

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Introduction to Extremal and Probabilistic Combinatorics

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High Girth and High Chromatic

Random Regular

3-Flow Conjecture





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Definitions



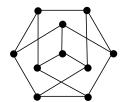
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Definition (Graph)

A graph is an ordered pair G = (V, E) consisting of a vertex set V and set of edges E (2-element subsets of V).





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Extremal Graph Theory

How much of something can you have, given a certain constraint?

Probabilistic Methods



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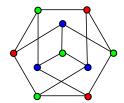
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Definition (Proper Coloring)

A proper coloring of G is an assignment of labels (or colors) to vertices such that no edge connects two vertices with the same color.





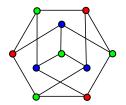
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Definition (k-coloring)

A proper coloring of G with k (or fewer colors) is a k-coloring.



Definition (Chromatic Number)

The chromatic number of G is the smallest k such that there is a k-coloring of G.





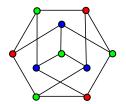
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Definition (Cycle of length *k*)

A cycle of length k consists of a closed walk (no repetitions of vertices or edges) through k vertices.



Definition (Girth)

The girth of a graph G is the length of a shortest cycle contained in G.

Observe: a triangle-free graph has girth ≥ 4 .





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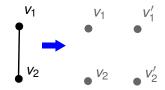




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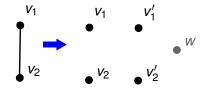




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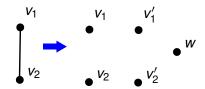




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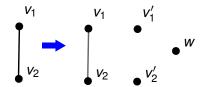




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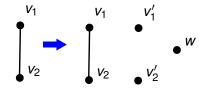




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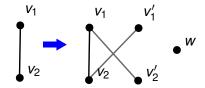




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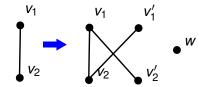




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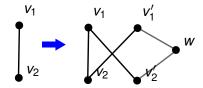




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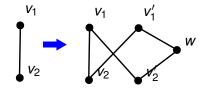




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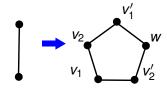




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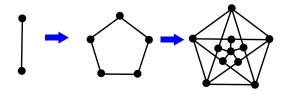




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High Girth and High Chromatic Number

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High Girth and High Chromatic Number



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What about for higher girth?

Can we find graphs with high girth and arbitrarily high chromatic number?

Yes, breakthrough using probabilistic combinatorics



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GRAPH THEORY AND PROBABILITY

P. ERDŐS

A well-known theorem of Ramay (θ_1 , θ_2) states that to every θ_1 there exist a smallest integer (θ_2) to that every graph of (g_1) vertices contains either a set of a independent points or a complete graph of order θ_1 , but there exists a graph of $(g_1^0) - 1$ vertices which does not contain a complete subgraph of a vertices and also does not contain a set of θ_1 independent points. (A graph is called complete if every two of its vertices are connected by an edge; a set of θ_2) of the containing of θ_3 or the containing of θ_3 or θ_4 or θ_3 or θ_4 or θ_3 or θ_4 or θ_4 or θ_3 or θ_4 or θ

(1)
$$2^{in} < g(n) < {2n-2 \choose n-1}$$

It is not even known that $g(n)^{1/6}$ tends to a limit. The lower bound in (1) has been obtained by combinatorial and probabilistic arguments without an explicit construction.

In our paper (5) with Szekeres f(k, l) is defined as the least integer so that every graph having f(k, l) vertices contains either a complete graph of order k or a set of l independent points (f(k, k) - g(k)). Szekeres proved

(2)
$$f(k, l) \le {k + l - 2 \choose k - 1}$$
.

Thus for

$$k=3, f(3,l) \leqslant \binom{l+1}{2}\,.$$

I recently proved by an explicit construction that $f(3,l)>l^{1+\epsilon_1}$ (4). By probabilistic arguments I can prove that for k>3

(3)
$$f(k, l) > l {k+l-2 \choose k-1}^{c_2}$$

which shows that (2) is not very far from being best possible.

Define now h(k, l) as the least integer so that every graph of h(k, l) vertices contains either a closed circuit of k or fewer lines, or that the graph contains a set of l independent points. Clearly h(3, l) = f(3, l).

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(4) $h(k, l) > l^{k+1/2k}$.

Further we shall prove that Received December 13, 1957.

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GRAPH THEORY AND PROBABILITY

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3-Flow Conjectur "It is not enough to be in the right place at the right time.

You should also have an open mind at the right time."

-Paul Erdős



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Theorem (Erdős 1959)

For any integers ℓ and k, there is a graph of girth $> \ell$ and chromatic number > k.

Idea: use random graphs

How do we generate random graphs on *n* vertices?



