

Stable marriage with general preferences

Extended Abstract

Linda Farczadi, Konstantinos Georgiou, Jochen Könemann

University of Waterloo, Waterloo, Canada

Abstract. We propose a generalization of the classical stable marriage problem. In our model, the preferences on one side of the partition are given in terms of arbitrary binary relations, which need not be transitive nor acyclic. This generalization is practically well-motivated, and as we show, encompasses the well studied hard variant of stable marriage where preferences are allowed to have ties and to be incomplete. As a result, we prove that deciding the existence of a stable matching in our model is NP-complete. Complementing this negative result we present a polynomial-time algorithm for the above decision problem in a significant class of instances where the preferences are asymmetric. We also present a linear programming formulation whose feasibility fully characterizes the existence of stable matchings in this special case. Finally, we use our model to study a long standing open problem regarding the existence of cyclic 3D stable matchings. In particular, we prove that the problem of deciding whether a fixed 2D perfect matching can be extended to a 3D stable matching is NP- complete, showing this way that a natural attempt to resolve the existence (or not) of 3D stable matchings is bound to fail.

1 Introduction

The Stable Marriage (SM) problem is a classical bipartite matching problem first introduced by Gale and Shapley [8]. An instance of the problem consists of a set n of men, and a set of n women. Each man (woman) has a preference list that is a total order over the entire set of women (men). The goal is to find a stable matching between the men and women, meaning that there is no (man, woman) pair that both prefer each other to their current partners in the matching. Since its introduction, the stable marriage problem has become one of the most popular combinatorial problems with several books being dedicated to its study [10, 15, 21] and more recently [17]. The popularity of this model arises not only from its nice theoretical properties but also from its many applications. In particular, a wide array of allocation problems from many diverse fields can be analyzed within its context. Some well known examples include the labour market for medical interns, auction markets, the college admissions market, the organ donors market, and many more [21].

In their seminal work Gale and Shapley showed that every instance of SM admits a solution and such a solution can be computed efficiently using the so-called

Gale-Shapley (or man- proposing) algorithm. Among the many new variants of this classical problem, two extensions have received most of the attention: incomplete preference lists and ties in the preferences. Introducing either one of these extensions on its own does not pose any new challenges, meaning that solutions are still guaranteed to exist, all solutions have the same size, and they can be computed using a modification of the original Gale-Shapley algorithm [9, 10]. However, the same cannot be said about the Stable Marriage problem with Ties and Incomplete Lists (SMTI) that incorporates both extensions. In this variant stable matchings no longer need to be of the same size, even though they are still guaranteed to exist. In fact, deciding whether a given instance admits a stable matching of a given size is NP-hard [18], even in the case where ties occur only on one side of the partition. Several papers have studied the approximate variants of this problem (see [17] for a more complete account).

A central assumption in most variants of SM is that agents' preferences are transitive (i.e., if x is preferred to y , and y is preferred to z then x is also preferred to z). However, there are several studies [1, 6, 3, 19] that suggest that non-transitive, and even cyclic preferences arise naturally. Cyclicity, for example, may be introduced in the context of multi-attribute comparisons [7]; e.g., consider the following study from [19] where 62 college students were asked to make binary comparisons between three potential marriage partners x, y and z according to the following three criteria: intelligence, looks and wealth. The candidates had the following attributes: candidate x was very intelligent, plain, and well off; candidate y was intelligent, very good looking, and poor; and candidate z was fairly intelligent, good looking, and rich. From the 62 participants, 17 displayed the following cyclic preference: x was preferred to y , y was preferred to z , and z was preferred to x . In order to better capture such situations there is a need for a model that allows for more general preferences.

Addressing this need we propose the Stable Marriage with General Preferences (SMG) problem. As in SM, in an instance of SMG we are given n men, and n women, and the preferences of men are complete total orders over the set of women. The preferences of women, on the other hand, are given in terms of arbitrary binary relations over the men. Each of these binary relations will be represented by a set of ordered pairs of men. We say that a woman prefers man x at least as much as man y if the ordered pair (x, y) is part of her preference set. A matching is then stable as long as for every unmatched (man, woman) pair at least one member prefers her mate in the matching at least as much as the other member of the pair.

This introduction of non-transitive preferences, even when restricted to just one side of the partition, changes the properties of the model drastically. Like in SM any solution must be a perfect matching. However, solutions are no longer guaranteed to exist. We show that non-transitive preferences generalize both incomplete lists and ties by reducing the SMTI to SMG. In doing so, we prove that the SMG problem is also NP-hard. In addition, we provide results on the structural properties of the SMG problem and give sufficient conditions for the problem to be solvable in polynomial time.

