

Smoothed Aggregation Multigrid for Slowly Mixing Markov Chains

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1. Simple Markov Chain Example

- start in one state with probability 1: what is the stationary probability vector after ∞ number of steps?

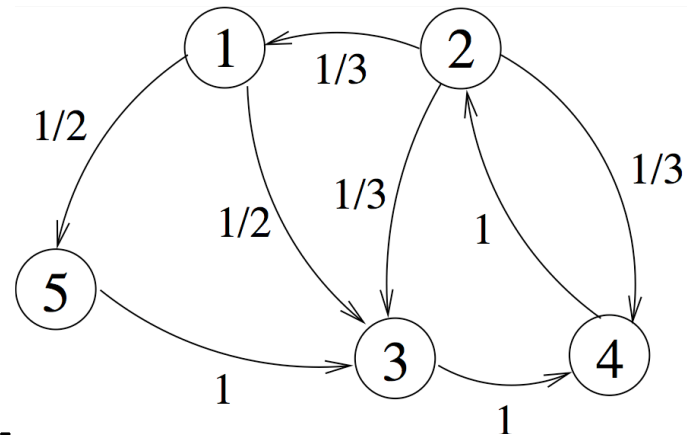
$$B = \begin{bmatrix} 0 & 1/3 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 1/2 & 1/3 & 0 & 0 & 1 \\ 0 & 1/3 & 1 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\mathbf{x}_{i+1} = B \mathbf{x}_i$$

- stationary probability:

$$B \mathbf{x} = \mathbf{x} \quad \|\mathbf{x}\|_1 = 1$$

$$\mathbf{x}^T = [2/19 \ 6/19 \ 4/19 \ 6/19 \ 1/19]$$



2. Problem Statement

$$B \mathbf{x} = \mathbf{x} \quad \|\mathbf{x}\|_1 = 1 \quad x_i \geq 0 \forall i$$

- B is column-stochastic

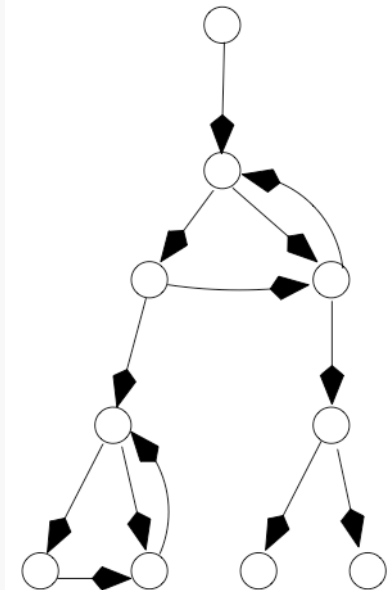
$$0 \leq b_{ij} \leq 1 \quad \forall i, j \quad \mathbf{1}^T B = \mathbf{1}^T$$

- B is irreducible (every state can be reached from every other state in the directed graph)

\Rightarrow

$$\exists! \mathbf{x} : \quad B \mathbf{x} = \mathbf{x} \quad \|\mathbf{x}\|_1 = 1 \quad x_i > 0 \quad \forall i$$

(no probability sinks!)



3. Power Method

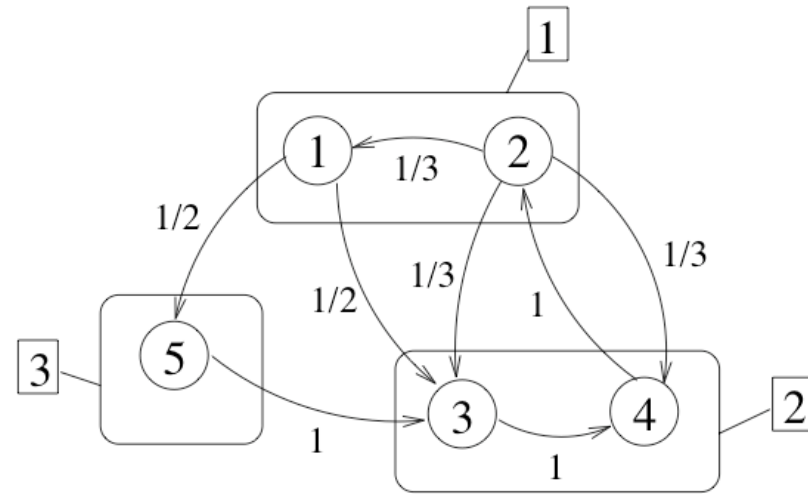
$$B \mathbf{x} = \mathbf{x} \quad \text{or} \quad (I - B) \mathbf{x} = 0 \quad \text{or} \quad A \mathbf{x} = 0$$

- largest eigenvalue of B : $\lambda_1 = 1$
- power method: $\mathbf{x}_{i+1} = B \mathbf{x}_i$
 - convergence factor: $|\lambda_2|$
 - convergence is very slow when
$$|\lambda_2| \approx 1$$
(slowly mixing Markov chain) (JAC, GS also slow)

4. Aggregation for Markov Chains

$$B_c \mathbf{x}_c = \mathbf{x}_c$$

$$b_{c,IJ} = \frac{\sum_{j \in J} x_j \left(\sum_{i \in I} b_{ij} \right)}{\sum_{j \in J} x_j}$$



$$B_c = Q^T B \text{diag}(\mathbf{x}) Q \text{diag}(Q^T \mathbf{x})^{-1}$$

$$x_{c,I} = \sum_{i \in I} x_i$$

$$\mathbf{x}_c = Q^T \mathbf{x}$$

$$Q = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

(Krieger, Horton, ... 1990s)

5. Error Equation

- error equation - coarse grid correction:

$$\mathbf{x} = \text{diag}(\mathbf{x}_i) \mathbf{e}_i$$

$$A \text{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$$

$$Q^T A \text{diag}(\mathbf{x}_i) Q \mathbf{e}_c = 0$$

$$A_c \mathbf{e}_c = 0$$

$$R = Q^T \quad P = \text{diag}(\mathbf{x}_i) Q$$

$$A_c = R A P$$

$$\mathbf{x}_{i+1} = P \mathbf{e}_c$$

Error Equation

- important properties of A_c :

$$\mathbf{x} = \text{diag}(\mathbf{x}_i) \mathbf{e}_i$$

$$A \text{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$$

$$Q^T A \text{diag}(\mathbf{x}_i) Q \mathbf{e}_c = 0$$

$$A_c \mathbf{e}_c = 0$$

$$R = Q^T \quad P = \text{diag}(\mathbf{x}_i) Q$$

$$A_c = R A P$$

$$(1) \mathbf{1}_c^T A_c = 0 \quad \forall \mathbf{x}_i$$

$$(\text{since } \mathbf{1}_c^T R = \mathbf{1}^T \text{ and } \mathbf{1}^T A = 0)$$

$$(2) A_c \mathbf{1}_c = 0 \quad \text{for } \mathbf{x}_i = \mathbf{x}$$

$$A_c (\text{diag}(P^T \mathbf{1}))^{-1}$$

$$= R (I - B) P (\text{diag}(P^T \mathbf{1}))^{-1}$$

$$= I_c - B_c$$

6. Multilevel Aggregation Algorithm

Algorithm: Multilevel Adaptive Aggregation
method (V-cycle)

$\mathbf{x} = \text{AM_V}(A, \mathbf{x}, \nu_1, \nu_2)$

begin

$\mathbf{x} \leftarrow \text{Relax}(A, \mathbf{x}) \quad \nu_1 \text{ times}$

build Q based on \mathbf{x} and A (Q is rebuilt every level and cycle)

