Markov Chains and Web Ranking: a Multilevel Adaptive Aggregation Method

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1. Simple Markov Chain Example



 each outgoing edge same probability (random walk on directed graph)





Simple Markov Chain Example

 start in one state with probability 1: what is the stationary probability vector after ∞ number of steps?

$$\mathbf{x}_{i+1} = B \, \mathbf{x}_i$$

• stationary probability:

$$B\mathbf{x} = \mathbf{x} \qquad \|\mathbf{x}\|_1 = 1$$

$$B = \begin{bmatrix} 0 & 1/3 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 1/2 & 1/3 & 0 & 0 & 1 \\ 0 & 1/3 & 1 & 0 & 0 \\ 1/2 & 0 & 0 & 0 & 0 \end{bmatrix}$$



$$\mathbf{x}^T = [2/19 \, 6/19 \, 4/19 \, 6/19 \, 1/19]$$



2. Problem Statement

 $B\mathbf{x} = \mathbf{x} \qquad \|\mathbf{x}\|_1 = \mathbf{1}$

• *B* is column-stochastic

 $b_{i,j} \ge 0, \qquad \sum_i b_{i,j} = 1 \,\,\forall j$

• *B* is irreducible (every state can be reached from every other state in the directed graph)

 \Rightarrow $\exists ! \mathbf{x} : B \mathbf{x} = \mathbf{x} \text{ and } \|\mathbf{x}\|_1 = 1, \qquad x_i > 0 \ \forall i$ (no probability sinks!)





Problem Statement

$$B\mathbf{x} = \mathbf{x} \qquad \|\mathbf{x}\|_1 = 1$$

- *B* is column-stochastic •
- *B* is irreducible
- *B* is a-periodic

 \Rightarrow

$$\mathbf{x} = \lim_{n \to \infty} B^n \mathbf{x}_0$$
 for

$$\operatorname{im}_{n\to\infty} B^n \mathbf{x}_0$$
 for any \mathbf{x}_0

• largest eigenvalue of B:
$$\lambda_1 = 1$$





3. Traditional, One-Level Iterative Methods

Power Method

$$\mathbf{x}_{i+1} = B\mathbf{x}_i$$

• Weighted Jacobi Method (WJAC)

$$B \mathbf{x} = \mathbf{x}$$
$$(B - I) \mathbf{x} = 0 \qquad A \mathbf{x} = 0$$
$$A = L + D + U$$

$$\mathbf{x}_{i+1} = N((1-w)\mathbf{x}_i + wD^{-1}(L+U)\mathbf{x}_i)$$

normalization:

$$N(\mathbf{x}) = \frac{\mathbf{x}}{\|\mathbf{x}\|_1} \qquad (\mathbf{x} \neq \mathbf{0})$$



Traditional, One-Level Iterative Methods

- Power Method
 - convergence factor: $|\lambda_2|$
 - convergence is very slow when $|\lambda_2| \approx 1$
- WJAC: similar



4. Aggregation

• form three coarse, aggregated states

$$x_{c,I} = \sum_{i \in I} x_i$$

$$\mathbf{x}_c^T = [8/19 \ 10/19 \ 1/19]$$

 $B_c \mathbf{x}_c = \mathbf{x}_c$







$$B_c = \begin{bmatrix} 1/4 & 3/5 & 0\\ 5/8 & 2/5 & 1\\ 1/8 & 0 & 0 \end{bmatrix}$$

(Simon and Ando, 1961)

Aggregation





$$\mathbf{x}_c = P^T \mathbf{x}$$

$$B_c = P^T B \operatorname{diag}(\mathbf{x}) P \operatorname{diag}(P^T \mathbf{x})^{-1}$$

(Krieger, Horton, ...)





5. Two-Level Acceleration by Aggregation

Algorithm: Two-level acceleration by agglomeration

choose initial guess ${\bf x}$

repeat

$$\begin{aligned} \mathbf{x} &\leftarrow N(\mathsf{Relax}(A, \mathbf{x})) & \nu \text{ times} \\ A_c &= P^T A \operatorname{diag}(\mathbf{x}) P \operatorname{diag}(P^T \mathbf{x})^{-1} \\ \mathbf{x}_c &\leftarrow \text{solve } A_c \, \mathbf{x}_c = 0, \ \|\mathbf{x}_c\|_1 = 1 \quad \text{(coarse-level solve)} \\ \mathbf{x} &= N(\operatorname{diag}(P \operatorname{diag}(P^T \mathbf{x})^{-1} \mathbf{x}_c) \, \mathbf{x}) \quad \text{(coarse-level correction)} \\ \end{aligned}$$
end

≈ Iterative Aggregation/Disaggregation Algorithm (IAD) (Takahashi 1975, Koury et al 1984, Schweitzer and Kindle 1986, Krieger 1990, Marek and Mayer 1998, ...)