The second half of this paper focuses on three-dimensional stable matching models whose study was initiated by Knuth [15]. We will be particularly interested in the Cyclic 3-Dimensional Stable Matching problem (c3DSM), where we are given a set of n men, a set of n women and in addition a set of n dogs. The preferences of the men are complete total orders over the set of women. Similarly the women have preferences over the dogs, and the dogs have preferences over the men. A 3D matching is said to be stable if there is no (man, woman, dog) triple that is strictly preferred to their current triples in the matching by each of its members. A prominent open question is whether every instance of c3DSM admits a stable matching, and whether it can be computed efficiently.

A natural avenue for attacking c3DSM is to solve the following problem which we refer to as Stable Extension (SE): suppose we fix a perfect matching M on dogs and men, can we efficiently determine whether M is extendible to a 3D stable matching? Recall that women have preferences over dogs only, but note that the given matching M induces preferences over their male owners as well! In essence, this allows us to state the SE problem as a two dimensional bipartite matching problem, and we show in Theorem 4 that SE can be seen as a special case of SMG. We then prove that SE remains NP-complete.

Contributions. In Section 3.1 we show the following result.

Theorem 1. *SMG is NP-complete.*

We then identify a significant class of instances that are solvable in polynomial time: those where the preferences are asymmetric, meaning that for every pair of men x, y , each woman prefers at most one to the other. We then prove the following result.

Theorem 2. *For instances of SMG with asymmetric preferences, there exists a polynomial time algorithm that finds a solution if and only if one exists.*

We provide two different proofs. The first (given in Section 3.2) employs an adaptation of the classical Gale-Shapley man-proposing algorithm. The second (given in Section 3.3) relies on a polyhedral characterization: we define a polytope that is non-empty if and only if the instance admits a stable matching. We also develop an efficient rounding algorithm for its fractional points. Despite displaying stronger structural properties than SMG, we show that SE remains hard to solve.

Theorem 3. *SE is NP-complete.*

The proof of the above theorem is given in Section 4. At a high level, its strategy resembles that of the proof of Theorem 1. The details are however significantly more intricate, mainly due to the fact that SE instances correspond to SMG instances in which preferences are induced by a given 3D matching instance. As an interesting consequence for the c3DSM, Theorem 3 rules out the natural algorithmic strategy of fixing and extending a 2D perfect matching on two of the input sets.

Related work. To the best of our knowledge, the stable marriage problem with non-transitive preferences has not been studied before. However, there is a rich literature about SM and its variants. In particular, there has been significant work concerning the approximation variant of the SMTI problem, where the goal is to find a maximum size stable matching. When ties are allowed on both sides, the problem is NP-hard to approximate within $33/29$ [23] and the currently best known ratio is $3/2$ [20]. When ties are only allowed on the side of the women the problem is NP-hard to approximate within $21/19$ [11] and the currently best known ratio is $25/17$ [14].

A related model known as Stable Marriage with Indifference [13, 16], allows for preferences to be given in the form of partial orders. In this case, it is possible for a pair of agents to be incomparable, but transitivity is still assumed to hold. This model allows for several definitions of stability, and depending on which definition is used, solutions might not always exist. However, there is an efficient algorithm for computing each type of stable matchings, whenever they exist.

For c3DSM, it is known that every instance admits a stable matching for $n \leq 4$ [4]. The authors conjectured that this result can be extended to general instances. In [2] it was shown that if we allow unacceptable partners, the existence of a stable matching becomes NP-complete. In the same paper, and also independently in [12], it was shown that the c3DSM problem under a different notion of stability known as strong stability is also NP-complete.

2 Definitions and notation

Throughout this paper we denote the set of men by B and the set of women by C . In the 3D setting, we have an additional set of n dogs A . A 3D perfect matching \mathcal{M} is a set of n disjoint triples from $A \times B \times C$. For every dog $a \in A$ we denote by $\mathcal{M}(a)$ the man that a is matched to in \mathcal{M} . Similarly for every man $b \in B$, $\mathcal{M}(b)$ denotes the woman that b is matched to in \mathcal{M} , and for every woman $c \in C$, $\mathcal{M}(c)$ denotes the dog that c is matched to in \mathcal{M} . A 3D perfect matching can also be induced by fixing perfect matchings on any two of the following sets $A \times B$, $B \times C$ or $C \times A$. We will use M to denote a perfect matching on $A \times B$ and N to denote a perfect matching on $B \times C$. We then define the 3D matching $M \circ N$ by setting $(a, b, c) \in M \circ N$ if and only if $(a, b) \in M$ and $(b, c) \in N$. For each $q \in A \cup B$ we denote by $M(q)$ the partner of q in M , and similarly for each $q \in B \cup C$ we denote by $N(q)$ the partner of q in N . If the preferences of an agent q are given in terms of an ordering $P(q)$ (with or without ties) over a set A then for all $x, y \in A$ we write $x \succ_q y$ to denote that q strictly prefers x to y and $x \succeq_q y$ to denote that q prefers x at least as much as y .