$R = Q^T$ and $P = \text{diag}(\mathbf{x}) Q$

$A_c = R A P$

$\mathbf{x}_c = \text{AM_V}(A_c \text{diag}(P^T \mathbf{1})^{-1}, P^T \mathbf{1}, \nu_1, \nu_2)$ (coarse-level solve)

$\mathbf{x} = P (\text{diag}(P^T \mathbf{1}))^{-1} \mathbf{x}_c$ (coarse-level correction)

$\mathbf{x} \leftarrow \text{Relax}(A, \mathbf{x}) \quad \nu_2 \text{ times}$

end

(Krieger, Horton 1994, but no good
way to build Q , convergence not good)

7. Well-posedness: Singular M-matrices

- singular M-matrix:

$A \in \mathbb{R}^{n \times n}$ is a singular M-matrix \Leftrightarrow

$\exists B \in \mathbb{R}^{n \times n}, b_{ij} \geq 0 \forall i, j : A = \rho(B)I - B$

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

- our $A=I-B$ is a singular M-matrix on all levels

(1) Irreducible singular M-matrices have a unique solution to the problem $A\mathbf{x} = 0$, up to scaling. All components of \mathbf{x} have strictly the same sign (i.e., scaling can be chosen s.t. $x_i > 0 \forall i$). (This follows directly from the Perron-Frobenius theorem.)

(3) Irreducible singular M-matrices have nonpositive off-diagonal elements, and strictly positive diagonal elements ($n > 1$).

(4) If A has a strictly positive element in its left or right nullspace and the off-diagonal elements of A are nonpositive, then A is a singular M-matrix (see also [21]).

Well-posedness: Unsmoothed Method

THEOREM 3.1 (Singular M-matrix property of AM coarse-level operators). A_c is an irreducible singular M-matrix on all coarse levels, and thus has a unique right kernel vector \mathbf{e}_c with strictly positive components (up to scaling) on all levels.

$$(1) \quad \mathbf{1}_c^T A_c = 0 \quad \forall \mathbf{x}_i$$

(since $\mathbf{1}_c^T R = \mathbf{1}^T$ and $\mathbf{1}^T A = 0$)

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

THEOREM 3.2 (Fixed-point property of AM V-cycle). *Exact solution \mathbf{x} is a fixed point of the AM V-cycle.*

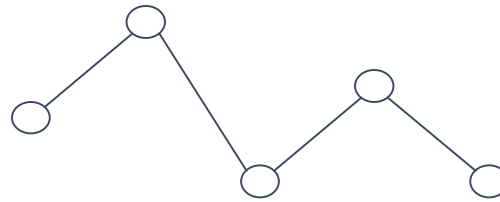
$$(2) \quad A_c \mathbf{1}_c = 0 \quad \text{for } \mathbf{x}_i = \mathbf{x}$$

$$A_c \mathbf{e}_c = 0$$

$$\mathbf{x}_{i+1} = P \mathbf{e}_c$$

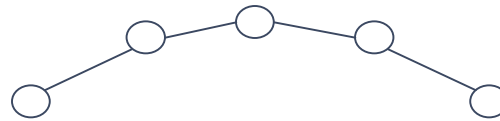
8. We Need 'Smoothed Aggregation' ...

(Vanek, Mandel, and Brezina, Computing, 1996)

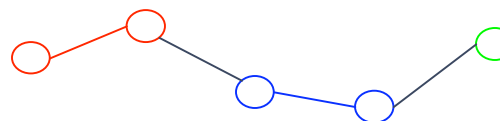


$$Q = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

after smoothing:

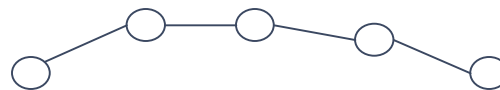


coarse grid
correction with Q :



$$Q_s = \begin{bmatrix} \times & 0 & 0 \\ \times & \times & 0 \\ \times & \times & 0 \\ 0 & \times & \times \\ 0 & \times & \times \end{bmatrix}$$

coarse grid
correction with Q_s :



Smoothed Aggregation

$$A = D - (L + U)$$

- smooth the columns of P with weighted Jacobi:

$$P_s = (1 - w) \text{diag}(\mathbf{x}_i) Q + w D^{-1} (L + U) \text{diag}(\mathbf{x}_i) Q$$

- smooth the rows of R with weighted Jacobi:

$$R_s = R (1 - w) + R w (L + U) D^{-1}$$

Smoothed Aggregation

- smoothed coarse level operator:

$$\begin{aligned} A_{cs} &= R_s (D - (L + U)) P_s \\ &= R_s D P_s - R_s (L + U) P_s \end{aligned}$$

$$\begin{aligned} \mathbf{1}_c^T A_{cs} &= 0 \quad \forall \mathbf{x}_i, \\ A_{cs} \mathbf{1}_c &= 0 \quad \text{for } \mathbf{x}_i = \mathbf{x} \end{aligned}$$

- problem: A_{cs} is not a singular M-matrix (signs wrong)

- solution: lumping approach on S in

$$A_{cs} = S - G$$

$$\hat{A}_{cs} = \hat{S} - G$$

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

Smoothed Aggregation

$$A_{cs} = S - G$$

$$\hat{A}_{cs} = \hat{S} - G$$

- we want as little lumping as possible
- only lump 'offending' elements (i,j) :

$$s_{ij} \neq 0, i \neq j \text{ and } s_{ij} - g_{ij} \geq 0$$

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

$$\mathbf{1}_c^T \hat{A}_{cs} = 0 \quad \forall \mathbf{x}_i,$$

$$\hat{A}_{cs} \mathbf{1}_c = 0 \quad \text{for } \mathbf{x}_i = \mathbf{x}$$

(we consider both off-diagonal signs and reducibility here!)

- for 'offending' elements (i,j) , add $S_{\{i,j\}}$ to S :

$$S_{\{i,j\}} = \begin{bmatrix} & & i & & j & & \\ & \ddots & \vdots & & \vdots & & \\ i & \dots & \beta_{\{i,j\}} & \dots & -\beta_{\{i,j\}} & \dots & \\ & & \vdots & & \vdots & & \\ j & \dots & -\beta_{\{i,j\}} & \dots & \beta_{\{i,j\}} & \dots & \\ & & \vdots & & \vdots & & \end{bmatrix}$$

$$s_{ij} - g_{ij} - \beta_{\{i,j\}} < 0$$

$$s_{ji} - g_{ji} - \beta_{\{i,j\}} < 0$$

conserves both row and column sums

9. Lumped Smoothed Method is Well-posed

THEOREM 4.1 (Singular M-matrix property of lumped SAM coarse-level operators). \hat{A}_{cs} is an irreducible singular M-matrix on all coarse levels, and thus has a unique right kernel vector \mathbf{e}_c with strictly positive components (up to scaling) on all levels.

