±1 QUADRATIC PROGRAMMING

Formulations

$$(P) \ \ \mu^* := \max \ q(x) := x^t Q x + c^t x, \ x \in F := \{-1, 1\}^n,$$

 $Q n \times n$ symmetric matrix, $c \in \mathbb{R}^n$.

Relaxations, $K \supset F$

$$(RP) \quad f(u) = \max_{x \in K} q_u(x) := x^t (Q - \operatorname{diag}(u))x + u^t e + c^t x,$$

Solve Tractable Problem $B := \min_{u \in L} f(u)$ Upper Bound $\mu^* \leq B$

Three Relaxations

relaxed problem RP1

$$(RP_v^1) \quad f_1(v) := \max_{-1 \le x \le 1} q_v(x).$$

bound $B_1 := \min_{Q-\operatorname{diag}(v) \leq 0} f_1(v)$.

relaxed problem RP2

$$(RP_u^2)$$
 $f_2(u) := \max_{||y||^2 = n} q_u(y).$

bound $B_2 := \min_{u^t e=0} f_2(u) = \min_u f_2(u)$.

Bound 3

$$Q^c := egin{bmatrix} 0 & rac{1}{2}c^t\ rac{1}{2}c & Q \end{bmatrix} \cdot$$
 $q^c_u(y) := y^t(Q^c - ext{diag}(u))y + u^t\epsilon$

relaxed problem RP3

$$(RP_u^3) \quad \begin{array}{rcl} f_3(u) & := & \max_{||y||^2 = n+1} q_u^c(y) \ & = & (n+1)\lambda_{\max}(Q^c - \operatorname{diag}(u)) + u^t e_{x} \end{array}$$

bound $B_3 := \min_{u^t e=0} f_3(u) = \min_u f_3(u)$.

BOUND 1 - Convex Quadratic Programming

Consider the shifted function

$$q_{\nu}(x) := x^t(Q - \operatorname{diag}(\nu))x + \nu^t e + c^t x,$$

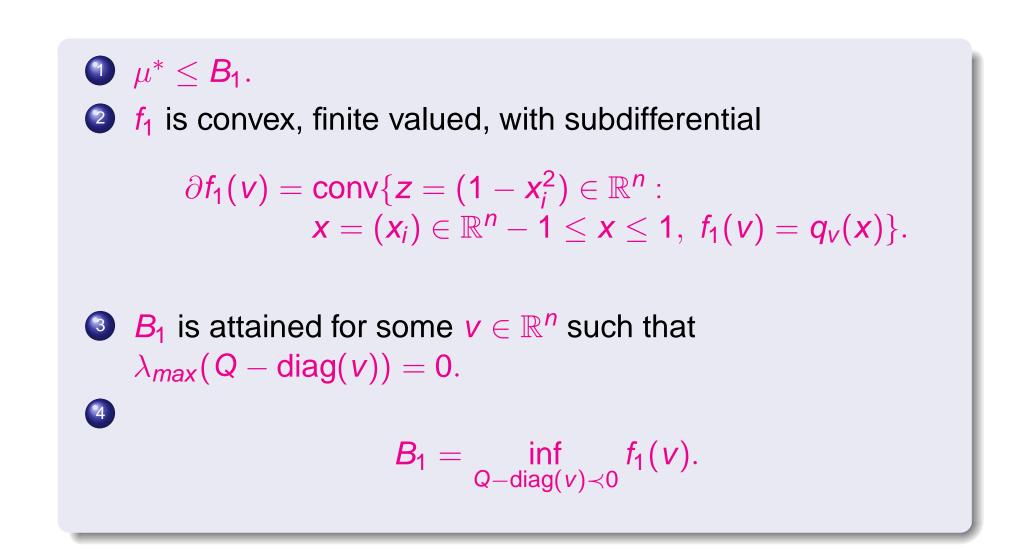
and the relaxed problem

$$(RP_{v}^{1}) \quad f_{1}(v) := \max_{-1 \le x \le 1} q_{v}(x).$$

Then a bound for (P) is

$$B_1 := \min_{\substack{Q-\text{diag}(v) \leq 0}} f_1(v).$$

Properties for Bound B_1 .



