

- **unconstrained problem:**  $x$  is optimal if and only if

$$x \in \mathbf{dom} f_0, \quad \nabla f_0(x) = 0$$

- **equality constrained problem**

$$\text{minimize } f_0(x) \quad \text{subject to } Ax = b$$

$x$  is optimal if and only if there exists a  $\nu$  such that

$$x \in \mathbf{dom} f_0, \quad Ax = b, \quad \nabla f_0(x) + A^T \nu = 0$$

- **minimization over nonnegative orthant**

$$\text{minimize } f_0(x) \quad \text{subject to } x \succeq 0$$

$x$  is optimal if and only if

$$x \in \mathbf{dom} f_0, \quad x \succeq 0, \quad \begin{cases} \nabla f_0(x)_i \geq 0 & x_i = 0 \\ \nabla f_0(x)_i = 0 & x_i > 0 \end{cases}$$

## Equivalent convex problems

two problems are (informally) **equivalent** if the solution of one is readily obtained from the solution of the other, and vice-versa

some common transformations that preserve convexity:

- **eliminating equality constraints**

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b \end{array}$$

is equivalent to

$$\begin{array}{ll} \text{minimize (over } z) & f_0(Fz + x_0) \\ \text{subject to} & f_i(Fz + x_0) \leq 0, \quad i = 1, \dots, m \end{array}$$

where  $F$  and  $x_0$  are such that

$$Ax = b \iff x = Fz + x_0 \text{ for some } z$$

- **introducing equality constraints**

$$\begin{array}{ll} \text{minimize} & f_0(A_0x + b_0) \\ \text{subject to} & f_i(A_ix + b_i) \leq 0, \quad i = 1, \dots, m \end{array}$$

is equivalent to

$$\begin{array}{ll} \text{minimize (over } x, y_i) & f_0(y_0) \\ \text{subject to} & f_i(y_i) \leq 0, \quad i = 1, \dots, m \\ & y_i = A_ix + b_i, \quad i = 0, 1, \dots, m \end{array}$$

- **introducing slack variables for linear inequalities**

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & a_i^T x \leq b_i, \quad i = 1, \dots, m \end{array}$$

is equivalent to

$$\begin{array}{ll} \text{minimize (over } x, s) & f_0(x) \\ \text{subject to} & a_i^T x + s_i = b_i, \quad i = 1, \dots, m \\ & s_i \geq 0, \quad i = 1, \dots, m \end{array}$$

- **epigraph form:** standard form convex problem is equivalent to

$$\begin{array}{ll}
 \text{minimize (over } x, t) & t \\
 \text{subject to} & f_0(x) - t \leq 0 \\
 & f_i(x) \leq 0, \quad i = 1, \dots, m \\
 & Ax = b
 \end{array}$$

- **minimizing over some variables**

$$\begin{array}{ll}
 \text{minimize} & f_0(x_1, x_2) \\
 \text{subject to} & f_i(x_1) \leq 0, \quad i = 1, \dots, m
 \end{array}$$

is equivalent to

$$\begin{array}{ll}
 \text{minimize} & \tilde{f}_0(x_1) \\
 \text{subject to} & f_i(x_1) \leq 0, \quad i = 1, \dots, m
 \end{array}$$

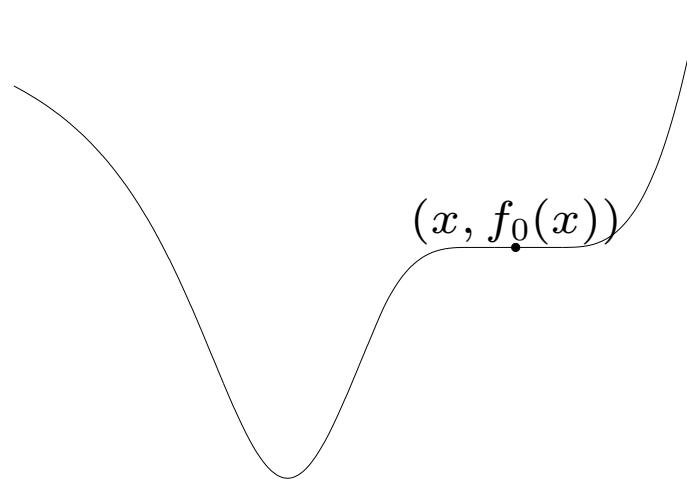
where  $\tilde{f}_0(x_1) = \inf_{x_2} f_0(x_1, x_2)$