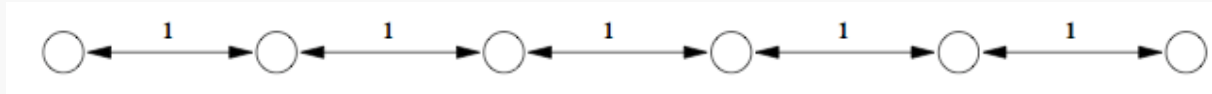
THEOREM 4.2 (Fixed-point property of lumped SAM V-cycle). *Exact solution \mathbf{x} is a fixed point of the SAM V-cycle (with lumping).*

$$\begin{aligned}\mathbf{1}_c^T \hat{A}_{cs} &= 0 \quad \forall \mathbf{x}_i, \\ \hat{A}_{cs} \mathbf{1}_c &= 0 \quad \text{for } \mathbf{x}_i = \mathbf{x}\end{aligned}$$

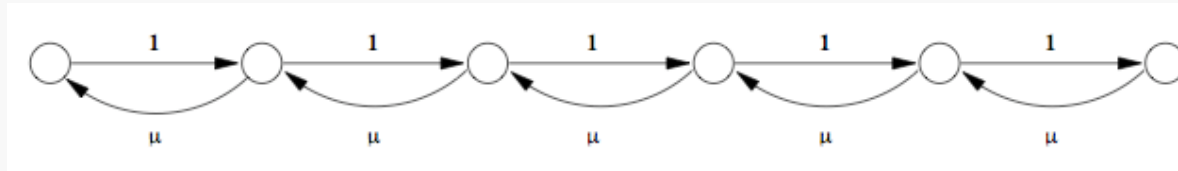
$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

10. Numerical Results: Test Problems

- uniform chain

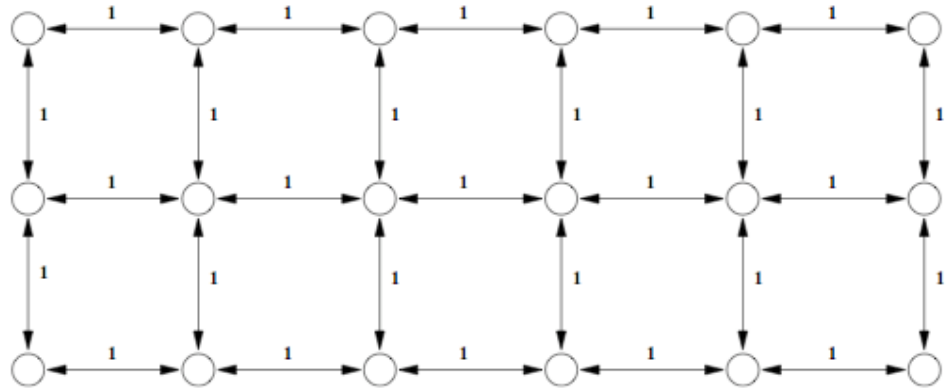


- birth-death chain

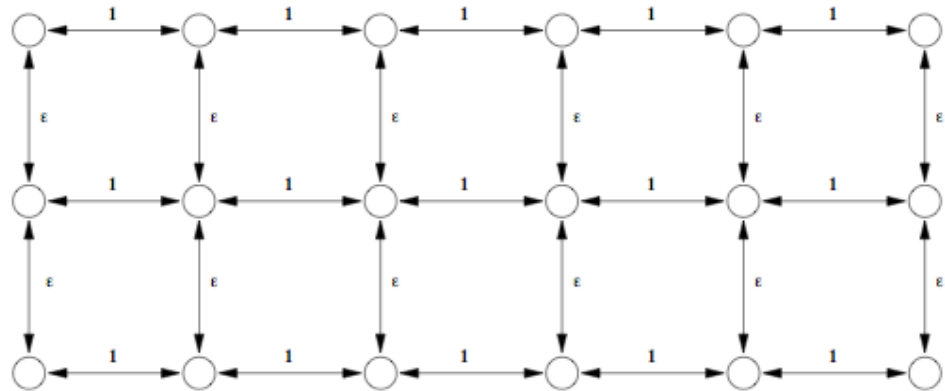


Test Problems

- uniform 2D lattice



- anisotropic 2D lattice



Test Problems

- tandem queueing network



FIG. 5.6. Tandem queueing network.

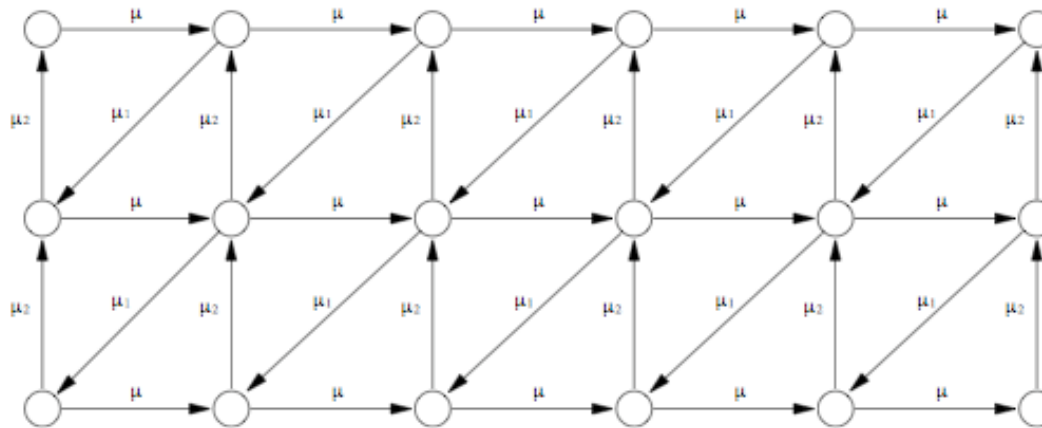
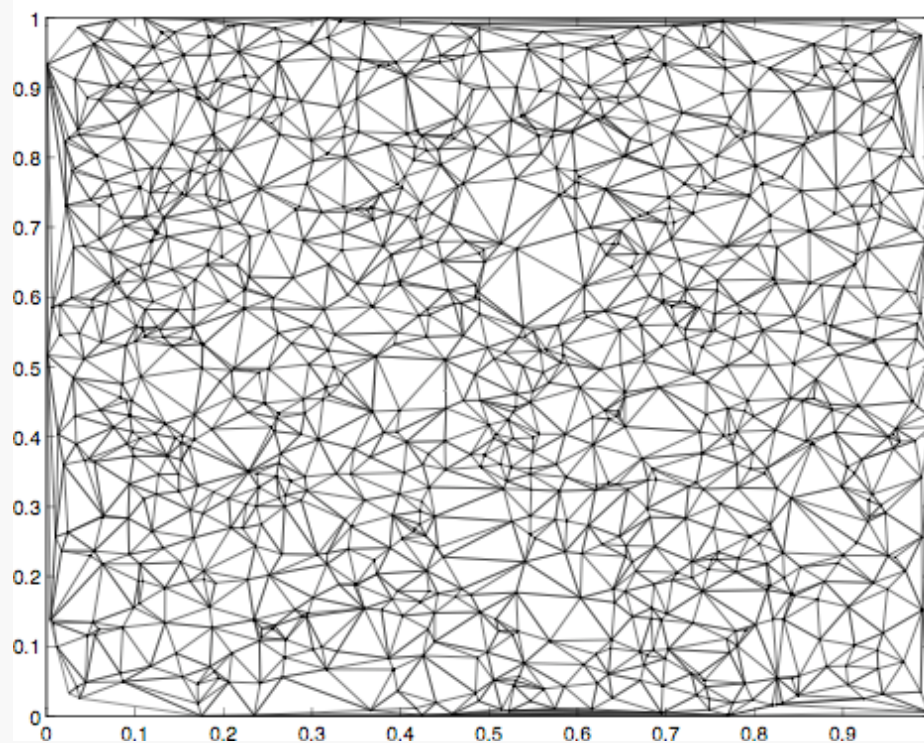


FIG. 5.7. Graph for tandem queueing network.