2.1 Stable marriage with ties and incomplete lists (SMTI)

An instance \mathcal{I} of SMTI consists of a set B of n men and a set C of n women. Each man $b \in B$ has a preference list $P(b)$ that is an ordering over a subset of C and is allowed to contain ties. Similarly each woman $c \in C$ has a preference

list $P(c)$ that is an ordering over a subset of B and is also allowed to contain ties. A pair (b, c) is said to be *acceptable* if b appears in $P(c)$ and c appears in $P(b)$. It is assumed that a woman c is acceptable to a man b if and only if man b is acceptable to woman c . A pair (b, c) is *blocking* with respect to a matching N if (b, c) is an acceptable pair that is not in N , $c \succ_b N(b)$ and $b \succ_c N(c)$. A matching N is *stable* if it uses only acceptable pairs and it has no blocking pairs. In that case, we also say that N is a solution to \mathcal{I} . If a solution is of size n we refer to it as a perfect stable matching. While all instances of SMTI admit a stable matching, not all instances admit a perfect stable matching. In this paper we use SMTI to refer to the decision problem of whether a given instance admits a perfect stable matching. This problem is known to be NP-complete [18], even when the ties occur only in the preference lists of the women.

2.2 Stable marriage with general preferences (SMG)

An instance \mathcal{I} of SMG consists of a set B of n men and a set C of n women. Each man $b \in B$ has a preference list $P(b)$ that is complete total order over C . Each woman $c \in C$ has a preference relation given in terms of a set of ordered pairs $\mathcal{R}_c \subseteq B \times B$. For a given pair of men $b, b' \in B$ and woman $c \in C$ we interpret $(b, b') \in \mathcal{R}_c$ as woman c preferring man b at least as much as man b' . Note that whether (b, b') is in \mathcal{R}_c is completely independent of whether $(b', b) \in \mathcal{R}_c$. We say that a pair (b, c) is *blocking* with respect to a matching N , if b and c are not matched to each other and neither one prefers its partner in N at least as much as the other. Formally, (b, c) is blocking if $(b, c) \notin N$, $c \succ_b N(b)$ and $(N(c), b) \notin \mathcal{R}_c$. A matching N is *stable* if it has no blocking pairs. It follows from this definition that any stable matching is a perfect matching. In this paper we use SMG to refer to the decision problem of whether a given instance admits a stable matching (which we also call a solution). An instance \mathcal{I} of SMG is said to have *asymmetric preferences* if for every $b_1, b_2 \in B$ and $c \in C$ at most one of the following two conditions holds: $(b_1, b_2) \in \mathcal{R}_c$ or $(b_2, b_1) \in \mathcal{R}_c$.

Note that we could have obtained an alternate definition of stability by saying that a pair (b, c) is blocking if $(b, c) \notin N$, $c \succ_b N(b)$ and $(b, N(c)) \in \mathcal{R}_c$. However, the two models are equivalent via the following correspondence: create a new instance \mathcal{I}' with sets \mathcal{R}'_c where $(b, b') \in \mathcal{R}'_c$ if and only if $(b', b) \notin \mathcal{R}_c$. Then the solutions that are stable for \mathcal{I} under the definition of stability used in this paper, are exactly those that are stable for \mathcal{I}' using the alternate definition of stability. Hence, we can use our definition of stability without loss of generality.

2.3 Stable extension (SE)

An instance \mathcal{I} of SE consists of a set of n dogs A , a set of n men B , and a set of n women C , together with a fixed perfect matching M on $A \times B$. The preferences are defined cyclically (A over B , B over C , and C over A) and are complete total orders over the corresponding sets. A triple (a, b, c) is *blocking* with respect to a 3D matching \mathcal{M} if $(a, b, c) \notin \mathcal{M}$, $b \succ_a \mathcal{M}(a)$, $c \succ_b \mathcal{M}(b)$ and $a \succ_c \mathcal{M}(c)$. Note that if (a, b, c) is a blocking triple then a, b and c must be part of three disjoint

triples in \mathcal{M} . A 3D matching \mathcal{M} is stable if it has no blocking pairs. It follows from this definition that any stable 3D matching must be a perfect matching. We say that a perfect matching N on $B \times C$ is a *stable extension*, or a solution to \mathcal{I} , if $M \circ N$ is a 3D stable matching, and we use SE to refer to the decision problem of whether a given instance admits a stable extension.

