1. Introduction

- mathematical optimization
- least-squares and linear programming
- convex optimization
- example
- course goals and topics
- nonlinear optimization
- brief history of convex optimization

Mathematical optimization

(mathematical) optimization problem

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, \quad i = 1, \dots, m$

- $x = (x_1, \dots, x_n)$: optimization variables
- $f_0: \mathbf{R}^n \to \mathbf{R}$: objective function
- $f_i: \mathbf{R}^n \to \mathbf{R}, i = 1, \dots, m$: constraint functions

optimal solution x^* has smallest value of f_0 among all vectors that satisfy the constraints

Examples

portfolio optimization

- variables: amounts invested in different assets
- constraints: budget, max./min. investment per asset, minimum return
- objective: overall risk or return variance

device sizing in electronic circuits

- variables: device widths and lengths
- constraints: manufacturing limits, timing requirements, maximum area
- objective: power consumption

data fitting

- variables: model parameters
- constraints: prior information, parameter limits
- objective: measure of misfit or prediction error

Solving optimization problems

general optimization problem

- very difficult to solve
- ullet methods involve some compromise, e.g., very long computation time, or not always finding the solution

exceptions: certain problem classes can be solved efficiently and reliably

- least-squares problems
- linear programming problems
- convex optimization problems

Least-squares

minimize
$$||Ax - b||_2^2$$

solving least-squares problems

- analytical solution: $x^* = (A^T A)^{-1} A^T b$
- reliable and efficient algorithms and software
- computation time proportional to n^2k $(A \in \mathbf{R}^{k \times n})$; less if structured
- a mature technology

using least-squares

- least-squares problems are easy to recognize
- a few standard techniques increase flexibility (e.g., including weights, adding regularization terms)

Linear programming

minimize
$$c^T x$$

subject to $a_i^T x \leq b_i, \quad i = 1, \dots, m$

solving linear programs

- no analytical formula for solution
- reliable and efficient algorithms and software
- computation time proportional to n^2m if $m \ge n$; less with structure
- a mature technology

using linear programming

- not as easy to recognize as least-squares problems
- a few standard tricks used to convert problems into linear programs (e.g., problems involving ℓ_1 or ℓ_∞ -norms, piecewise-linear functions)

Convex optimization problem

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, \quad i = 1, \dots, m$

• objective and constraint functions are convex:

$$f_i(\alpha x + \beta y) \le \alpha f_i(x) + \beta f_i(y)$$

if
$$\alpha + \beta = 1$$
, $\alpha \ge 0$, $\beta \ge 0$

• includes least-squares problems and linear programs as special cases

solving convex optimization problems

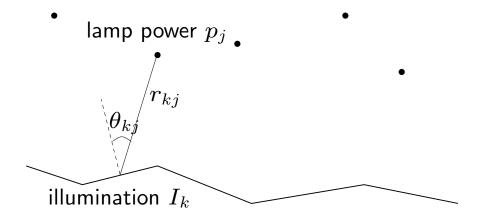
- no analytical solution
- reliable and efficient algorithms
- computation time (roughly) proportional to $\max\{n^3, n^2m, F\}$, where F is cost of evaluating f_i 's and their first and second derivatives
- almost a technology

using convex optimization

- often difficult to recognize
- many tricks for transforming problems into convex form
- surprisingly many problems can be solved via convex optimization

Example

m lamps illuminating n (small, flat) patches



intensity I_k at patch k depends linearly on lamp powers p_j :

$$I_k = \sum_{j=1}^m a_{kj} p_j, \qquad a_{kj} = r_{kj}^{-2} \max\{\cos \theta_{kj}, 0\}$$

problem: achieve desired illumination I_{des} with bounded lamp powers

minimize
$$\max_{k=1,...,n} |\log I_k - \log I_{\text{des}}|$$
 subject to $0 \le p_j \le p_{\text{max}}, \quad j=1,\ldots,m$

how to solve?

- 1. use uniform power: $p_j = p$, vary p
- 2. use least-squares:

minimize
$$\sum_{k=1}^{n} (I_k - I_{des})^2$$

round p_j if $p_j > p_{\text{max}}$ or $p_j < 0$

3. use weighted least-squares:

minimize
$$\sum_{k=1}^{n} (I_k - I_{\text{des}})^2 + \sum_{j=1}^{m} w_j (p_j - p_{\text{max}}/2)^2$$

iteratively adjust weights w_j until $0 \le p_j \le p_{\text{max}}$

4. use linear programming:

$$\begin{array}{ll} \text{minimize} & \max_{k=1,\ldots,n} |I_k - I_{\text{des}}| \\ \text{subject to} & 0 \leq p_j \leq p_{\text{max}}, \quad j=1,\ldots,m \\ \end{array}$$

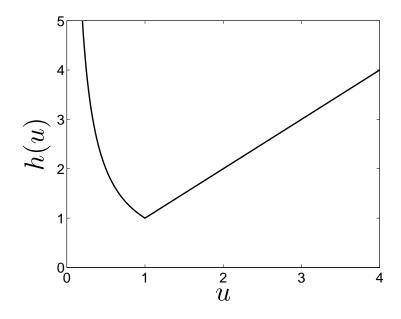
which can be solved via linear programming of course these are approximate (suboptimal) 'solutions'

5. use convex optimization: problem is equivalent to

minimize
$$f_0(p) = \max_{k=1,...,n} h(I_k/I_{\text{des}})$$

subject to $0 \le p_j \le p_{\text{max}}, \quad j=1,\ldots,m$

with $h(u) = \max\{u, 1/u\}$



 f_0 is convex because maximum of convex functions is convex

 \mathbf{exact} solution obtained with effort pprox modest factor imes least-squares effort

additional constraints: does adding 1 or 2 below complicate the problem?

- 1. no more than half of total power is in any 10 lamps
- 2. no more than half of the lamps are on $(p_i > 0)$
- answer: with (1), still easy to solve; with (2), extremely difficult
- moral: (untrained) intuition doesn't always work; without the proper background very easy problems can appear quite similar to very difficult problems

Course goals and topics

goals

- 1. recognize/formulate problems (such as the illumination problem) as convex optimization problems
- 2. develop code for problems of moderate size (1000 lamps, 5000 patches)
- 3. characterize optimal solution (optimal power distribution), give limits of performance, etc.

topics

- 1. convex sets, functions, optimization problems
- 2. examples and applications
- 3. algorithms

Nonlinear optimization

traditional techniques for general nonconvex problems involve compromises local optimization methods (nonlinear programming)

- ullet find a point that minimizes f_0 among feasible points near it
- fast, can handle large problems
- require initial guess
- provide no information about distance to (global) optimum

global optimization methods

- find the (global) solution
- worst-case complexity grows exponentially with problem size

these algorithms are often based on solving convex subproblems

Brief history of convex optimization

theory (convex analysis): ca1900–1970

algorithms

- 1947: simplex algorithm for linear programming (Dantzig)
- 1960s: early interior-point methods (Fiacco & McCormick, Dikin, . . .)
- 1970s: ellipsoid method and other subgradient methods
- 1980s: polynomial-time interior-point methods for linear programming (Karmarkar 1984)
- late 1980s—now: polynomial-time interior-point methods for nonlinear convex optimization (Nesterov & Nemirovski 1994)

applications

- before 1990: mostly in operations research; few in engineering
- since 1990: many new applications in engineering (control, signal processing, communications, circuit design, . . .); new problem classes (semidefinite and second-order cone programming, robust optimization)