6. Relation with Multigrid for PDEs

$$-u_{xx} - u_{yy} = f(x, y)$$



ο

0.5

0 0

• high-frequency error is removed by relaxation

0

0

0

0.5

1

0.5

• low-frequency-error needs to be removed by coarse-grid correction

0.5



0

0.5

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0.5

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0.5

0.5

0

0

multigrid hierarchy: V-cycle



- multigrid V-cycle:
 - relax (=smooth) on various grids
 - ♣ transfer error using restriction (P^T) and interpolation (P)



7. Multilevel Adaptive Aggregation (MAA) Method

Algorithm: Multilevel Adaptive Aggregation method (V-cycle)

$$\mathbf{x} = \mathsf{MAA}_{-}\mathsf{V}(A, \mathbf{x}, \nu_1, \nu_2)$$

begin

 $\mathbf{x} \leftarrow N(\operatorname{Relax}(A, \mathbf{x})) \qquad \nu_1 \text{ times}$ build *P* based on **x** and *A* (*P* is rebuilt every level and cycle) $A_c = P^T A \operatorname{diag}(\mathbf{x}) P \operatorname{diag}(P^T \mathbf{x})^{-1}$ $\mathbf{x}_c = \mathsf{MAA}_- \mathsf{V}(A_c, N(P^T \mathbf{x}), \nu_1, \nu_2) \quad \text{(coarse-level solve)}$ $\mathbf{x} = N(\operatorname{diag}(P \operatorname{diag}(P^T \mathbf{x})^{-1} \mathbf{x}_c) \mathbf{x}) \quad \text{(coarse-level correction)}$ $\mathbf{x} \leftarrow N(\operatorname{Relax}(A, \mathbf{x})) \qquad \nu_2 \text{ times}$ end

(Krieger, Horton 1994, but no satisfactory way to build P) Waterloo KiciAM 2007 hdesterck@uwaterloo.ca

8. Choosing Aggregates Based on Strength

- error equation: $A \operatorname{diag}(\mathbf{x}_i) \mathbf{e}_i = 0$
- use strength of connection in $A \operatorname{diag}(\mathbf{x}_i)$
- define row-based strength (determine all states that strongly influence a row's state)
- state that has largest value in x_i is seed point for new aggregate, and all unassigned states influenced by it join its aggregate
- repeat

(similar to Brandt, McCormick and Ruge, 1983)



- web adjacency matrix G, then B=N(G)
- needs to be made irreducible, a-periodic
- need to add in extra links, with 'coupling factor' α





• PageRank (used by Google):

$$B_{PR} = (1 - \alpha) N(G + e d^T) + \alpha N(e e^T)$$

(\alpha = 0.15)





• BackButton (add reverse of each link):

$$B_{BB} = N((1 - \alpha - \epsilon) N(G) + \alpha N(G^T) + \epsilon I)$$





• second eigenvalue for PageRank: $\leq 1-\alpha$ (α =0.15 and α =0.01)



• second eigenvalue for BackButton (α =0.15 and α =0.01): (c) BackButton





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• second eigenvalue for BackButton:





10. Performance of MAA

• total efficiency factor of MAA relative to WJAC:

$$f_{MAA-WJAC}^{(tot)} = \frac{\log(r_{MAA})/t_{MAA}}{\log(r_{WJAC})/t_{WJAC}},$$

• $f^{(tot)} = 2$: MAA 2 times more efficient than WJAC

• $f^{(tot)} = 1/2$: MAA 2 times *less* efficient than WJAC



Performance of MAA for PageRank

	1						
n	γ_{MAA}	it_{MAA}	$c_{grid,MAA}$	γ_{WJAC}	$f_{MAA-WJAC}^{(tot)}$	$f_{MAA-WJAC}^{(as)}$	
PageRank, $\alpha = 0.15$							
2000	0.355124	10	1.67	0.815142	1/2.74	1/3.13	
4000	0.335889	9	1.67	0.805653	1/2.52	1/3.59	
8000	0.387411	9	1.65	0.821903	1/2.79	1/4.14	
16000	0.554686	12	1.78	0.836429	1/4.07	1/6.89	
32000	0.502008	11	1.83	0.833367	1/3.94	1/6.20	
64000	0.508482	11	1.75	0.829696	1/3.86	1/6.21	
128000	0.532518	12	1.75	0.829419	1/4.31	1/7.01	
PageRank, $\alpha = 0.01$							
2000	0.321062	10	1.77	0.956362	3.42	1.32	
4000	0.658754	20	1.75	0.980665	2.16	1.03	
8000	0.758825	22	1.65	0.976889	1.88	1/1.65	
16000	0.815774	27	1.77	0.979592	1.45	1/2.31	
32000	0.797182	29	1.82	0.979881	1.35	1/2.09	
64000	0.786973	33	1.79	0.980040	1.19	1/1.96	
128000	0.854340	38	1.72	0.980502	1.05	1/2.88	