Test Problems

- random walk on triangulation of random points in unit square



11. Numerical Results: Geometric Aggregation (size 3)

n	lev	it	C_{op}	γ	γ_{eff}	it	C_{op}	γ	γ_{eff}
27	2	32	1.32	0.66	0.73	32	1.63	0.66	0.78
81	3	85	1.43	0.87	0.91	52	2.07	0.80	0.90
243	4	>100	1.47	0.95	0.97	73	2.37	0.87	0.94
729	5	>100	1.49	0.98	0.98	>100	2.58	0.92	0.97
2187	6	>100	1.50	0.98	0.98	>100	2.72	0.95	0.98

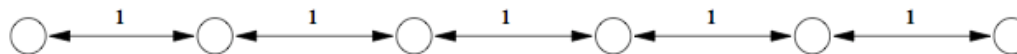
TABLE 5.1

Uniform chain. G-AM with V-cycles (left) and W-cycles (right). (Size-three aggregates, no smoothing.)

n	it	C_{op}	lev	γ	γ_{eff}	R_{lump}
27	13	1.32	2	0.27	0.37	0
81	13	1.43	3	0.27	0.40	0
243	13	1.47	4	0.27	0.41	0
729	13	1.49	5	0.27	0.41	0
2187	13	1.50	6	0.27	0.42	0
6561	13	1.50	7	0.27	0.42	0
19683	13	1.50	8	0.27	0.42	0
59049	13	1.50	9	0.27	0.42	0

TABLE 5.2

Uniform chain. G-SAM with V-cycles and size-three aggregates. (Smoothing with lumping.)

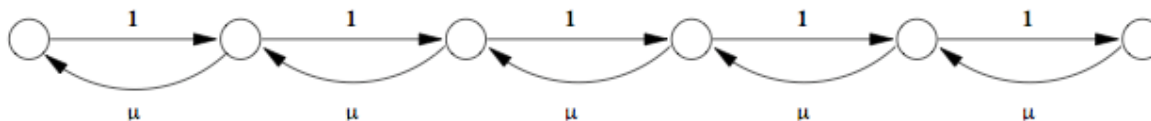


Numerical Results: Geometric Aggregation (size 3)

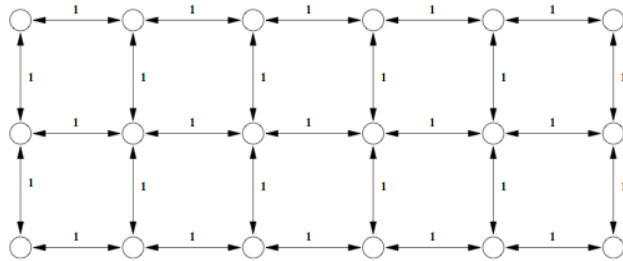
n	C_{op}	lev	it	γ	γ_{eff}	it	γ	γ_{eff}	R_{lump}
27	1.32	2	33	0.66	0.73	13	0.27	0.37	0
81	1.43	3	95	0.88	0.92	12	0.27	0.40	0
243	1.47	4	>100	0.95	0.97	13	0.26	0.40	0
729	1.49	5	>100	0.97	0.98	12	0.24	0.38	0

TABLE 5.4

Birth-death chain ($\mu = 0.96$). (left) G-AM (no smoothing) and (right) G-SAM (smoothing with lumping). (V-cycles and size-three aggregates.)



Numerical Results: Geometric Aggregation (3x3)



n	γ_{res}	iter	C_{op}	levels
64	0.71	43	1.11	2
100	0.85	72	1.17	3
169	0.86	85	1.15	3
400	0.89	>100	1.13	3
900	0.95	>100	1.12	4

TABLE 5.9

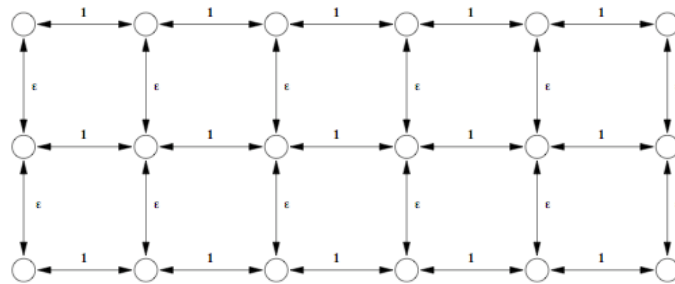
Uniform 2D lattice. G-AM with V-cycles and size-three aggregates. (No smoothing.)

n	lev	it	C_{op}	γ	γ_{eff}	R_{lump}	it	C_{op}	γ	γ_{eff}	R_{lump}
64	2	22	1.17	0.49	0.55	0	17	1.34	0.39	0.49	0
256	3	24	1.22	0.54	0.61	0	17	1.47	0.42	0.55	0
1024	4	25	1.22	0.57	0.63	6.6e-4	18	1.49	0.44	0.58	4.3e-3
4096	4	26	1.23	0.58	0.64	9.7e-4	19	1.52	0.45	0.59	3.6e-3
16384	5	27	1.22	0.61	0.67	4.0e-5	19	1.51	0.45	0.59	1.3e-4
65536	6	27	1.23	0.62	0.67	1.5e-5	19	1.52	0.45	0.59	2.6e-4

TABLE 5.8

Uniform 2D lattice. G-SAM with V-cycles (left) and W-cycles (right), using three-by-three aggregates. (Smoothing with lumping.)

Numerical Results: Geometric Aggregation (3x3)

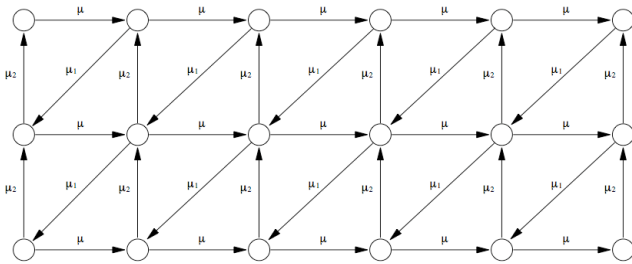


n	lev	it	C_{op}	γ	γ_{eff}	it	C_{op}	γ	γ_{eff}	R_{lump}
64	2	>100	1.11	1.00	1.00	>100	1.17	1.00	1.00	0
256	3	>100	1.14	0.97	0.97	>100	1.22	0.95	0.96	0
1024	4	>100	1.13	0.99	0.99	>100	1.22	0.99	0.99	0
4096	4	>100	1.13	0.99	0.99	>100	1.23	0.99	0.99	0

TABLE 5.9

Anisotropic 2D lattice ($\epsilon = 1e-6$). (left) G-AM (no smoothing) and (right) G-SAM (smoothing with lumping) with V-cycles and three-by-three aggregates.

Numerical Results: Geometric Aggregation (3x3)



n	it	C_{op}	lev	γ	γ_{eff}
256	>100	1.19	3	0.91	0.93
1024	>100	1.19	4	0.96	0.97
4096	>100	1.19	4	0.96	0.97

TABLE 5.10

Tandem queueing network. G-AM with V-cycles and three-by-three aggregates. (No smoothing.)

n	lev	it	C_{op}	γ	γ_{eff}	R_{lump}	it	C_{op}	γ	γ_{eff}	R_{lump}
256	3	19	1.28	0.44	0.52	2.4e-2	14	1.60	0.33	0.50	3.9e-2
1024	4	19	1.27	0.46	0.55	2.0e-3	14	1.62	0.32	0.50	5.9e-3
4096	4	20	1.29	0.49	0.57	7.4e-3	14	1.65	0.33	0.51	1.4e-2
16384	5	21	1.28	0.56	0.63	8.1e-4	14	1.63	0.33	0.51	2.0e-3
65536	6	21	1.28	0.51	0.59	1.6e-3	14	1.64	0.33	0.51	3.8e-3

TABLE 5.11

Tandem queueing network. G-SAM with V-cycles (left) and W-cycles (right), using three-by-three aggregates. (Smoothing with lumping.)

