C&O 355 Lecture 16

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# Topics

- Review of Fourier-Motzkin Elimination
- Linear Transformations of Polyhedra
- Convex Combinations
- Convex Hulls
- Polytopes & Convex Hulls



#### Fourier-Motzkin Elimination



Joseph Fourier

Given a polyhedron Q ⊆ ℝ<sup>n</sup>,
 we want to find the set Q' ⊆ ℝ<sup>n-1</sup> satisfying

 $(\mathbf{x}_{1},\ldots,\mathbf{x}_{n-1}) \in \mathbf{Q'} \quad \Leftrightarrow \quad \exists \mathbf{x}_{n} \text{ s.t. } (\mathbf{x}_{1},\ldots,\mathbf{x}_{n-1},\mathbf{x}_{n}) \in \mathbf{Q}$ 

- Q' is called the projection of Q onto first n-1 coordinates
- Fourier-Motzkin Elimination constructs Q' by generating (finitely many) constraints from the constraints of Q.
- **Corollary:** Q' is a polyhedron.

#### **Elimination Example**



• Project Q onto coordinates {x<sub>1</sub>, x<sub>2</sub>}...



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- Fourier-Motzkin: Q' is a polyhedron.
- Of course, the ordering of coordinates is irrevelant.

# Elimination Example



- Of course, the ordering of coordinates is irrevelant.
- Fourier-Motzkin: Q" is also a polyhedron.
- I can also apply Elimination twice...



• Fourier-Motzkin: Q''' is also a polyhedron.

# Projecting a Polyhedron Onto Some of its Coordinates

• Lemma: Given a polyhedron  $\mathbf{Q} \subseteq \mathbb{R}^n$ .

Let  $S=\{s_1,...,s_k\} \subseteq \{1,...,n\}$  be any subset of the coordinates. Let  $Q_S = \{(x_{s1},...,x_{sk}) : x \in Q\} \subseteq \mathbb{R}^k$ .

In other words,  $Q_S$  is projection of Q onto coordinates in S. Then  $Q_S$  is a polyhedron.

• Proof:

Direct from Fourier-Motzkin Elimination. Just eliminate all coordinates not in S.

#### Linear Transformations of Polyhedra

 Lemma: Let P = { x : Ax≤b } ⊆ ℝ<sup>n</sup> be a polyhedron. Let M be any matrix of size p<sub>x</sub>n. Let Q = { Mx : x∈P } ⊆ ℝ<sup>p</sup>. Then Q is a polyhedron.



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Geometrically obvious, but not easy to prove...



#### Linear Transformations of Polyhedra

Lemma: Let P = { x : Ax≤b } ⊆ ℝ<sup>n</sup> be a polyhedron.
 Let M be any matrix of size pxn.
 Let Q = { Mx : x∈P } ⊆ ℝ<sup>p</sup>. Then Q is a polyhedron.

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...unless you know Fourier-Motzkin Elimination!

#### • Proof:

Let P' = { (x,y) : Mx=y, Ax $\leq$ b }  $\subseteq \mathbb{R}^{n+p}$ , where x $\in \mathbb{R}^{n}$ , y $\in \mathbb{R}^{p}$ .

P' is obviously a polyhedron.

Note that Q is projection of P' onto y-coordinates.

By previous lemma, Q is a polyhedron.

#### **Convex Sets**

- Let  $S \subseteq \mathbb{R}^n$  be any set.
- Recall: S is convex if  $\lambda x + (1 \lambda)y \in S \quad \forall x, y \in S \text{ and } \forall \lambda \in [0, 1]$
- **Proposition:** The intersection of any collection of convex sets is itself convex.
- Proof: Exercise.

#### **Convex Combinations**

- Let  $S \subseteq \mathbb{R}^n$  be any set.
- **Definition:** A convex combination of points in S is any point  $p = \sum_{i=1}^{k} \lambda_i s_i$ where k is finite and  $s_i \in S \ \forall i, \ \lambda_i \geq 0 \ \forall i, \ \sum_{i=1}^{k} \lambda_i = 1$
- Theorem: (Carathéodory's Theorem) It suffices to take k≤n+1.
- **Proof:** This was Assignment 2, Problem 4.