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3-Flow Conjectur 290

On random graphs I.

Dedicated to O. Varga, at the occasion of his 50° birthday. By P. ERDÓS and A. RÉNYI (Budapost).

Let us consider a "random graph" $\Gamma_{\sim N}$ having n possible (labelled) vertices and N edges; in other words, let us choose at random (with equal (fn 1)).

probabilities) one of the $\binom{\binom{n}{2}}{N}$ possible graphs which can be formed from

the n (tableted) vertices P_n, P_n, \dots, P_n by selecting N edges from the $\{j\}$ $\{j\}$ compositile edges, P_n $\{j\}$ ($i \le j \le j n$). Thus the effective number of vertices of P_n , was be less than n, as some points P_n may be not connected in P_n , with any other point P_n ; we shall call size objust P_n inducted points. We consider the isolated points also as belonging to P_n , x_n , x_n , is called completely connected it is deflectly content and a points P_n P_n , P_n (n). If it has no isolated points and is connected in the columny sense, in the present page we consider appropriate a subtracting properties of standard points P_n .

1. What is the probability of $\Gamma_{v,N}$ being completely connected?

What is the probability that the greatest connected component (subgraph) of F_{n,N} should have effectively n−k points? (k−0, 1, ...).
 What is the probability that f_{n,N} should consist of exactly k+1

3. What is the probability that \(\Gamma_{v,N} \) should consist of exactly \(k+1 \) connected components? \((k=0,1,...) \).
4. If the edges of a graph with \(n \) vertices are chosen successively so that after each step every edge which has not yet been chosen has the same

probability to be chosen as the next, and if we continue this process until the graph becomes completely connected, what is the probability that the number of necessary steps P will be equal to a given number I? As (partial) answers to the above questions we prove the following

four theorems. In Theorems 1, 2, and 3 we use the notation

 $N_c = \left[\frac{1}{2} n \log n + c n\right]$

where c is an arbitrary fixed real number ([x] denotes the integer part of x).





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Regular Graphs

3-Flow Conjectur "A mathematician is a device for turning coffee into theorems."

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"If I feel unhappy, I do mathematics to become happy.

If I am happy, I do mathematics to keep happy." —Alfréd Rényi



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$$n = 6$$

Model (Erdős and Rényi 1959)

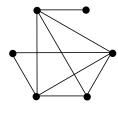
G(n,p) model (Erdős–Rényi model)

- Begin with n vertices.
- 2

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$$n = 6, p = \frac{1}{2}$$

Model (Erdős and Rényi 1959)

G(n, p) model (Erdős–Rényi model)

- Begin with n vertices.
- Include each edge independently with probability p.



High Girth and High Chromatic Number

Random Regular Graphs

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Theorem (Erdős 1959)

For any integers ℓ and k, there is a graph of girth $> \ell$ and chromatic number > k.

ldea:

- for n large and p carefully chosen, $G_{n,p}$ has "few" short cycles (at least half the time)
- for n large $G_{n,p}$ has high chromatic number (at least half the time)
- combining these and deleting some problem vertices we get graphs with high chromatic number and no short cycles



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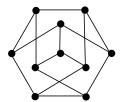
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Definition (Regular)

G is regular if all vertices have the same degree.



How do we generate random *d*-regular graphs on *n* vertices?



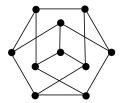
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Earse, J. Combingsories (1980) 1, 311-316

A Probabilistic Proof of an Asymptotic Formula for the Number of Labelled Regular Graphs

BÉLA BOLLOBÁS

Let Δ and a be entared numbers such that $\Delta a = 2m$ is even and $\Delta \leq (2 \log n)^{\frac{1}{2}} = 1$. Then $a : n = \infty$, the number of labelled Δ -regular graphs on n vertices is asymptotic to

$$e^{-n-kT} \frac{(2m)!}{m!2^m (4!)^{2m}}$$

where $\lambda=(\Delta-1)/2$. As a consequence of the method we determine the asymptotic distribution of the number of abort cycles is graphs with a given degree sequence, and give analogous formulae for hypergraphs.

In 1979 Next [6] determined as exact formula for the number of stabilité d -equit and [1], and except possible d is a final d -equit and d -equit d -equit

Theorem 1. Let $d_1 \cong d_2 \cong \cdots \cong d_n$ be notical numbers with $\sum_{i=1}^n d_i = 2m$ even.

