

CO 367 – Notes on Unconstrained Optimization I

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[Reference: Boyd and Vandenberghe – textbook and lecture slides]

This week we consider unconstrained minimization problems: given a C^2 function $f : \mathbb{R}^n \rightarrow \mathbb{R}$, solve

$$(*) \quad \min_{x \in \mathbb{R}^n} f(x),$$

that is, find the global minimizer(s) of f .

We assume that $p^* > -\infty$ and there exists some x^* such that $p^* = f(x^*)$.

Recall:

1. If x^* solves $(*)$, then the *first order necessary condition* holds:

$$\nabla f(x^*) = 0. \tag{1}$$

2. Suppose, in addition, that f is a convex function.

If x^* satisfies (1), then x^* solves $(*)$.

Hence

solving $(*) \equiv$ solving the equation $\nabla f(x) = 0$ for x .

1 How to solve $(*)$?

- Use *iterative methods*:
 - Start with an initial point $x^{(0)}$.
 - Produce a sequence $\{x^{(k)}\}_{k=1, \dots}$ such that
 1. $f(x^{(k)}) \leq f(x^{(0)})$ for all k .
 2. $f(x^{(k)}) \rightarrow p^*$ as $k \rightarrow \infty$.

We would hope that the sequence $\{x^{(k)}\}$ at least has a limit point, which we would take as the solution of $(*)$.

- But we usually don't know p^* a priori — then how do we know if the objective values of our iterates $x^{(k)}$ are getting close to p^* ?
- Usually we require $\{x^{(k)}\}$ to satisfy

$$\|\nabla f(x^{(k)})\| \rightarrow 0 \tag{2}$$

instead.

- **If $\{x^{(k)}\}$ has a limit point**, then by continuity of f and ∇f ,

$$\begin{aligned} f(x^{(k)}) \rightarrow p^* &\implies f(x^*) = p^*, \\ \|\nabla f(x^{(k)})\| \rightarrow 0 &\implies \nabla f(x^*) = 0. \end{aligned}$$

2 Strongly convex functions

So for what kind of function f can we be sure that, if we produce a sequence $\{x^{(k)}\}$ satisfying

1. $f(x^{(k)}) \leq f(x^{(0)})$ for all k , and
2. $\|\nabla f(x^{(k)})\| \rightarrow 0$ as $k \rightarrow \infty$,

then

1. $\{x_k\}_{k=1, \dots}$ has a limit point x^* , and
2. The limit point x^* solves (*) ?

One class of functions that enjoy this property is the class of *strongly convex functions*. A function is strongly convex on a set S if there exists some $m > 0$ such that

$$\nabla^2 f(x) - mI \text{ is positive semidefinite for all } x \in S,$$

or, equivalently,

$$f(y) \geq f(x) + \nabla f(x)^T(y - x) + \frac{m}{2} \|\nabla f(x)\|^2$$

for all $x, y \in S$.

- Let $x^{(0)} \in \text{dom } f$. If f is strongly convex on the set

$$S := \{x \in \mathbb{R}^n : f(x) \leq f(x_0)\} = f^{-1}((-\infty, f(x_0)]),$$

which is a closed set, then

1. S is a bounded set. (So S is a compact set.)

This ensures that the sequence $\{x^{(k)}\}$ has a limit point.

2. $p^* > -\infty$.

3. $f(x) - p^* \leq \frac{1}{2m} \|\nabla f(x)\|_2^2$.

This means that even if we don't know p^* , we can look at $\|\nabla f(x^{(k)})\|$ to determine how close $f(x^{(k)})$ is to the optimal value (provided we know m).

4. The limit point x^* of $\{x^{(k)}\}$ satisfies $\nabla f(x^*) = 0$ and therefore solves (*) by convexity of f .

• Simple example/non-examples of strongly convex function:

– $f(x, y) = x^2 + y^2$ is strongly convex.

– $f(x, y) = x^4 + y^4$ is not strongly convex. In fact, it is not even strictly convex (at $(0, 0)$).

– $f(x) = e^x$ is strictly convex, but not strongly convex. Note that f does not have a minimizer on \mathbb{R} : $\inf_{\mathbb{R}} f = 0$ but $f(x) > 0$ for all $x \in \mathbb{R}$.

3 Descent methods

An (iterative) *descent method* generates a sequence $\{x^{(k)}\}$ of point such that

$$f(x^{(k+1)}) < f(x^{(k)}).$$

Each iterate $x^{(k)}$ is obtained via the following: Three ingredients needed:

Algorithm 3.1: General descent method

1 Given(*initial point* $x^{(0)}$);

2 repeat

 1. Determine a *search direction* Δx .

 2. *Line search*. Choose a *step size* $t > 0$.

 3. $x \leftarrow x + t\Delta x$.

3 until *stopping criterion is satisfied*. ;

1. search direction,

2. step size, and

3. stopping criterion.

3.1 Search direction Δx

- Usually require that

$$\nabla f(x)^T \Delta x < 0. \tag{3}$$

A vector satisfying (3) is called a *descent direction*.

- Reason: f convex; if $\nabla f(x)^T \Delta x < 0$ does not hold, then for all $t > 0$

$$f(x + t\Delta x) \geq f(x) + t\nabla f(x)^T \Delta x \geq f(x),$$

so $x + t\Delta x$ does not give a lower objective value.

- **How to pick Δx ?**

A natural choice if $\Delta x = -\nabla f(x)$.

This corresponds to *steepest descent method*.

3.2 Line search

Once we fix the search direction, we need to pick a point on the ray $\{x + t\Delta x : t > 0\}$. This amounts to choosing a positive parameter t , which is called the step size. There are two common types of line search:

- *Exact line search.*

We find t such that $f(x + t\Delta x)$ is the smallest possible on the ray. To find this t is equivalent to solving the single-variate minimization problem

$$\min_{t>0} f(x + t\Delta x).$$

It can be quite expensive to perform though. So usually we use the other alternative: backtracking.

- *Backtracking line search.*

Backtracking is the computationally less expensive option for choosing a step size that guarantees sufficient objective value decrease at the new iterate.

We have two parameters $\alpha \in (0, 0.5)$ and $\beta \in (0, 1)$.

- α controls the (relative) level of objective value decrease at each iteration.
- β is the backtracking parameter.

The backtracking procedure is as follows:

- We start with $t = 1$.

– Repeat

$$t \leftarrow \beta t$$

until

$$f(x + t\Delta x) < f(x) + \alpha t \nabla f(x)^T \Delta x. \quad (4)$$

This inequality is equivalent to

$$\frac{f(x) - f(x + t\Delta x)}{t} > \alpha \cdot (-\nabla f(x)^T \Delta x),$$

so the choice of t is such that the objective value decrease (given by the LHS) is to be at least a factor α of $-\nabla f(x)^T \Delta x$.

3.3 Stopping criterion

We usually use the first order necessary condition of optimality (1): we stop at an iterate $x^{(k)}$ if $\nabla f(x^{(k)}) \sim 0$, in the sense that

$$\left\| \nabla f(x^{(k)}) \right\| \leq \varepsilon, \quad (5)$$

where $\varepsilon > 0$ is a very small number (the user-defined zero-tolerance).