# Quasiconvex optimization

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & Ax = b \end{array}$$

with  $f_0 : \mathbf{R}^n \rightarrow \mathbf{R}$  quasiconvex,  $f_1, \dots, f_m$  convex

can have locally optimal points that are not (globally) optimal



## convex representation of sublevel sets of $f_0$

if  $f_0$  is quasiconvex, there exists a family of functions  $\phi_t$  such that:

- $\phi_t(x)$  is convex in  $x$  for fixed  $t$
- $t$ -sublevel set of  $f_0$  is 0-sublevel set of  $\phi_t$ , *i.e.*,

$$f_0(x) \leq t \iff \phi_t(x) \leq 0$$

### example

$$f_0(x) = \frac{p(x)}{q(x)}$$

with  $p$  convex,  $q$  concave, and  $p(x) \geq 0$ ,  $q(x) > 0$  on  $\text{dom } f_0$

can take  $\phi_t(x) = p(x) - tq(x)$ :

- for  $t \geq 0$ ,  $\phi_t$  convex in  $x$
- $p(x)/q(x) \leq t$  if and only if  $\phi_t(x) \leq 0$

## quasiconvex optimization via convex feasibility problems

$$\phi_t(x) \leq 0, \quad f_i(x) \leq 0, \quad i = 1, \dots, m, \quad Ax = b \quad (1)$$

- for fixed  $t$ , a convex feasibility problem in  $x$
- if feasible, we can conclude that  $t \geq p^*$ ; if infeasible,  $t \leq p^*$

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*Bisection method for quasiconvex optimization*

**given**  $l \leq p^*$ ,  $u \geq p^*$ , tolerance  $\epsilon > 0$ .

**repeat**

1.  $t := (l + u)/2$ .
2. Solve the convex feasibility problem (1).
3. **if** (1) is feasible,  $u := t$ ; **else**  $l := t$ .

**until**  $u - l \leq \epsilon$ .

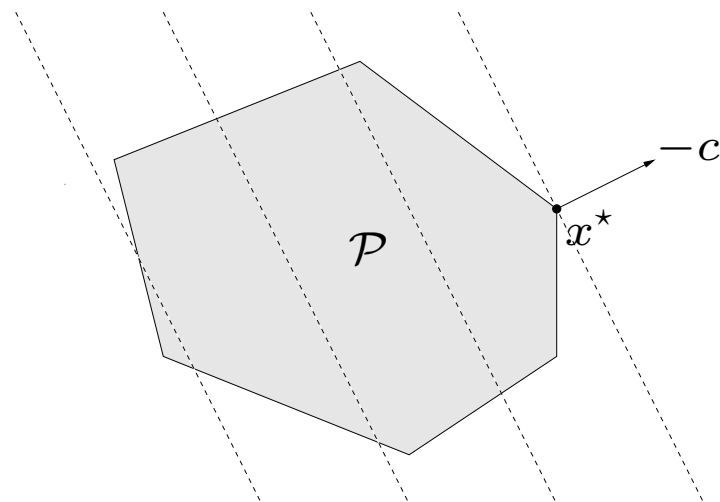
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requires exactly  $\lceil \log_2((u - l)/\epsilon) \rceil$  iterations (where  $u, l$  are initial values)

# Linear program (LP)

$$\begin{array}{ll} \text{minimize} & c^T x + d \\ \text{subject to} & Gx \preceq h \\ & Ax = b \end{array}$$

- convex problem with affine objective and constraint functions
- feasible set is a polyhedron