12. Numerical Results: Algebraic Aggregation

- error equation: $A \text{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$
 - use strength of connection in $A \text{diag}(\mathbf{x}_i)$
 - define row-based strength (determine all states that strongly influence a row's state, similar to AMG)
 - state that has largest value in \mathbf{x}_i is seed point for new aggregate, and all unassigned states influenced by it join its aggregate
 - repeat
- (our paper on PageRank in SISC, 2008)

Numerical Results: Algebraic Aggregation

n	it	C_{op}	lev	γ	γ_{eff}
27	39	1.71	3	0.74	0.84
81	83	1.85	4	0.87	0.93
243	>100	1.96	6	0.96	0.98
729	>100	1.98	7	1.00	1.00

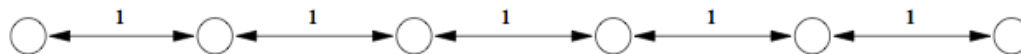
TABLE 6.1

Uniform chain. A-AM with V-cycles and distance-one aggregation. (No smoothing.)

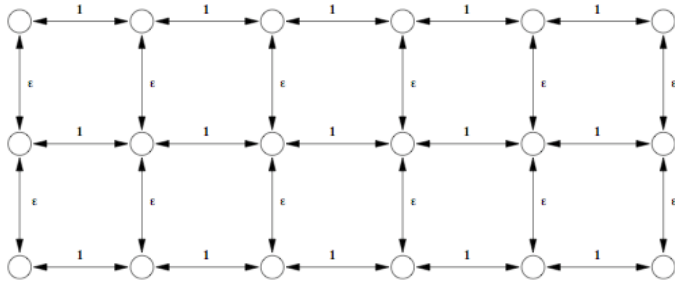
n	it	C_{op}	lev	γ	γ_{eff}	R_{lump}	it	C_{op}	lev	γ	γ_{eff}	R_{lump}
27	10	2.09	3	0.20	0.46	2.4e-2	13	1.33	2	0.32	0.43	0
81	10	2.09	4	0.18	0.44	0.0e-0	13	1.44	3	0.32	0.45	0
243	10	2.20	5	0.21	0.49	1.3e-3	12	1.46	4	0.32	0.46	0
729	11	2.23	6	0.21	0.50	3.3e-3	12	1.49	5	0.31	0.45	0
2187	11	2.24	6	0.23	0.52	2.5e-3	12	1.49	6	0.31	0.46	0
6561	11	2.24	8	0.24	0.53	1.4e-3	12	1.49	7	0.31	0.46	0
19683	11	2.25	8	0.26	0.55	1.4e-3	12	1.49	8	0.32	0.47	0
59049	12	2.25	9	0.30	0.59	1.0e-3	12	1.50	9	0.32	0.47	0

TABLE 6.2

Uniform chain. A-SAM with V-cycles using distance-one aggregation (left) and distance-two aggregation (right). (Smoothing with lumping.)



Numerical Results: Algebraic Aggregation

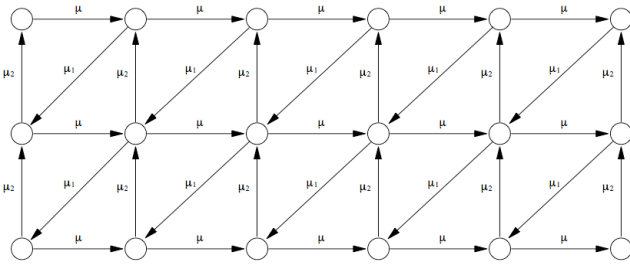


n	it	C_{op}	lev	γ	γ_{eff}	it	C_{op}	lev	γ	γ_{eff}	R_{lump}
64	27	1.78	4	0.55	0.72	17	1.76	3	0.40	0.59	0.0e-0
256	45	1.93	6	0.74	0.86	15	2.23	4	0.33	0.61	7.4e-4
1024	80	2.01	8	0.86	0.93	14	2.81	5	0.33	0.68	1.6e-3
4096	>100	2.04	10	0.93	0.96	14	3.43	7	0.33	0.73	4.9e-4
16384	>100	2.06	11	0.96	0.98	13	4.17	7	0.33	0.77	2.5e-4
65536	>100	2.16	13	0.97	0.99	13	4.80	9	0.32	0.79	7.6e-5

TABLE 6.7

Anisotropic 2D lattice ($\epsilon = 1e - 6$). (left) A-AM (no smoothing, distance-one aggregation) and (right) A-SAM (smoothing with lumping, distance-two aggregation) with V-cycles.

Numerical Results: Algebraic Aggregation



n	it	C_{op}	lev	γ	γ_{eff}
256	>100	1.86	5	0.92	0.96
1024	>100	1.96	6	0.91	0.95
4096	>100	2.03	7	0.98	0.99
16384	>100	2.09	9	0.98	0.99

TABLE 6.8

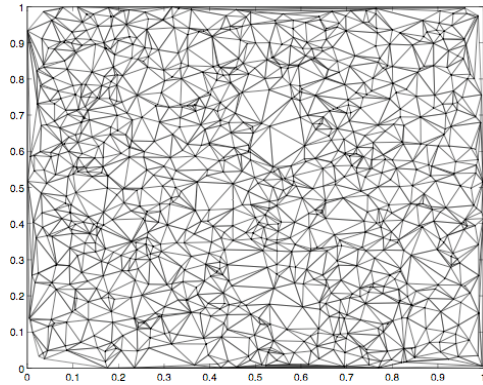
Tandem queueing network. A-AM with V-cycles and distance-one aggregation. (No smoothing.)

n	it	C_{op}	lev	γ	γ_{eff}	R_{lump}
256	18	1.94	4	0.39	0.61	1.1e-1
1024	20	2.04	4	0.41	0.64	7.6e-2
4096	24	2.12	5	0.45	0.69	5.5e-2
16384	30	2.18	6	0.56	0.77	5.3e-2
65536	37	2.37	6	0.71	0.86	1.3e-1

TABLE 6.9

Tandem queueing network. A-SAM with V-cycles and distance-two aggregation. (Smoothing with lumping.)

Numerical Results: Algebraic Aggregation



$n = 32768$	it	C_{op}	lev	γ	γ_{eff}
multilevel A-AM	>100	1.26	8	0.9827	0.9862
two-level A-AM	>100	1.20	2	0.9833	0.9861
one-level method	>100	1.00	1	0.9846	0.9846

TABLE 6.11

Unstructured planar graph with $n = 32768$ nodes. Comparison of multilevel A-AM, two-level A-AM (aggregation-disaggregation), and one-level relaxation. (V-cycles and distance-one aggregation, no smoothing. The $V(1,1)$ -cycles have two relaxations per level per cycle.)

n	it	C_{op}	lev	γ	γ_{eff}	R_{lump}
1024	20	1.69	5	0.5265	0.6848	2.6e-02
2048	19	1.68	5	0.5200	0.6779	2.1e-02
4096	21	1.80	5	0.6069	0.7578	2.4e-02
8192	22	1.92	7	0.6360	0.7904	2.5e-02
16384	30	2.03	7	0.7610	0.8744	2.4e-02
32768	28	2.08	7	0.7399	0.8649	2.4e-02

TABLE 6.10

Unstructured planar graph. A-SAM with V-cycles and distance-one aggregation. (Smoothing with lumping.)