We now demonstrate how an instance \mathcal{I} of SE can be reduced to an SMG instance. First, for each man $b \in B$ we define A_b to be the set of dogs in A that prefer b to the man assigned to them in the fixed perfect matching M . That is $A_b = \{a \in A : b \succ_a M(a)\}$. The set A_b contains exactly those dogs in A with whom man b can potentially be in a blocking triple when extending M to a 3D matching. It follows that if $A_b = \emptyset$ then man b cannot be in a blocking triple in any extension of M to a 3D matching. Now, for each pair (b, c) we define $\alpha(b, c)$ to be the dog in the set A_b that woman c prefers the most. That is $\alpha(b, c) = \max_{\succ_c} A_b$. If $A_b = \emptyset$ then we let $\alpha(b, c)$ be the dog in the last position in woman c 's preference list. We now define preferences for each woman $c \in C$

$$\mathcal{R}_c := \{(b, b') \mid b, b' \in B, b \neq b', M(b) \succeq_c \alpha(b', c)\}. \quad (1)$$

Note that if $(N(c), b) \in \mathcal{R}_c$ then in the 3D matching $M \circ N$ the woman c will be matched to a dog that she prefers at least as much as any dog in A_b , therefore guaranteeing that the man b and woman c will never be part of the same blocking triple. Hence, in order for $M \circ N$ to be a 3D stable matching it suffices to ensure that for all $(b, c) \notin N$ we either have b matched to someone better than c , meaning $N(b) \succ_b c$, or we have c matched to some $N(c)$ such that $(N(c), b) \in \mathcal{R}_c$. But this is exactly the definition of a stable matching for an instance of SMG. Hence we have the following theorem.

Theorem 4. *SE can be reduced in polynomial time to SMG.*

Note that Theorem 4, together with Theorem 3 imply NP-hardness of SMG, and hence prove Theorem 1 (modulo containment in NP which is straightforward). Nevertheless, we choose to present first the proof of Theorem 1 as a warm-up, as it shares many similarity with that of Theorem 3.

3 Results for SMG

3.1 NP-completeness of SMG (Proof of Theorem 1)

Containment in NP is straightforward, and is based on the observation that deciding whether an edge not in a perfect matching of an instance of SMG is blocking or not can be done in polynomial time in n . The rest of our argument focuses on hardness.

Our proof uses a polynomial time reduction from SMTI where ties occur only in the preference lists of the women. This problem is known to be NP-complete [18]. Let \mathcal{I} be an instance of SMTI where ties occur only on the side of the women. We let $B = \{b_1, \dots, b_n\}$ denote the set of men and $C = \{c_1, \dots, c_n\}$

denote the set of women for the instance \mathcal{I} . For each person $q \in B \cup C$ we let $P(q)$ denote their preference list.

We now describe how to construct an instance \mathcal{J} of SMG. The set of men for our instance will be given by $B' = B \cup \{b_{n+1}, b_{n+2}\}$ and the set of women by $C' = C \cup \{c_{n+1}, c_{n+2}\}$. The preferences of the men are defined as follows: each original man $b \in B$ ranks the women in $P(b)$ first, in the same order as in $P(b)$, followed by the woman c_{n+1} , and the remaining women of C' ranked arbitrarily; each new man b_{n+i} for $i \in \{1, 2\}$ ranks the woman c_{n+i} first, and the remaining women of C' arbitrarily. Now, for each original woman $c \in C$ we define the binary relation $\mathcal{R}_c \subseteq B \times B$ as follows $\mathcal{R}_c := \{(b, b') \mid b, b' \in P(b), b \succeq_c b'\}$. That is, \mathcal{R}_c contains the ordered pair (b, b') when both b and b' are acceptable to c under the instance \mathcal{I} and c prefers b at least as much as b' . Finally, for each extra woman c_{n+i} for $i \in \{1, 2\}$ we set $\mathcal{R}_{c_{n+i}} := \emptyset$. This completes the definition of the instance \mathcal{J} . The new agents are used to establish the following property.

Lemma 1. *In any solution N to \mathcal{J} every man $b \in B$ is matched to a woman from the set $P(b)$.*

Note that each $b \in B$ ranks c_{n+1} immediately after all acceptable partners and before any unacceptable partners. Now if b is matched to someone outside of $P(b)$, the pair (b, c_{n+1}) will be blocking, since b will prefer c_{n+1} to its current partner, and c_{n+1} cannot be matched to someone she prefers at least as much as b since we defined $\mathcal{R}_{c_{n+1}} = \emptyset$. The full proof can be found in [5]. The following lemma completes the proof of Theorem 1.

