modulo the equations of \mathcal{P} , so if q is \mathcal{P} -equivalent to a variable then \mathcal{U} satisfies the one-variable equation $e(x) = x$. \mathcal{W} must also satisfy this equation. However, the equations of \mathcal{P} together with $e(x) = x$ imply that $e = f = E = F$, $s = 0$ and $x + y = x \oplus y$. \mathcal{W} must now satisfy all of these equations and this is enough to imply that \mathcal{W} is a definitionally equivalent to a variety of Boolean groups. Since the variety of all Boolean groups is a minimal variety and \mathcal{U} is a proper subvariety of \mathcal{W} , therefore we must have that \mathcal{U} is the trivial variety. But this contradicts the assumption that \mathcal{U} and \mathcal{W} satisfy the same one-variable equations, since now \mathcal{U} satisfies $x = 0$ and \mathcal{W} does not. We conclude that q is not \mathcal{P} -equivalent to a variable. Therefore $q(\mathbf{x})$ is \mathcal{P} -equivalent to $q_1(x_1) + \dots + q_n(x_n)$ or the same with $+$ replaced by \oplus .

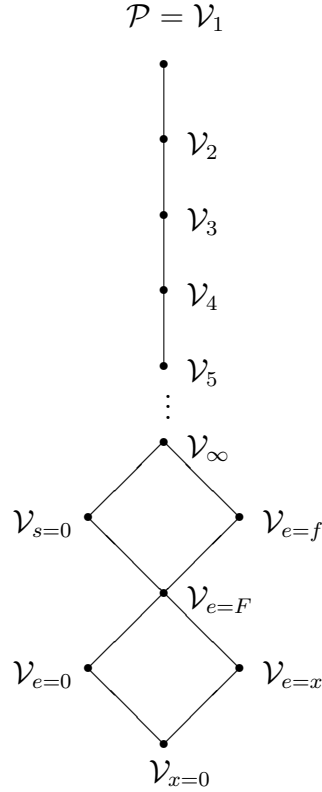
By substituting zeros into the equation $p = q$ we can see that for each i

$$p_i(x_i) := p(0, 0, \dots, x_i, \dots, 0) = q(0, 0, \dots, x_i, \dots, 0) =: q_i(x_i)$$

is a one-variable equation of \mathcal{U} and therefore of \mathcal{W} . It follows that each of p and q is \mathcal{P} -equivalent to either $p_1(x_1) + \dots + p_n(x_n)$ or $p_1(x_1) \oplus \dots \oplus p_n(x_n)$. Since $p = q$ fails to hold in \mathcal{W} , it must be that $p = q$ is \mathcal{P} -equivalent to

$$p_1(x_1) + \dots + p_n(x_n) = p_1(x_1) \oplus \dots \oplus p_n(x_n).$$

Because $x + y = e(x) + e(y)$ and $x \oplus y = f(x) \oplus f(y)$ hold in \mathcal{P} , it follows that $ep_i(x_i) = fp_i(x_i)$ is a one-variable equation of \mathcal{U} , therefore of \mathcal{W} . The only way for this to be true is if $p_i \in \{0, s(x_i)\}$ for all i . Since we may assume that each $p_i(x_i)$ depends on its variable, we may conclude that $p_i(x_i) = s(x_i)$ for all i . Thus, $p = q$ is \mathcal{P} -equivalent to E_n . We have shown that if $\mathcal{U} \subset \mathcal{W} \subseteq \mathcal{P}$ and \mathcal{U} and \mathcal{W} satisfy the


 FIGURE 1. Subvariety Lattice of \mathcal{P}

same one-variable equations, then \mathcal{U} is axiomatized relative to \mathcal{W} by a collection of the E_n 's. Thus every subvariety of \mathcal{P} is axiomatizable relative to \mathcal{P} by one-variable equations and some of the E_n 's.

The one-variable equations which fail to hold in \mathcal{P} are easy to locate since there are only seven \mathcal{P} -inequivalent unary terms: $\{0, x, e(x), f(x), s(x), E(x), F(x)\}$. It is a simple matter to show that each one-variable equation which fails in \mathcal{P} has either $s(x) = 0$ or $e(x) = f(x)$ as a consequence. The first clearly entails all E_n while the second entails $x + y = x \oplus y$ which clearly entails all E_n . Therefore, every one-variable equation which fails in \mathcal{P} entails all E_n . Combining this fact with what we have previously established, we obtain that any subvariety of \mathcal{P} which is not one of $\mathcal{P} = \mathcal{V}_1, \mathcal{V}_2, \dots$ or \mathcal{V}_∞ must be a subvariety of \mathcal{V}_∞ . Moreover, any subvariety of \mathcal{V}_∞ must be axiomatizable relative to \mathcal{P} by one-variable equations. Since there are so few nontrivial one-variable equations it is easy to determine that the subvarieties of \mathcal{V}_∞ are: $\mathcal{V}_{s=0}$, $\mathcal{V}_{e=f}$, $\mathcal{V}_{e=F}$, $\mathcal{V}_{e=x}$, $\mathcal{V}_{e=0}$, and $\mathcal{V}_{x=0}$. The notation $\mathcal{V}_{p=q}$ means that $\mathcal{V}_{p=q}$ is axiomatized by $p(x) = q(x)$ and the equations of \mathcal{P} . See Figure 1. \square

Lemma 3.6. \mathcal{V}_∞ has a six-element generator.

Proof. Borrowing notation from the previous proof, we must show that there is a six-element algebra in \mathcal{V}_∞ which is not in $\mathcal{V}_{s=0}$ or $\mathcal{V}_{e=f}$. That is, we must produce a six-element algebra \mathbf{A} for which

- (1) $\mathbf{A} \in \mathcal{P}$,
- (2) $\mathbf{A} \not\models s(x) = 0$,
- (3) $\mathbf{A} \not\models e(x) = f(x)$, and
- (4) $\mathbf{A} \models E_n$ for all n .

Observe that condition (4) says precisely that $s(A)$ generates a subgroup which is contained in $e(A) \cap f(A)$.

Let $\mathbf{B} = \mathbf{Z}_2^2$ and let $\mathbf{1} = (1, 1) \in B$. Let \mathbf{C} be the group obtained from \mathbf{B} by replacing the element $\mathbf{1}$ with a new element $\mathbf{1}'$. Let $Q = \{q\}$. Let $\iota : \mathbf{C} \rightarrow \mathbf{B}$ be the isomorphism which fixes $B \cap C$ and maps $\mathbf{1}'$ to $\mathbf{1}$. Define $e_Q(q) = s_Q(q) = (0, 1) \in B$. This yields a three-sorted structure which is associated to the algebra \mathbf{A} which has six-element universe $B \cup \{\mathbf{1}'\} \cup \{q\}$. Note that the subgroup generated by $s(A)$ is just $\{(0, 0), (0, 1)\} \subseteq e(A) \cap f(A)$. We have that $s(q) \neq 0$ and $e(\mathbf{1}) = \mathbf{1} \neq \mathbf{1}' = f(\mathbf{1})$ so the conditions listed above are met. \square

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