13. AMG version of algorithm also works

$$\mathbf{v} \leftarrow V(A, \mathbf{x})$$

1. Relax a_1 times on $A\mathbf{u} = \mathbf{0}$ with initial guess \mathbf{x} .
2. If you have reached the coarsest level, then go to step 4.

Else

Set $\bar{A} \leftarrow A \text{diag}(\mathbf{x})$.

Compute the set of coarse grid points C .

Construct the interpolation operator P .

Construct the coarse grid operator $\bar{A}_c \leftarrow P^T \bar{A} P$.

Obtain lumped coarse grid operator $\hat{A}_c \leftarrow \text{Lump}(\bar{A}_c, \eta)$.

Recursively call the algorithm $\mathbf{e}_c \leftarrow V(\hat{A}_c, \mathbf{1}_c)$.

3. Correct $\mathbf{x} \leftarrow \text{diag}(\mathbf{x}) P \mathbf{e}_c$.
4. Relax a_2 times on $A\mathbf{u} = \mathbf{0}$ with initial guess \mathbf{x} .

14. Conclusions

- SAM (Smoothed Aggregation for Markov Chains): algorithm for stationary vector of slowly mixing Markov chains with near-optimal complexity
- smoothing is essential for aggregation for many problems
- pretty good convergence results, operator complexity may grow
- AMG works too
- appropriate theoretical framework (well-posedness)
- are there other ways for choosing R , P , lumping?
- no theory yet on optimal convergence (non-symmetric matrices)

- Questions?

Smoothed Aggregation

$$A_{cs} = S - G$$

$$\hat{A}_{cs} = \hat{S} - G$$

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

$$s_{ij} - g_{ij} - \beta_{\{i,j\}} < 0$$

$$s_{ji} - g_{ji} - \beta_{\{i,j\}} < 0$$

- for 'offending' elements (i,j) , choose $\eta \in (0,1]$ s.t.

$$s_{ij} - g_{ij} - \beta_{\{i,j\}}^{(1)} = -\eta g_{ij}$$

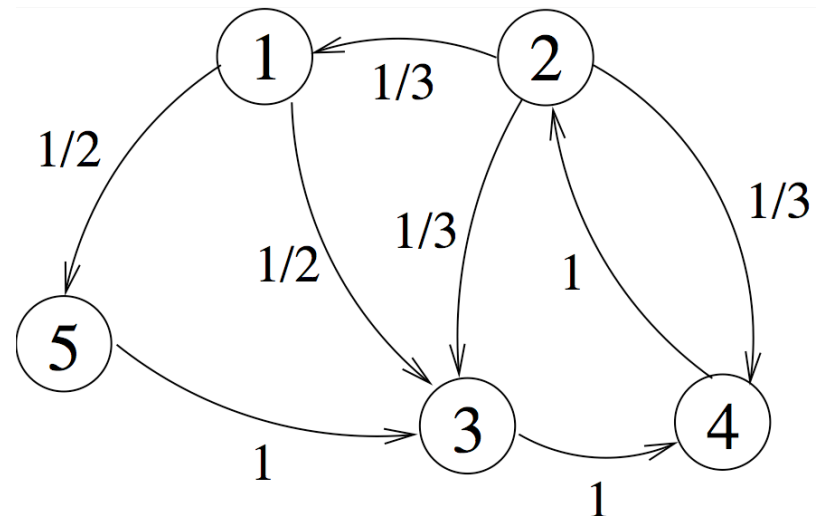
$$s_{ji} - g_{ji} - \beta_{\{i,j\}}^{(2)} = -\eta g_{ji}$$

$$\text{with } \beta_{\{i,j\}} = \max(\beta_{\{i,j\}}^{(1)}, \beta_{\{i,j\}}^{(2)})$$

- $\eta=1$ means lump full value of offending elements of S ($\hat{s}_{ij} = 0$)

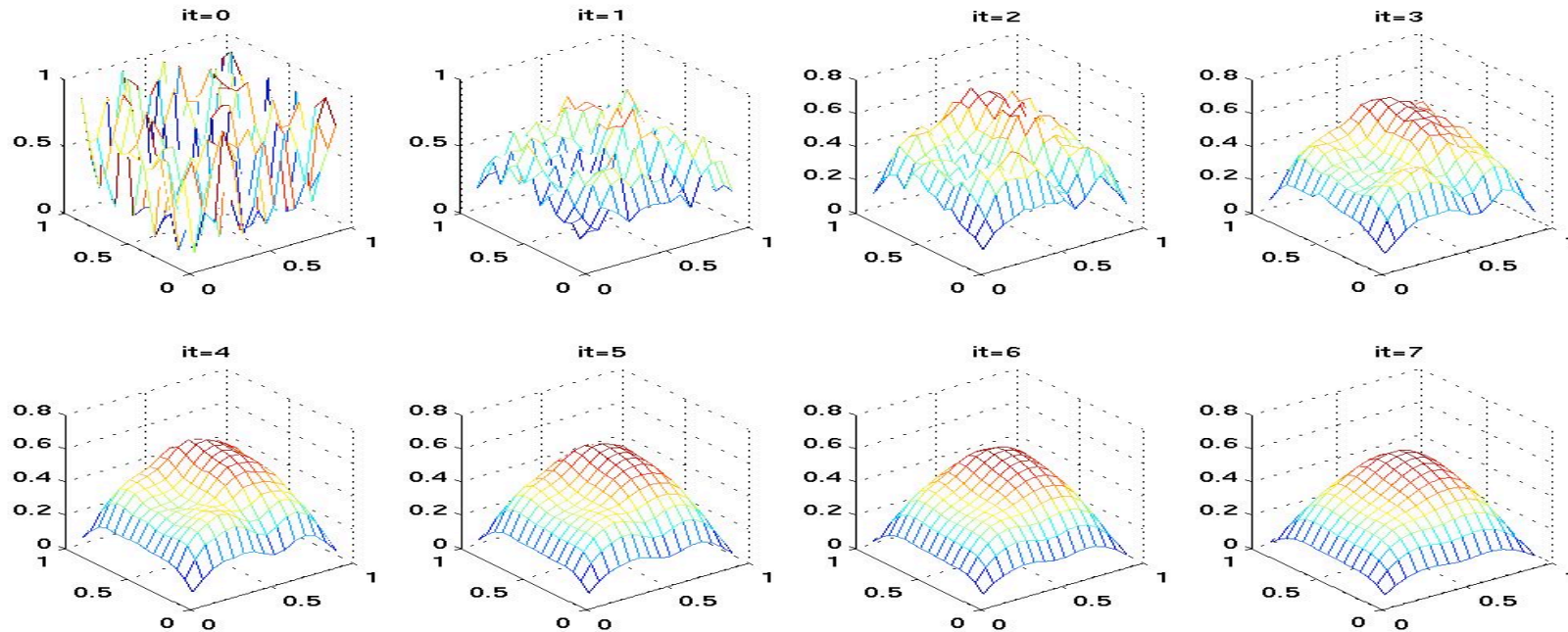
1. Simple Markov Chain Example

- 5 states
- each outgoing edge same probability (random walk on directed graph)