Lemma 2. *\mathcal{I} admits a perfect stable matching if and only if \mathcal{J} admits a stable matching.*

Proof. Suppose that N is a perfect stable matching for \mathcal{I} . Then complete N to a perfect matching on $B' \cup C'$ by matching b_{n+i} to c_{n+i} for every $i \in \{1, 2\}$. To see that this is a stable matching for \mathcal{J} , note that both men b_{n+1} and b_{n+2} are matched to their most preferred woman in C' , hence they cannot be part of any blocking pairs. It remains to show that no man in B can be part of a blocking pair. Consider a man $b \in B$, and suppose that c is a woman that b strictly prefers to $N(b)$ according to the preferences in \mathcal{J} . Then it must be the case that $c \in P(b)$ and b also strictly prefers c to $N(b)$ in \mathcal{I} . Since N is a solution to \mathcal{I} , woman c prefers $N(c)$ at least as much as b in \mathcal{I} . Hence from the way we defined the set \mathcal{R}_c we have $(N(c), b) \in \mathcal{R}_c$, implying that (b, c) is not blocking.

To see the other direction suppose that \mathcal{J} admits a stable matching, and let N be the part of this stable matching obtained by restricting it to the sets $B \cup C$. It follows from Lemma 1 that N is a perfect matching and every man is matched to an acceptable woman. To see that there are no blocking pairs, consider any pair $(b, c) \notin N$ such that (b, c) is an acceptable pair, that is $b \in P(c)$ and $c \in P(b)$. Assume now that in \mathcal{I} man b strictly prefers c to $N(b)$. Since $c \in P(b)$ it follows that b also strictly prefers c to $N(b)$ in \mathcal{J} . Hence, since N is a stable matching for \mathcal{J} , we must have $(N(c), b) \in \mathcal{R}_c$. From the way we defined \mathcal{R}_c this implies that $N(c)$ is acceptable to c , and c prefers $N(c)$ at least as much as b in \mathcal{I} . Therefore N does not have any blocking pairs, and is a stable matching in \mathcal{I} .

3.2 Algorithmic results

In this section we introduce a variant of the Gale-Shapley man-proposing algorithm for instances of SMG that have asymmetric preferences. Let \mathcal{I} be an SMG instance as defined in Section 2.2. Like in the classical algorithm, each man in B is originally declared single and is given a list containing all the women of C in order of preference. In each round, every man b that is still single proposes to its most preferred woman in C that is still in his list. If a woman c accepts a proposal from a man b then they become engaged, and b 's status changes from single to engaged. On the other hand, if c rejects b 's proposal then b removes woman c from his list and remains single. The difference from the original Gale-Shapley algorithm is in the way that the women decide to accept or reject incoming proposals. A woman c accepts a proposal from a man b if and only if $(b, b') \in \mathcal{R}_c$ for all other men b' that have proposed to c up to that point in the algorithm. This will ensure that whenever a woman c rejects a proposal from a man b , c is guaranteed to be matched at the end of the algorithm to some b' such that $(b', b) \in \mathcal{R}_c$ therefore ensuring that (b, c) will not be a blocking pair. The description of the algorithm is given below.

Algorithm 1 A deferred acceptance algorithm for SMG

```
1: while there a single man in  $B$  with a non-empty list do
2:   for all  $b$  single with a non-empty list do
3:      $b$  proposes to the top  $c$  in its list
4:   for all  $b, c$  such that  $b$  proposed to  $c$  do
5:     if  $(b, b') \in \mathcal{R}_c$  for all  $b' \neq b$  that proposed to  $c$  then
6:        $c$  accepts  $b$  ( $(b, c)$  become an engaged pair)
7:     else
8:        $c$  rejects  $b$ 
9: If the set of engaged pairs forms a perfect matching return this solution, else conclude that  $\mathcal{I}$  does not admit a stable matching.
```

It is easy to see that Algorithm 1 terminates and runs in polynomial time since each man in B proposes to every woman in C at most once. The following lemma is easy to establish and its proof is found in [5].