Optimality conditions

Lagrangian

$$L_1(v, \Lambda) := f_1(v) + \operatorname{trace} \Lambda(Q - \operatorname{diag}(v)), \qquad \Lambda \succeq 0$$

Optimality conditions:

 $\begin{array}{ll} 0 \in \partial f_1(v) - \operatorname{diag}(\Lambda) & (\text{stationarity}) \\ \operatorname{trace} \Lambda(Q - \operatorname{diag}(v)) = 0 & (\text{complementary slack}) \\ \Lambda \succeq 0 & (\text{multiplier sign}). \end{array}$

BOUND 2 - Optimization Over Sphere

$$q_u(y) := y^t(Q - \operatorname{diag}(u))y + u^t e + c^t y,$$

the second relaxed problem

$$(RP_u^2) \quad f_2(u) := \max_{||y||^2 = n} q_u(y).$$

Now a bound for (P) is

$$B_2 := \min_{u^t e=0} f_2(u) = \min_{u} f_2(u).$$

Properties for bound B_2

µ^{*} ≤ B₂.
f₂ is convex, finite valued, with subdifferential
∂f₂(u) = conv{z = (1 - y_i²) ∈ ℝⁿ : y = (y_i) ∈ ℝⁿ, ||y||² = n, f₂(u) = q_v(y)}.
The bound B₂ is attained for some u ∈ ℝⁿ. Moreover, if B₂ > μ^{*}, then the hard case holds for (RP²_u).

THEOREM: The bound $B_1 = B_2$

In fact

$$egin{array}{rll} f_2(u)&=&q_u(y_u)\ &\geq&\max_{(-1\leq y\leq 1)}q_{(u+\lambda e)}(y)\ &=&f_1(u+\lambda e) \end{array}$$

and

 $B_2 = f_2(u)$, with Lagrange multiplier λ for (RP_u^2) if and only if $B_2 = f_2(u) = f_1(u + \lambda e) = B_1$.

THEOREM:

$$B_1=B_2=B_3=B_4$$

Suppose that $f_4(v) := \max_{1 \le x \le 1} q_v^c(x)$ and

$$B_4 := \min_{\substack{Q^c - \operatorname{diag}(v) \leq 0}} f_4(v).$$

Then

$$B_1 = B_2 = B_3 = B_4.$$

 $(B_{53} - \alpha)$

Corollary

Let *B* be an $(n - 1) \times (n - 1)$ real symmetric matrix, and consider the perturbation of *B* given by

±1 Quadratic Programming Problem

(P) $\mu^* := \max q(x) \quad x \in F \cap S$

where: $F = \{\pm 1\}^n$, $S \subset \mathbb{R}^n$, $F \cap S \neq \emptyset$ quadr. obj.: $q(x) := x^t Q x - 2c^t x$, $Q \in S^n$, $c \in \mathbb{R}^n$

EXAMPLES

Quadratic Assignment Problem, QAP, (in trace formulation)

$$(QAP) \max_{X \in \Pi} q(X) := \operatorname{trace}(AXB - 2C)X^{t}$$

Max-Cut Problem, MC

$$(MC) \max \frac{1}{2} \sum_{i < j} w_{ij} (1 - x_i x_j), x \in F$$

HIDDEN SEMIDEFINITE CONSTRAINTS

Trust Region Subproblem (TRS)

$$\mu^* = \min_{x} q(x) \text{ s.t. } x^t x = s^2 (\leq s^2)$$

- $= \min_{\boldsymbol{x}} \max_{\boldsymbol{\lambda}} L(\boldsymbol{x}, \boldsymbol{\lambda})$
- $\geq \max_{\lambda} \min_{\mathbf{x}} L(\mathbf{x}, \lambda)$
- $= \max_{Q-\lambda \succeq 0} \min_{x} L(x,\lambda)$

$$= \max_{\mathsf{Q}-\lambda\succeq 0} \textit{h}(\lambda) = \mu^*$$

where

 $x_{\lambda} = (Q - \lambda I)^{\dagger} c;$ $h(\lambda) = L(x, \lambda) = -c^{t}(Q - \lambda I)^{\dagger} c + \lambda s^{2}$ nonconvex objective but min-max = max-min hidden convexity provides the hidden convex dual program

Relaxations of ± 1 -Quadratic Progr.

use perturbations

$$q_u(x) := q(x) + x^t \operatorname{Diag}(u)x - u^t e$$

Relaxation 0:

$$f_0(u) := \max_x q_u(x)$$

$$\mu^* \leq B_0 := \min_{u^t e=0} f_0(u) = \min_u f_0(u)$$
$$= \min_{Q+\text{Diag}(u) \leq 0} f_0(u)$$