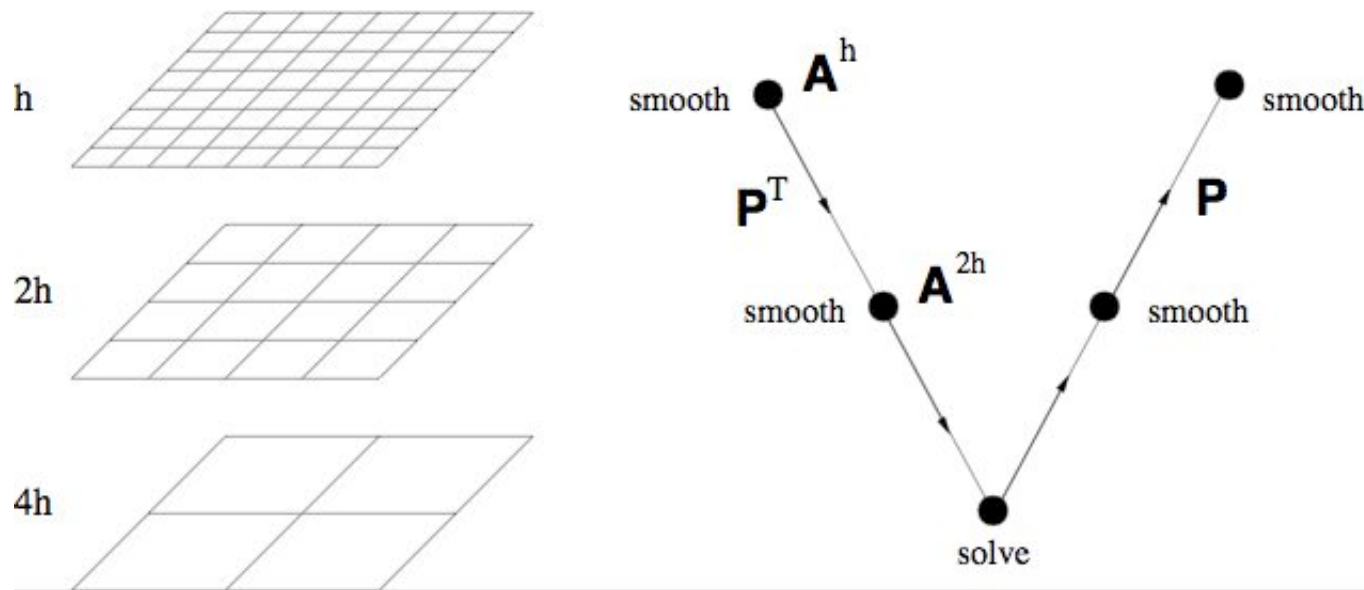
Principle of Multigrid (for PDEs)

$$-u_{xx} - u_{yy} = f(x, y) \quad Ax = b$$



- **high-frequency error** is removed by **relaxation** (weighted Jacobi, Gauss-Seidel, ...)
- **low-frequency-error** needs to be removed by **coarse-grid correction**

Multigrid Hierarchy: V-cycle



- multigrid V-cycle:
 - **relax** (=smooth) on successively coarser grids
 - transfer error using **restriction** ($R=P^T$) and **interpolation** (P)
- $W=O(n)$

Smoothed Aggregation

$$\begin{aligned} A_{cs} &= R_s (D - (L + U)) P_s \\ &= R_s D P_s - R_s (L + U) P_s \end{aligned}$$

$$A = \begin{bmatrix} + & - & - & - & - \\ - & + & - & - & - \\ - & - & + & - & - \\ - & - & - & + & - \\ - & - & - & - & + \end{bmatrix}$$

$$A_{cs} = S - G \qquad \hat{A}_{cs} = \hat{S} - G$$

- we want to retain crucial properties

$$\begin{aligned} \mathbf{1}_c^T \hat{A}_{cs} &= 0 \quad \forall \mathbf{x}_i, \\ \hat{A}_{cs} \mathbf{1}_c &= 0 \quad \text{for } \mathbf{x}_i = \mathbf{x} \end{aligned}$$

- we can lump to diagonal in symmetric way, conserving both row and column sums

Numerical Results: Geometric Aggregation (size 3)

n	γ_{res}	iter	C_{op}	levels
54	0.86	75	1.43	3
162	0.95	>100	1.47	4
486	0.97	>100	1.49	5
1458	0.98	>100	1.50	6

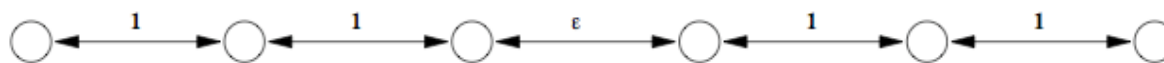
TABLE 5.5

Uniform chain with two weak links ($\epsilon = 0.001$). G-AM with V-cycles and size-three aggregates. The two weak links occur between aggregates at all levels. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	R_{lump}
54	0.26	12	1.43	3	0
162	0.27	13	1.47	4	0
486	0.27	13	1.49	5	0
1458	0.27	13	1.50	6	0
4374	0.27	13	1.50	7	0

TABLE 5.6

Uniform chain with two weak links ($\epsilon = 0.001$). G-SAM with V-cycles and size-three aggregates. The two weak links occur between aggregates at all levels. (Smoothing with lumping.)



Numerical Results: Geometric Aggregation (size 3)

n	γ_{res}	iter	C_{op}	levels
27	1.00	>100	1.32	2
81	1.00	>100	1.43	3
243	0.98	>100	1.47	4
729	0.98	>100	1.49	5

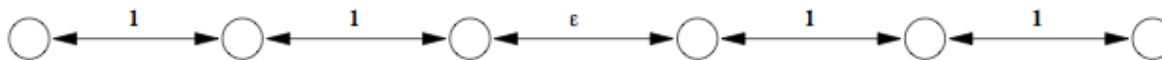
TABLE 5.7

Uniform chain with two weak links ($\epsilon = 0.001$). G-AM with V-cycles and size-three aggregates. The two weak links occur inside an aggregate on the finest level. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	R_{lump}
27	1.00	>100	1.32	2	0
81	1.00	>100	1.43	3	0
243	1.00	>100	1.47	4	0
729	1.00	>100	1.49	5	0

TABLE 5.8

Uniform chain with two weak links ($\epsilon = 0.001$). G-SAM with V-cycles and size-three aggregates. The two weak links occur inside an aggregate on the finest level. (Smoothing with lumping.)



Numerical Results: effect of η

n	γ_{res}	iter	C_{op}	levels
27	0.75	43	1.71	3
81	0.87	87	1.85	4
243	0.96	>100	1.96	6
729	0.99	>100	1.98	7

TABLE 6.1

Uniform chain. A-AM with V-cycles and distance-one aggregation. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	γ_{res}	iter	C_{op}	levels
27	0.20	10	2.03	3	0.19	10	2.03	3
81	0.27	11	2.69	4	0.25	11	2.59	4
243	0.32	12	3.33	6	0.32	12	3.33	6
729	0.51	14	3.73	8	0.64	17	3.75	8
2187	0.75	21	3.95	9	0.77	24	3.94	9

TABLE 6.3

Uniform chain. A-SAM with V-cycles using distance-one aggregation, lumping only the off-diagonal elements of $R_s DP_s$ that cause nonnegative off-diagonal elements of A_{cs} . Lumping their full value ($\eta = 1$, left), and part of their value ($\eta = 0.75$, right).