Lemma 3. *Any solution returned by Algorithm 1 is a stable matching.*

Proof of Theorem 2. Using Lemma 3 it suffices to show that if Algorithm 1 does not find a solution then the given instance does not admit a stable matching. Suppose by contradiction that there exists a stable matching N but Algorithm 1 does not find a solution. First note that since the preferences are asymmetric, no woman $c \in C$ accepts more than one proposal at any point in the algorithm. Hence every woman is engaged to at most one man. Now since Algorithm 1 does not find a solution there is a man b that is rejected by every woman in C . In particular, b is rejected by $N(b)$. Among all pairs $(b, c) \in N$ such that b

proposed to c and c rejected b , let (b_0, c_0) be the one that corresponds to the earliest rejection. As observed earlier, no man is rejected because of an arbitrary choice. Therefore, if b_0 was rejected at c_0 then there must be some man b_1 that also proposed to c_0 and $(b_0, b_1) \notin \mathcal{R}_{c_0}$. Let c_1 be the partner of b_1 in N . We must have $c_1 \succ_{b_1} c_0$, since otherwise (b_1, c_0) would be a blocking pair for the stable matching N . Thus b_1 proposed to c_1 before proposing to c_0 and c_1 rejected b_1 before c_0 rejected b_0 . But this contradicts our choice of (b_0, c_0) . This concludes the proof of the theorem. \square

3.3 A polyhedral characterization

In this section we provide a polyhedral description for SMG, that is an analogue of the well studied stable marriage polytope, first introduced in [22]. It is well known that the latter polytope is integral, meaning that the optimization version of SM can be solved in polynomial time. For our setting, we show that our polytope can be used to efficiently decide the feasibility of an SMG instance with asymmetric preferences, thus giving an alternative proof of Theorem 2. We remark however that our polytope is not integral for this class of instances. Indeed, one can easily find instances with asymmetric preferences for which our polytope has fractional extreme points.

Given an instance \mathcal{I} of SMG we associated with each pair $(b, c) \in B \times C$ a variable x_{bc} , with the intended meaning that $x_{bc} = 1$ if b and c are matched to each other and $x_{bc} = 0$ otherwise. We then consider the following relaxation of the problem and let $P(\mathcal{I})$ denote the set of all vectors satisfying the constraints below

$$\sum_c x_{bc} = 1 \quad \forall b \in B \quad (2)$$

$$\sum_b x_{bc} = 1 \quad \forall c \in C \quad (3)$$

$$x_{bc} + \sum_{c' \succ_b c} x_{bc'} + \sum_{(b', b) \in \mathcal{R}_c} x_{b'c} \geq 1 \quad \forall b \in B, c \in C \quad (4)$$

$$x_{bc} \geq 0 \quad \forall b \in B, c \in C \quad (5)$$

It is easy to check that x is the incidence vector of a stable matching for \mathcal{I} if and only if x is an integer vector in $P(\mathcal{I})$. Our main result is the following.

Theorem 5. *Let \mathcal{I} be an instance of SMG with asymmetric preferences. Then $P(\mathcal{I}) \neq \emptyset$ if and only if \mathcal{I} admits a stable matching. Furthermore any fractional point $x \in P(\mathcal{I})$ can be efficiently rounded to a stable matching solution for \mathcal{I} .*

Proof. The first direction is trivial since if \mathcal{I} admits a stable matching then the incidence vector corresponding to this stable matching is clearly in P . Now assume $P \neq \emptyset$ and let x be any point in P . We will show how to efficiently round x to a stable matching, thus completing the proof of the theorem. For each $b \in B$ let $f(b)$ be b 's most preferred woman in the set $\{c \in C : x_{bc} > 0\}$.

Define $N = \{(b, f(b)) : b \in B\}$. We first show that N is a perfect matching. Since each man selects exactly one woman it suffices to show that no two men select the same woman. Suppose by contradiction that $f(b_1) = f(b_2) = c$ for some $b_1 \neq b_2$. Note that

$$\begin{aligned}
f(b) = c &\Rightarrow \sum_{c' \succ_b c} x_{bc'} = 0 \quad \text{from the definition of } f(b) \\
&\Rightarrow x_{bc} + \sum_{(b', b) \in \mathcal{R}_c} x_{b'c} \geq 1 \quad \text{from the stability constraint for } (b, c) \\
&\Rightarrow x_{bc} + \sum_{(b', b) \in \mathcal{R}_c} x_{b'c} = 1 \quad \text{from the matching constraint for } c \\
&\Rightarrow \sum_{(b', b) \notin \mathcal{R}_c} x_{b'c} = 0.
\end{aligned}$$