Performance of MAA for BackButton

		<i>it</i>			f(tot)	f(as)	
n	γMAA	u_{MAA}	$c_{grid,MAA}$	$\gamma WJAC$	JMAA-WJAC	JMAA-WJAC	
BackButton, $\alpha = 0.15$							
2000	0.746000	35	1.74	0.981331	2.36	1/1.41	
4000	0.800454	39	1.64	0.982828	2.70	1/1.36	
8000	0.786758	40	1.53	0.992129	3.15	1.17	
16000	0.851671	50	1.62	0.992330	3.00	1/1.38	
32000	0.988423	214	1.64	0.998366	4.92	1/2.88	
64000	0.973611	185	1.59	0.999013	9.95	1.40	
128000	0.943160	116	1.55	0.999693	34.64	9.90	
BackButton, $\alpha = 0.01$							
2000	0.658032	23	1.68	0.999563	106.02	46.05	
4000	0.794123	29	1.71	0.999345	73.02	19.78	
8000	0.841182	39	1.70	0.997624	23.49	2.64	
16000	0.835592	44	1.78	0.998696	19.72	4.42	
32000	0.845457	56	1.83	0.999114	39.58	8.22	
64000	0.959561	81	1.75	0.999660	75.05	5.74	
128000	0.921870	42	1.70	0.999963	816.62	103.79	



11. Web Matrix Regularizations as a Function of Coupling Factor α



Web Matrix Regularizations as a Function of Coupling Factor α



12. Conclusions

• PageRank regularization of webgraph with α =0.15 seems to be a pretty good model for web ranking



Conclusions

 MAA cannot beat Power Method speed for PageRank

reasons:

o λ_2 =0.85, independent of n

- o there are global connections on all scales
- o Power Method scales perfectly (also in parallel)
- o 0.85¹⁰=0.2 0.85²⁰=0.04

still:

o MAA provides information about how web pages are clustered, at all levels



Conclusions

 MAA can dramatically outperform Power Method for Markov Chains for which |λ₂(n)| → 1 for large n

reason:

o multilevel nature of MAA allows to bridge scales (Markov Chain has only local connections in this case)



Comparison: Strength-based Aggregation

- major differences with previously considered multilevel methods for Markov chains (Horton, Krieger,...):
 - o we use AMG strength of connection based on scaled problem matrix $A \operatorname{diag}(\mathbf{x}_i)$
 - o our aggregates based on columns of strength matrix (~ AMG coarsening)
- IAD methods
 - o normally only 2-level
 - o aggregates fixed, based on previously known, regular structure of Markov chain



Two-Level Acceleration by Aggregation

• multiplicative correction: error equation

 $\mathbf{x} = \mathsf{diag}(\mathbf{x}_i) \, \mathbf{e}_i$

 $A \operatorname{diag}(\mathbf{x}_i) \mathbf{e}_i = \mathbf{0}$

$$P^T A \operatorname{diag}(\mathbf{x}_i) P \mathbf{e}_c = \mathbf{0}$$

 $\mathbf{x}_c = \operatorname{diag}(P^T \mathbf{x}_i) \mathbf{e}_c$

$$A_c \mathbf{x}_c = \mathbf{0}$$

 $P^T A \operatorname{diag}(\mathbf{x}) P \operatorname{diag}(P^T \mathbf{x})^{-1} \mathbf{x}_c = 0$



• BackLink (to root page):

$$B_{BL} = N((1 - \alpha - \epsilon) N(G) + \alpha e^{(1)} e^{T} + \epsilon I)$$





 second eigenvalue for BackLink (α=0.15 and α =0.01):
(b) BackLink





Performance of MAA

n	$\gamma_{\Lambda I \Lambda \Lambda}$	it NA AA	Consid MAA		f(tot)	$f^{(as)}$	
			gria,MAA	/WJAC	J MAA-W JAC	J MAA-WJAC	
Backlink, $\alpha = 0.15$							
2000	0.331226	11	1.67	0.839540	1/3.11	1/3.04	
4000	0.344225	11	1.75	0.851397	1/3.30	1/3.18	
8000	0.361255	11	1.69	0.858532	1/3.04	1/3.24	
16000	0.358282	11	2.03	0.866344	1/3.75	1/4.11	
32000	0.369351	11	2.26	0.868116	1/3.99	1/4.39	
64000	0.368789	11	1.88	0.868889	1/3.30	1/3.53	
128000	0.369744	11	1.78	0.871525	1/3.07	1/3.22	
BackLink, $\alpha = 0.01$							
2000	0.452383	16	1.89	0.952865	2.01	1/1.21	
4000	0.778003	28	1.76	0.953782	1.41	1/4.23	
8000	0.749847	20	1.72	0.970096	2.23	1/2.23	
16000	0.745776	23	1.96	0.976919	1.87	1/2.11	
32000	0.855323	28	1.93	0.981223	1.66	1/3.04	
64000	0.868049	32	1.96	0.983076	1.45	1/3.15	
128000	0.837747	31	1.83	0.985161	1.65	1/2.09	



Web Matrix Regularizations as a Function of Coupling Factor α



Web Matrix Regularizations as a Function of Coupling Factor α