Numerical Results: effect of η

n	γ_{res}	iter	C_{op}	levels	γ_{res}	iter	C_{op}	levels
27	0.19	10	2.03	3	0.19	10	1.96	3
81	0.19	10	2.51	4	0.20	10	2.61	4
243	0.24	11	2.96	5	0.22	11	3.20	5
729	0.37	12	3.63	7	0.29	12	3.29	6
2187	0.44	14	3.84	8	0.28	12	3.35	7

TABLE 6.4

Uniform chain. A-SAM with V-cycles using distance-one aggregation, lumping part of the value of the off-diagonal elements of $R_s D P_s$ that cause nonnegative off-diagonal elements of A_{CS} : $\eta = 0.25$, left, and, $\eta = 0.1$, right.

n	γ_{res}	iter	C_{op}	levels	γ_{res}	iter	C_{op}	levels
27	0.20	10	2.03	3	0.20	10	2.03	3
81	0.18	10	2.68	4	0.18	10	2.71	4
243	0.19	10	2.78	5	0.20	10	3.03	5
729	0.24	11	3.41	7	0.24	11	3.50	7
2187	0.27	11	3.75	8	0.26	11	3.81	8

TABLE 6.5

Uniform chain. A-SAM with V-cycles using distance-one aggregation, lumping part of the value of the off-diagonal elements of $R_s D P_s$ that cause nonnegative off-diagonal elements of A_{CS} : $\eta = 0.01$, left, and, $\eta = 1e - 6$, right.

Numerical Results: Algebraic Aggregation

n	γ_{res}	iter	C_{op}	levels
27	0.75	50	1.71	3
81	0.87	92	1.85	4
243	0.96	>100	1.96	6
729	0.97	>100	1.99	8

TABLE 6.6

Birth-death chain ($\mu = 0.96$). A-AM with V-cycles and distance-one aggregation. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	R_{lump}
27	0.27	12	1.32	2	0
81	0.35	15	1.43	3	0
243	0.35	15	1.47	4	0
729	0.35	15	1.49	5	0

TABLE 6.7

Birth-death chain ($\mu = 0.96$). A-SAM with V-cycles and distance-two aggregation. (Smoothing with lumping.)

Numerical Results: Algebraic Aggregation

n	γ_{res}	iter	C_{op}	levels
64	0.73	41	1.73	3
100	0.80	56	1.83	4
169	0.85	77	1.85	4
400	0.89	>100	1.96	6
900	0.96	>100	1.96	6

TABLE 6.10

Uniform 2D lattice. A-AM with V-cycles and distance-one aggregation. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	R_{lump}
64	0.42	18	1.30	3	0
100	0.45	19	1.26	3	3.46e-03
169	0.44	18	1.33	3	9.47e-03
400	0.47	20	1.46	4	9.27e-03
900	0.46	18	1.59	4	1.72e-02
1600	0.48	19	1.60	4	1.16e-02
2500	0.48	19	1.67	5	1.44e-02
4900	0.48	18	1.75	5	1.21e-02
6724	0.49	18	1.76	5	1.36e-02

TABLE 6.11

Uniform 2D lattice. A-SAM with V-cycles and distance-two aggregation. (Smoothing with lumping.)

4. Aggregation for Markov Chains

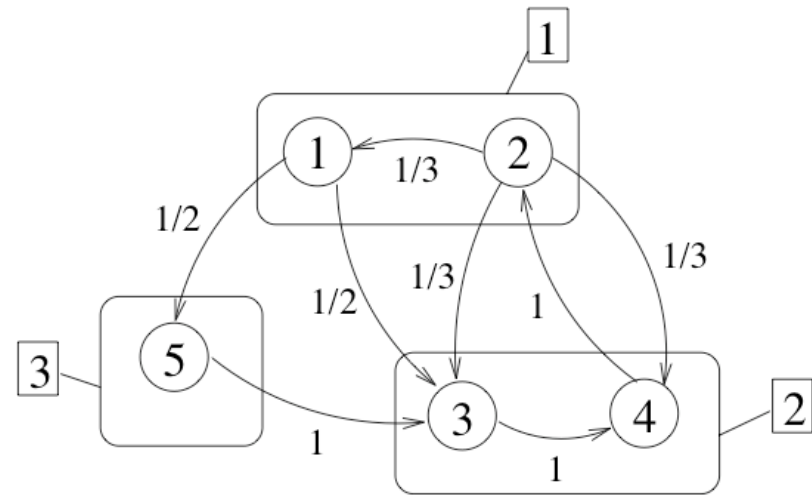
- form three coarse, aggregated states

$$\mathbf{x}_{c,I} = \sum_{i \in I} x_i$$

$$\mathbf{x}_c^T = [8/19 \quad 10/19 \quad 1/19]$$

$$B_c \mathbf{x}_c = \mathbf{x}_c$$

$$b_{c,IJ} = \frac{\sum_{j \in J} x_j \left(\sum_{i \in I} b_{ij} \right)}{\sum_{j \in J} x_j}$$



$$B_c = \begin{bmatrix} 1/4 & 3/5 & 0 \\ 5/8 & 2/5 & 1 \\ 1/8 & 0 & 0 \end{bmatrix}$$

(Simon and Ando, 1961)

5. Error Equation

- multiplicative correction: error equation, coarse level error equation, and coarse grid correction

$$\mathbf{x} = \text{diag}(\mathbf{x}_i) \mathbf{e}_i$$

$$A \text{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$$

$$Q^T A \text{diag}(\mathbf{x}_i) Q \mathbf{e}_c = 0$$

$$A_c \mathbf{e}_c = 0$$

$$R = Q^T \quad P = \text{diag}(\mathbf{x}_i) Q$$

$$A_c = R A P$$

5. Error Equation

- error equation - coarse grid correction:

$$\mathbf{x} = \text{diag}(\mathbf{x}_i) \mathbf{e}_i$$
$$A \text{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$$

$$Q^T A \text{diag}(\mathbf{x}_i) Q \mathbf{e}_c = 0$$

$$A_c \mathbf{e}_c = 0$$

$$R = Q^T \quad P = \text{diag}(\mathbf{x}_i) Q$$

$$A_c = R A P$$

$$\mathbf{x}_{i+1} = P \mathbf{e}_c$$

$$\mathbf{x}_c = \text{diag}(P^T \mathbf{1}) \mathbf{e}_c$$

$$A_c (\text{diag}(P^T \mathbf{1}))^{-1} \mathbf{x}_c = 0$$

$$\mathbf{x}_{i+1} = P (\text{diag}(P^T \mathbf{1}))^{-1} \mathbf{x}_c$$

Numerical Results: Algebraic Aggregation

n	γ_{res}	iter	C_{op}	levels
54	0.83	67	1.86	4
162	0.93	>100	1.91	5
486	0.96	>100	1.98	7
1458	0.97	>100	1.99	9

TABLE 6.8

Uniform chain with two weak links ($\epsilon = 0.001$). A-AM with V-cycles and distance-one aggregation. (No smoothing.)

n	γ_{res}	iter	C_{op}	levels	R_{lump}
54	0.33	14	1.51	3	0
162	0.33	13	1.50	4	5.51e-03
486	0.34	13	1.50	5	0
1458	0.29	12	1.49	6	3.07e-04
4374	0.27	11	1.50	7	3.05e-04

TABLE 6.9

Uniform chain with two weak links ($\epsilon = 0.001$). A-SAM with V-cycles and distance-two aggregation. (Smoothing with lumping.)