Therefore $f(b_1) = f(b_2) = c$ implies that $(b_1, b_2) \in \mathcal{R}_c$ and $(b_2, b_1) \in \mathcal{R}_c$. But this contradicts the assumption that the preferences are asymmetric. Hence N must be a perfect matching. To see that N satisfies the stability constraints consider any pair (b, c) . If $f(b) = c$ or $f(b) \succ_b c$ then (b, c) cannot be blocking, since b will be matched in N to someone he prefers at least as much as c . Hence it suffices to consider the case where $c \succ_b f(b)$. But then we must have $x_{bc} + \sum_{c' \succ_b c} x_{bc'} = 0$ and since $x \in P(\mathcal{I})$ this implies that $\sum_{(b', b) \in \mathcal{R}_c} x_{b'c} = 1$. Now, since N uses only edges in the support of x , it follows that $(N(c), b) \in \mathcal{R}_c$ and hence (b, c) is not blocking. Therefore N is a stable matching for \mathcal{I} . \square

4 NP-completeness of SE (Proof outline of Theorem 3)

In this section we outline the major ideas behind the proof of Theorem 3. The full version of the proof can be found in [5].

At a high level, our proof adopts a similar strategy as that of Theorem 1. In particular, we will provide a polynomial-time reduction for SMTI to SE. However, due to the additional structural properties of SE instances this reduction becomes significantly more intricate. As in the proof of Theorem 1 we start with an instance \mathcal{I} of SMTI where ties occur only on the side of the women. We then create an instance \mathcal{J} of SE that will consist of three sets A' , B' and C' . We will have $B \subset B'$ and $C \subset C'$. We will also define some additional agents and fix a perfect matching M on $A' \times B'$. Our goal will be to show that \mathcal{I} admits a perfect stable matching if and only if \mathcal{J} admits a stable extension. Using Theorem 4, we can show that \mathcal{J} admits a stable extension by arguing instead that the corresponding instance of SMG defined from \mathcal{J} admits a stable matching.

The major difference from the proof of Theorem 1 is that there we had the freedom of defining the sets \mathcal{R}_c directly, but here the sets \mathcal{R}_c are implicitly defined through the structure of the SE instance \mathcal{J} . This implicit definition entails a significant amount of inter-dependence between the \mathcal{R}_c sets. In particular it

now becomes more difficult to establish the following property of the sets \mathcal{R}_c : for every acceptable pair (b_j, c_i) of the instance \mathcal{I} we have $(b, b_j) \in \mathcal{R}_{c_i}$ if and only if b is a man that c_i prefers at least as much as b_j in \mathcal{I} . This property was crucial for establishing a correspondence between the pairs that are blocking in \mathcal{I} and those that are blocking in \mathcal{J} . In order to establish this property here we will need to introduce new agents in \mathcal{J} corresponding to each position of woman c_i 's preference list $P(c_i)$ from \mathcal{I} . In particular if $P(c_i)$ consists of t_i positions, then there will be t_i new dogs $a_{i,1}, \dots, a_{i,t_i}$. We will then define the preference of the dogs as well as the fixed perfect matching M in such a way as to ensure that if man b_j appears in position k in $P(c_i)$ then the only dog from the set $\{a_{i,1}, \dots, a_{i,t_i}\}$ that prefers man b_j over his partner in M will be $a_{i,k}$. Finally we will define the preference of woman c_i over the set of dogs in \mathcal{J} as follows: c_i will rank the dogs that are matched in M to men from the first position in $P(c_i)$ at the top of its list, in any arbitrary order among them, followed by the dog $a_{i,1}$; then the dogs that are matched in M to men from the second position in $P(c_i)$ followed by the dog $a_{i,2}$, and so on until the dogs that are matched in M to men from the last position in $P(c_i)$ followed by the dog a_{i,t_i} . The remainder of c_i 's preference list will be completed arbitrarily. We will then show how this will imply the desired structure for the set \mathcal{R}_{c_i} .

The other major difference from the proof of Theorem 1 lies in ensuring that the new agents will be matched between themselves in any solution to \mathcal{J} . Before, if we wanted to ensure that a given pair b_{n+1}, c_{n+1} are always matched to each other, it was sufficient to set c_{n+1} as the most preferred woman of b_{n+1} and $\mathcal{R}_{c_{n+1}} = \emptyset$. However, since our instance of SMG arises from the SE instance \mathcal{J} the sets \mathcal{R}_c are never empty. Indeed it follows from their definition that $(M(\alpha(b, c), b) \in \mathcal{R}_c$ for every pair $b \in B$ and $c \in C$ we have. Thus, in order to ensure that a given pair b_{n+1}, c_{n+1} are always matched to each other we now need to introduce an extra gadget consisting of two new men b_{n+2}, b_{n+3} and two new women c_{n+2}, c_{n+3} . We will have b_{n+i} rank c_{n+i} first in its preference list for all $i \in \{1, 2, 3\}$. It is then possible to ensure that $(b, b_{n+1}) \in \mathcal{R}_{c_{n+1}}$ if and only if $b = b_{n+2}$, $(b, b_{n+3}) \in \mathcal{R}_{c_{n+3}}$ if and only if $b = b_{n+2}$, and $(b, b_{n+2}) \in \mathcal{R}_{c_{n+2}}$ if and only if $b = b_{n+3}$. This will then guarantee that in any solution to \mathcal{J} the men b_{n+2} and b_{n+3} are matched to the women c_{n+2} and c_{n+3} (in any of the two possible ways) and that man b_{n+1} is always matched to woman c_{n+1} as desired. The full details of the proof can be found in [5].