 $A = d_1 \le (2 \log n)^4 - 1$ and $m \ge \max\{e.dn, (1 + e)n\}$ for some fixed $e \ge 0$. Then the number L(d) of labelled graphs with degree sequence $d = (d_1)$ satisfies

$$L(\mathbf{d}) = e^{-\lambda - \lambda^2} (2m)_m / \{2^m \prod_{i=1}^n d_i 1\},$$

where

constants

$$\lambda = \frac{1}{2\pi i} \sum_{i=1}^{n} \binom{d}{2}$$

PROOF. We shall represent our graphs as images of so called "configurations". Let $W=\bigcup_{i=1}^n W_i$ be a fixed set of $2m=\sum_{i=1}^n d_i$ abbelled vertices, where $|W_i|=d_i$. A configuration F is a partition of W into m pairs of vertices, calbed edges of F. Clearly there

$$N = N(m) = \left(\frac{2m!}{2}\right)\left(\frac{2m-2}{2}\right) \cdot \cdot \cdot \left(\frac{2}{2}\right) / m! = (2m)_m 2^{-m}$$
(1)
configurations. Furthermore, if we fix l independent (vertex disjoint) edges then there are

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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjectur Earse, J. Combinggories (1980) L. 311-316

A Probabilistic Proof of an Asymptotic Formula for the Number of Labelled Regular Graphs

BÉLA BOLLOBÁS

Let Δ and a be natural numbers such that $\Delta a = 2m$ is even and $\Delta \leq (2\log n)^{1} = 1$. Then as $n = \infty$, the number of labelled Δ -regular graphs on n vertices is asymptotic to

where $\lambda=(\Delta-1)/2$. As a consequence of the method we determine the asymptotic distribution of the number of short cycles is graphs with a given degree sequence, and give analogous bermulae for hypergraphs.

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Theorem 1. Let $d_1 \cong d_2 \cong \cdots \cong d_n$ be normal numbers with $\sum_{i=1}^n d_i = 2m$ even. Suppose $\Delta = d_1 \leqslant (2\log n)^i - 1$

and $m \ge \max\{edn, (1+e)n\}$ for some fixed e > 0. Then the number L(d) of labelled graphs with degree sequence $\mathbf{d} = (d_i)^n_1$ satisfies

$$L(\mathbf{d}) = e^{-\lambda^{-1}t}(2m)_{w} / \{2^{m} \hat{\prod} d_{t}!\}$$

where

constants.

$$\lambda = \frac{1}{2\pi} \hat{\Sigma} \begin{pmatrix} d \\ 2 \end{pmatrix}$$

PROOF. We shall represent our graphs as images of so called "configurations". Let $W = \bigcup_{i=1}^n W_i$ be a fixed set of $2m = \sum_{i=1}^n d_i$ labelled vertices, where $|W_i| = d_i$. A configuration F is a partition of W into m pairs of vertices, called edges of F. Clearly there

$$N = N(m) = \left(\frac{2m!}{2}\right)\left(\frac{2m-2}{2}\right) \cdot \cdot \cdot \left(\frac{2}{2}\right) / m! = (2m)_m 2^{-m}$$
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311

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Random Regular Graphs

3-Flow Conjectur "Erdős has an amazing ability to match problems with people.

Which is why so many mathematicians benefit from his presence."

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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture •

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$$n = 6$$

- Begin with n vertices.
- 2
- 3
- 4
- 5

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture











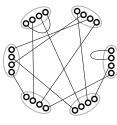
$$n = 6, d = 4$$

- Begin with n vertices.
- Create n "cells," each with d "points." (dn even)
- 3
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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



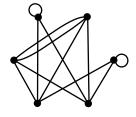
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- Begin with n vertices.
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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



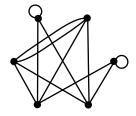
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High Girth and High Chromatic Number

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High Girth and High Chromatic Number

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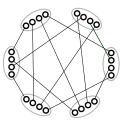
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Random Regular Graphs

3-Flow Conjectur

Bootstrap percolation on the random regular graph

József Balogh* and Boris G. Pittel †

December 8, 2005

Dedicated to Alan Frieze on the occasion of his 60-th birthday.

Abstract

The k-parameter bootstrap percolation on a graph is a model of an interacting particle system, which can also be viewed as a variant of a growth process of a cellular automata with threshold $k \ge 2$. At the start each of the graph vertices is active with probability p and inactive with probability 1 - p, independently of other vertices. Presence of active vertices triggers a percolation process controlled by the recursive rule: an active vertex remains active forever, and a currently inactive vertex becomes active when at least k of its neighbors are active. The basic problem is to identify, for a given graph, p^- , p^+ such that for $p < p^-$ ($p > p^+$ resp.) the probability that all vertices are eventually active is very close to 0 (1 resp.). The percolation process is a Markov chain on the space of subsets of the vertex set, which is easy to describe but hard to analyze rigorously in general. We study the percolation on the random d-regular graph, $d \ge 3$, via analysis of the process on its multigraph counterpart. Here, thanks to a "principle of deferred decisions", the percolation dynamics is described by a surprisinely simple Markov chain. Its generic state is formed by the counts of

1



^{*}University of Illinois at Urbana-Champaign, work was partially done while at the Ohio State University; email: jobal@math.uiuc.edu, research supported in part by NSF grant DMS-9039894 and OTKA grant 049399.