 B_0 is Lagrangian dual of original quadratic program in the following form with u_i as Lagrange multipliers

$$\min q(x) \text{ s.t. } x_i^2 = 1, \forall i$$

Relaxation 1

sphere of radius \sqrt{n}

$$f_1(u) := \max_{||x||^2 = n} q_u(x)$$

$$\mu^* \leq B_1 := \min_{u^t e = 0} f_1(u) = \min_{u} f_1(u)$$

$$B_1 = \min_u \max_{x^t x = n} q_u(x)$$

$$= \min_{u,\lambda} \max_{x} q_u(x) + \lambda(x^t x - n)$$

$$= \min_{v^t e=0} \max_{x} q_v(x), \text{ with } v = u + \lambda e$$

$$= B_0$$

Relaxation 2

Unit Box

$$f_1(u) := \max_{|x_i| \le 1} q_u(x)$$

$$\mu^* \leq B_2 := \min_{Q + \operatorname{Diag}(u) \leq 0} f_2(u) = \min_u f_2(u)$$

$$B_2 = \min_{\substack{u \ x_i^2 \le 1}} \max_{\substack{x_i^2 \le 1}} q_u(x)$$

= $\min_{\substack{u \ \lambda \ge 0}} \max_{\substack{x \ x}} q_u(x) + \sum_{\substack{i \ \lambda}} \lambda_i (1 - x_i^2)$
= B_0 after $v = u - \lambda$

i.e. same bound again

Relaxation 1^c

Homogenization; sphere radius = $\sqrt{n+1}$

$$\begin{aligned} Q^c &:= \begin{bmatrix} 0 & -c^t \\ -c & Q \end{bmatrix}; \quad q^c_u(y) &:= y^t (Q^c + \operatorname{diag}(u))y - u^t e \\ f^c_1(u) &:= \max_{||y||^2 = n+1} q^c_u(y) \\ &= (n+1)\lambda_{\max}(Q^c + \operatorname{diag}(u)) - u^t e \end{aligned}$$

$\mu^* \leq B_1^c := \min_{u^t e=0} f_1^c(u) = \min_u f_1^c(u)$

$$B_{1}^{c} = \min_{v} \max_{y^{t}y=n+1} q_{v}^{c}(y) = \min_{v} \max_{y} q_{v}^{c}(y)$$

= $\min_{u,u_{0}} \max_{x,x_{0}} u_{0}(x_{0}^{2}-1) + x^{t}(Q + \text{Diag}(u))x$
 $-2x_{0}c^{t}x - u^{t}e$

$$= \min_{u} \max_{x, x_0^2 = 1} x^t (Q + \text{Diag}(u)) x - 2x_0 c^t x - u^t e$$
$$= B_0$$

i.e. same bound again; similarly for B_2^c

Relaxation 3

Semidefinite SDP (c = 0)

$$q(x) = x^t Q x = \text{trace } Q x x^t, \qquad Y = x x^T$$

 $B_3 := \max \text{trace } Q Y \text{ s.t } \operatorname{diag}(Y) = e, Y \succeq 0$

The dual is

$$B_{3} = \min y^{t} e \text{ s.t. } Q - \text{Diag}(y) \leq 0$$

$$Q - \text{Diag}(y - \frac{e^{t}y}{n}e) \leq \frac{e^{t}y}{n}I, \text{ with } w = y - \frac{e^{t}y}{n}e \text{ and } z = \frac{e^{t}y}{n}$$

$$B_{3} = \min i mize \qquad nz$$
subject to
$$Q - \text{Diag}(w) \leq zI$$

$$w^{t}e = 0$$

So: $B_3 = B_1^c$, i.e. all bounds are equal to Lagrangian relaxation of the equivalent quadratically constrained program. The SDP relaxation is the dual of the Lagrangian dual.

Allow more General Perturbations?

$q_{V,d}(x) := x^t (Q + V) x + (c + d)^t x$

THEOREM: Suppose that $q_{V,d}(x) \ge q(x)$, $\forall x \in F$. Then

V = P + U, with $P \succeq 0$, U is diagonal, and trace U = 0.

Moreover, there exists a diagonal matrix W, with trace W = 0, such that

$$\max_{x} q_{V,d}(x) \geq \max_{x} q_{W,0}(x).$$

Therefore, we need only consider diagonal perturbations, i.e. we have the best quadratic approximation - by duality we have the best SDP relaxation.