Acknowledgement. We thank the reviewers for their thorough and careful review. We highly appreciate their insightful suggestions that led to a substantial improvement of the paper.

References

- [1] Birnbaum, M., Schmidt, U.: An experimental investigation of violations of transitivity in choice under uncertainty. *J Risk Uncertain* 37(1), 77–91 (2008)

- [2] Biró, P., McDermid, E.: Three-sided stable matchings with cyclic preferences. *Algorithmica* 58(1), 5–18 (2010)
- [3] Brams, S.J., Gehrlein, W.V., Roberts, F.S.: *The Mathematics of Preference, Choice and Order*. Springer (2009)
- [4] Eriksson, K., Sjöstrand, J., Strimling, P.: Three-dimensional stable matching with cyclic preferences. *Math. Soc. Sci.* 52(1), 77–87 (2006)
- [5] Farczadi, L., Georgiou, K., Könemann, J.: Stable marriage with general preferences. arXiv preprint arXiv:1407.1853 (2014)
- [6] Fishburn, P.C.: Nontransitive preferences in decision theory. *J Risk Uncertain* 4(2), 113–134 (1991)
- [7] Fishburn, P.C.: Preference structures and their numerical representations. *Theoretical Computer Science* 217(2), 359–383 (1999)
- [8] Gale, D., Shapley, L.S.: College admissions and the stability of marriage. *Amer. Math. Monthly* pp. 9–15 (1962)
- [9] Gale, D., Sotomayor, M.: Some remarks on the stable matching problem. *Discrete Appl. Math.* 11(3), 223–232 (1985)
- [10] Gusfield, D., Irving, R.W.: *The stable marriage problem: structure and algorithms*, vol. 54. MIT Press (1989)
- [11] Halldórsson, M.M., Iwama, K., Miyazaki, S., Yanagisawa, H.: Improved approximation results for the stable marriage problem. *TALG* 3(3), 30 (2007)
- [12] Huang, C.C.: Circular stable matching and 3-way kidney transplant. *Algorithmica* 58(1), 137–150 (2010)
- [13] Irving, R.W.: Stable marriage and indifference. *Discrete Appl. Math.* 48(3), 261–272 (1994)
- [14] Iwama, K., Miyazaki, S., Yanagisawa, H.: A 25/17-approximation algorithm for the stable marriage problem with one-sided ties. *Algorithmica* 68(3), 758–775 (2014)
- [15] Knuth, D.: *Stable marriage and its relation to other combinatorial problems: An introduction to the mathematical analysis of algorithms*. Amer. Math. Soc., Providence, RI (1997)
- [16] Manlove, D.F.: The structure of stable marriage with indifference. *Discrete Appl. Math.* 122(1), 167–181 (2002)
- [17] Manlove, D.F.: *Algorithmics of matching under preferences*. World Sci. Publishing (2013)
- [18] Manlove, D.F., Irving, R.W., Iwama, K., Miyazaki, S., Morita, Y.: Hard variants of stable marriage. *Theor. Comput. Sci.* 276(1), 261–279 (2002)
- [19] May, K.O.: Intransitivity, utility, and the aggregation of preference patterns. *Econometrica* pp. 1–13 (1954)
- [20] McDermid, E.: A 3/2-approximation algorithm for general stable marriage. In: *Automata, Languages and Programming*, pp. 689–700. Springer (2009)
- [21] Roth, A.E., Sotomayor, M.: *Two-sided matching: A study in game-theoretic modeling and analysis*. No. 18, Cambridge University Press (1992)
- [22] Vande Vate, J.H.: Linear programming brings marital bliss. *Oper. Res. Lett.* 8(3), 147–153 (1989)
- [23] Yanagisawa, H.: *Approximation algorithms for stable marriage problems*. Ph.D. thesis, Citeseer (2007)