# Examples

**diet problem:** choose quantities  $x_1, \dots, x_n$  of  $n$  foods

- one unit of food  $j$  costs  $c_j$ , contains amount  $a_{ij}$  of nutrient  $i$
- healthy diet requires nutrient  $i$  in quantity at least  $b_i$

to find cheapest healthy diet,

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \succeq b, \quad x \succeq 0 \end{array}$$

## piecewise-linear minimization

$$\text{minimize} \quad \max_{i=1, \dots, m} (a_i^T x + b_i)$$

equivalent to an LP

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & a_i^T x + b_i \leq t, \quad i = 1, \dots, m \end{array}$$

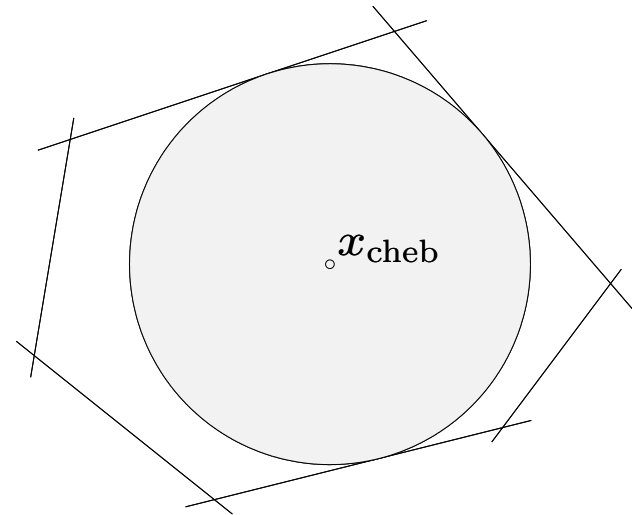
## Chebyshev center of a polyhedron

Chebyshev center of

$$\mathcal{P} = \{x \mid a_i^T x \leq b_i, \quad i = 1, \dots, m\}$$

is center of largest inscribed ball

$$\mathcal{B} = \{x_c + u \mid \|u\|_2 \leq r\}$$



- $a_i^T x \leq b_i$  for all  $x \in \mathcal{B}$  if and only if

$$\sup\{a_i^T (x_c + u) \mid \|u\|_2 \leq r\} = a_i^T x_c + r\|a_i\|_2 \leq b_i$$

- hence,  $x_c, r$  can be determined by solving the LP

$$\begin{array}{ll} \text{maximize} & r \\ \text{subject to} & a_i^T x_c + r\|a_i\|_2 \leq b_i, \quad i = 1, \dots, m \end{array}$$

## (Generalized) linear-fractional program

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & Gx \preceq h \\ & Ax = b \end{array}$$

### linear-fractional program

$$f_0(x) = \frac{c^T x + d}{e^T x + f}, \quad \text{dom } f_0(x) = \{x \mid e^T x + f > 0\}$$

- a quasiconvex optimization problem; can be solved by bisection
- also equivalent to the LP (variables  $y, z$ )

$$\begin{array}{ll} \text{minimize} & c^T y + dz \\ \text{subject to} & Gy \preceq hz \\ & Ay = bz \\ & e^T y + fz = 1 \\ & z \geq 0 \end{array}$$

## generalized linear-fractional program

$$f_0(x) = \max_{i=1,\dots,r} \frac{c_i^T x + d_i}{e_i^T x + f_i}, \quad \text{dom } f_0(x) = \{x \mid e_i^T x + f_i > 0, i = 1, \dots, r\}$$

a quasiconvex optimization problem; can be solved by bisection

**example:** Von Neumann model of a growing economy

$$\begin{array}{ll} \text{maximize (over } x, x^+) & \min_{i=1,\dots,n} x_i^+ / x_i \\ \text{subject to} & x^+ \succeq 0, \quad Bx^+ \preceq Ax \end{array}$$

- $x, x^+ \in \mathbf{R}^n$ : activity levels of  $n$  sectors, in current and next period
- $(Ax)_i, (Bx^+)_i$ : produced, resp. consumed, amounts of good  $i$
- $x_i^+ / x_i$ : growth rate of sector  $i$

allocate activity to maximize growth rate of slowest growing sector