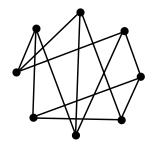
[†]The Ohio State University; email: bgp@math.ohio-state.edu, research supported in part by NSF grant DMS-0406024.



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Random Regular Graphs

3-Flow Conjectur



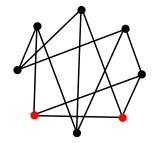
$$d = 3$$

Bootstrap Percolation on a Random Regular Graph

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjectur



$$d = 3, q = \frac{1}{4}$$

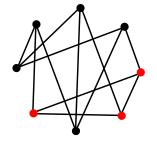
Bootstrap Percolation on a Random Regular Graph

 Infect vertices independently with some probability q

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjectur



$$d = 3, q = \frac{1}{4}$$

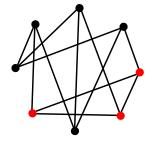
Bootstrap Percolation on a Random Regular Graph

- Infect vertices independently with some probability q
- Infection spreads to a vertex if > half neighbors are infected

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



$$d=3, q=\tfrac{1}{4}$$

- Infect vertices independently with some probability q
- Infection spreads to a vertex if > half neighbors are infected
- Iterate until stabilizes

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



$$d = 3, q = \frac{1}{4}$$

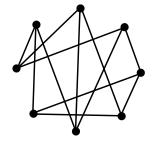
- Infect vertices independently with some probability q
- Infection spreads to a vertex if > half neighbors are infected
- Iterate until stabilizes
- Is the whole graph infected?



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Random Regular Graphs

3-Flow Conjecture



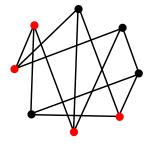
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3-Flow Conjecture



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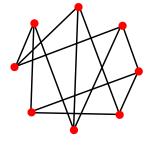
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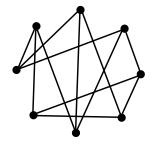
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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



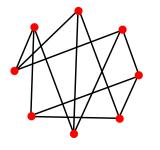
$$d = 3, q = 0$$

- Is the whole graph infected?
- For q = 0, no.

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture



$$d = 3, q = 1$$

- Is the whole graph infected?
- For q = 1, yes.

High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Theorem (Balogh and Pittel 2007)

Let $d \ge 3$. For random d-regular graphs, the dissemination threshold is a constant

$$p_d = \frac{d-2}{d-1}$$

asymptotically almost surely (a.a.s.).

$$p_3 = \frac{1}{2}, p_4 = \frac{2}{3},$$
 etc.

where an event X = X(n) holds a.a.s. if $\mathbb{P}[X(n)] \to 1$ as $n \to \infty$

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High Girth and High Chromatic

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High Girth and High Chromatic Number

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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture One of the most famous graph theory conjectures is the Tutte nowhere-zero 3-flow conjecture.

Conjecture (Equivalent Form, Tutte 1966)

Every 4-edge-connected, 5-regular graph has an edge orientation in which every out-degree is either 4 or 1.

Definition (k-edge-connected

G is k-edge-connected if G remains connected whenever any set of fewer than k edges are removed.



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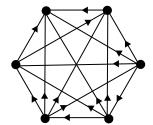
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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Question (Barát and Thomassen 2006)

Does every 4-edge-connected, 4-regular graph have an edge orientation in which every out-degree is either 4 or 1.

Answer: No!



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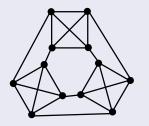


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3-Flow Conjecture

Example (Barát and Thomassen 2006)



Pigeonhole!

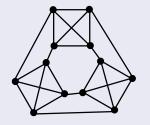


High Girth an High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

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High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Conjecture (Barát and Thomassen 2006)

If G is a planar 4-edge-connected, 4-regular graph such that 3|e(G), then G has an edge orientation in which every out-degree is either 4 or 1.



High Girth an High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Counterexample (Lai 2007)



High Girth and High Chromatic Number

Random Regular Graphs

3-Flow Conjecture

Theorem (Bollobás 1981, Wormald 1981)

A random d-regular graph is d-edge-connected asymptotically almost surely (a.a.s.).

Theorem (D. and Postle 2016+)

If 3|n, then a random 4-regular graph on n vertices has an edge orientation in which every out-degree is either 4 or 1 asymptotically almost surely (a.a.s.).



High Girth and High Chromatic Number

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Thank you for listening!