#### **Convex Combinations**

- Theorem: Let S⊆R<sup>n</sup> be a convex set.
  Let comb(S) = { p : p is a convex comb. of points in S }.
  Then comb(S) = S.
- **Proof:**  $S \subseteq comb(S)$  is trivial.

#### **Convex Combinations**

- Theorem: Let S⊆R<sup>n</sup> be a convex set.
  Let comb(S) = { p : p is a convex comb. of points in S }.
  Then comb(S) = S.
- **Proof:** We'll show comb(S)  $\subseteq$  S. Consider  $p = \sum_{i=1}^{k} \lambda_i s_i$ . Need to show  $p \in$ S. By induction on k. Trivial if k=1 or if  $\lambda_k$ =1. So assume k>1 and  $\lambda_k$ <1. Note:  $p = \sum_{i=1}^{k} \lambda_i s_i = (1 - \lambda_k) \left( \sum_{i=1}^{k-1} \frac{\lambda_i}{1 - \lambda_k} s_i \right) + \lambda_k s_k$

Call this point p'

Note  $\sum_{i=1}^{k-1} \frac{\lambda_i}{1-\lambda_k} = 1$  so p' is a convex combination of at most k-1 points in S. By induction, p' $\in$ S.

But  $p = (1 - \lambda_k)p' + \lambda_k s_k$ , so p $\in$ S since S is convex.

- Let  $S \subseteq \mathbb{R}^n$  be any set.
- **Definition:** The **convex hull** of S, denoted **conv(S)**, is the intersection of all convex sets containing S.



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- Let  $S \subseteq \mathbb{R}^n$  be any set.
- **Definition:** The **convex hull** of S, denoted **conv(S)**, is the intersection of all convex sets containing S.
- **Claim:** conv(S) is itself convex.
- Proof:

Follows from our earlier proposition "Intersection of convex sets is convex".

- Let  $S \subseteq \mathbb{R}^n$  be any set.
- **Definition:** The **convex hull** of S, denoted **conv(S)**, is the intersection of all convex sets containing S.
- Let comb(S) = { p : p is a convex comb. of points in S }.
- Theorem: conv(S)=comb(S).
- **Proof:** We'll show conv(S)⊆comb(S).

Claim: comb(S) is convex.

**Proof:** Consider  $p = \sum_{i=1}^{k} \lambda_i s_i$  and  $q = \sum_{j=1}^{\ell} \mu_j t_j$ , where  $\lambda \ge 0, \ \mu \ge 0, \ \sum_i \lambda_i = 1, \ \sum_j \mu_j = 1, \ s_i \in S, \ t_j \in S$ 

So p,q $\in$ comb(S). For  $\alpha \in$  [0,1], consider

$$\alpha p + (1 - \alpha)q = \sum_{i=1}^{k} \alpha \lambda_i s_i + \sum_{j=1}^{\ell} (1 - \alpha)\mu_j t_j$$
  
Thus  $\alpha p$  + (1- $\alpha$ )q  $\in$  comb(S). [

- Let  $S \subseteq \mathbb{R}^n$  be any set.
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- Let comb(S) = { p : p is a convex comb. of points in S }.
- Theorem: conv(S)=comb(S).
- **Proof:** We'll show conv(S)⊆comb(S).

Claim: comb(S) is convex.

Clearly comb(S) contains S.

But conv(S) is the intersection of **all** convex sets containing S, so conv(S) $\subseteq$ comb(S).

- Let  $S \subseteq \mathbb{R}^n$  be any set.
- **Definition:** The **convex hull** of S, denoted **conv(S)**, is the intersection of all convex sets containing S.
- Let comb(S) = { p : p is a convex comb. of points in S }.
- Theorem: conv(S)=comb(S).
- **Proof:** Exercise: Show comb(S)  $\subseteq$  conv(S).

#### Convex Hulls of Finite Sets

 Theorem: Let S={s<sub>1</sub>,...,s<sub>k</sub>}⊂ℝ<sup>n</sup> be a finite set. Then conv(S) is a polyhedron.

Geometrically obvious, but not easy to prove...



#### **Convex Hulls of Finite Sets**

- Theorem: Let S={s<sub>1</sub>,...,s<sub>k</sub>}⊂R<sup>n</sup> be a finite set. Then conv(S) is a polyhedron.
   Geometrically obvious, but not easy to prove... ...unless you know Fourier-Motzkin Elimination!
- Proof: Let M be the n<sub>x</sub>k matrix where M<sub>i</sub> = s<sub>i</sub>.
  By our previous theorem,

 $\operatorname{conv}(S) = \{ y : \exists x \text{ s.t. } Mx = y, \sum_{i=1}^{k} x_i = 1, x \ge 0 \}$ But this is a projection of the polyhedron

$$\{ (x,y) : Mx = y, \sum_{i=1}^{k} x_i = 1, x \ge 0 \}$$

By our lemma on projections of polyhedra, conv(S) is also a polyhedron.

- Let  $P \subset \mathbb{R}^n$  be a **polytope.** (i.e., a **bounded** polyhedron)
- Is P the convex hull of anything?
- Since P is convex, P = conv(P). Too obvious...
- Maybe P = conv( extreme points of P )?



- Theorem: Let P⊂ℝ<sup>n</sup> be a non-empty polytope.
  Then P = conv( extreme points of P ).
- Proof: First we prove conv( extreme points of P ) ⊆ P.
  We have:

conv( extreme points of P )  $\subseteq$  conv( P )  $\subseteq$  P.

Obvious

Our earlier theorem proved comb(P)  $\subseteq$  P

- Theorem: Let P⊂ℝ<sup>n</sup> be a non-empty polytope.
  Then P = conv( extreme points of P ).
- Proof: Now we prove P ⊆ conv( extreme points of P ). Let the extreme points be {v<sub>1</sub>,...,v<sub>k</sub>}. (Finitely many!) Suppose ∃b ∈ P \ conv( extreme points of P ). Then the following system has no solution:

$$\sum_{i=1}^{k} v_i \lambda_i = b$$
$$\sum_{i=1}^{k} \lambda_i = 1$$
$$\lambda_i \ge 0 \ \forall i$$

By Farkas' lemma,  $\exists u \in \mathbb{R}^n$  and  $\alpha \in \mathbb{R}$  s.t.

 $u^{\mathsf{T}}v_i + \alpha \ge 0 \ \forall i$  and  $u^{\mathsf{T}}b + \alpha < 0$ So  $-u^{\mathsf{T}}b > -u^{\mathsf{T}}v_i$  for every extreme point  $v_i$  of P.

- Theorem: Let P⊂ℝ<sup>n</sup> be a non-empty polytope.
  Then P = conv( extreme points of P ).
- **Proof:** Now we prove  $P \subseteq \text{conv}(\text{ extreme points of } P)$ . Let the extreme points be  $\{v_1, ..., v_k\}$ . (Finitely many!) Suppose  $\exists b \in P \setminus \text{conv}(\text{ extreme points of } P)$ . Then the following system has no solution:

By Farkas' lemma,  $\exists u \in \mathbb{R}^n$  and  $\alpha \in \mathbb{R}$  s.t.

 $u^{\mathsf{T}}v_i + \alpha \ge 0 \ \forall i$  and  $u^{\mathsf{T}}b + \alpha < 0$ So  $-u^{\mathsf{T}}b > -u^{\mathsf{T}}v_i$  for every extreme point  $v_i$  of P. Consider the LP max {  $-u^{\mathsf{T}}x : x \in P$  }. It is not unbounded, since P is bounded. Its optimal value is **not attained at an extreme point**. This is a contradiction.

# Generalizations

#### • Theorem:

Every polytope in  $\mathbb{R}^n$  is the convex hull of its extreme points.



- Theorem: [Minkowski 1911]
  Every compact, convex set in R<sup>n</sup> is the convex hull of its extreme points.
- Theorem: [Krein & Milman 1940] Every compact convex subset of a locally convex Hausdorff linear space is the closed convex hull of its extreme points.

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