C&O 355 Lecture 15

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# Topics

- Subgradient Inequality
- Characterizations of Convex Functions
- Convex Minimization over a Polyhedron
- (Mini)-KKT Theorem
- Smallest Enclosing Ball Problem

# Subgradient Inequality

- **Prop:** Suppose  $f : \mathbb{R} \to \mathbb{R}$  is differentiable. Then f is convex iff  $f(y) \ge f(x) + f'(x)(y-x) \qquad \forall x, y \in \mathbb{R}$
- Proof:

 $\Leftarrow$ : See Notes Section 3.2.

 $\Rightarrow$ : Exercise for Assignment 4.  $\Box$ 

#### **Convexity and Second Derivative**

- **Prop:** Suppose  $f : \mathbb{R} \to \mathbb{R}$  is twice-differentiable. Then f is convex iff  $f''(x) \ge 0 \forall x \in \mathbb{R}$ .
- **Proof:** See Notes Section 3.2.

## Subgradient Inequality in $\mathbb{R}^n$

- **Prop:** Suppose  $f : \mathbb{R}^n \to \mathbb{R}$  is differentiable. Then f is convex iff  $f(y) \ge f(x) + \nabla f(x)^T(y-x) \quad \forall x, y \in \mathbb{R}$
- Proof:
  - $\Leftarrow$ : Exercise for Assignment 4.
  - $\Rightarrow$ : See Notes Section 3.2.  $\Box$

### Minimizing over a Convex Set

- Prop: Let C⊆R<sup>n</sup> be a convex set.
   Let f : R<sup>n</sup>→R be convex and differentiable.
   Then x minimizes f over C iff ∇f(x)<sup>T</sup>(z-x)≥0 ∀z∈C.
- Proof: 

   direction
   direction

Direct from subgradient inequality.

$$f(z) \geq f(x) + \nabla f(x)^{\mathsf{T}}(z-x) \geq f(x)$$

Subgradient inequality

Our hypothesis

### Minimizing over a Convex Set

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   Let f : R<sup>n</sup>→R be convex and differentiable.
   Then x minimizes f over C iff ∇f(x)<sup>T</sup>(z-x)≥0 ∀z∈C.
- **Proof:**  $\Rightarrow$  direction

Let x be a minimizer, let  $z \in C$  and let y = z-x. Recall that  $\nabla f(x)^T y = f'(x;y) = \lim_{t \to 0} \frac{f(x+ty)-f(x)}{t}$ .

If limit is negative then we have f(x+ty) < f(x) for some  $t \in [0,1]$ , contradicting that x is a minimizer. So the limit is non-negative, and  $\nabla f(x)^{\mathsf{T}} y \ge 0$ .

#### Positive Semidefinite Matrices (again)

- Assume M is symmetric
- Old definition: M is PSD if  $\exists V \text{ s.t. } M = V^T V$ .
- New definition: M is PSD if  $y^TMy \ge 0 \ \forall y \in \mathbb{R}^n$ .
- Claim: Old  $\Rightarrow$  New.
- **Proof:**  $y^TMy = y^TV^TVy = ||Vy||^2 \ge 0.$
- Claim: New  $\Rightarrow$  Old.
- **Proof:** Based on spectral decomposition of M.

## **Convexity and Hessian**

- Prop: Let f: R<sup>n</sup>→R be a C<sup>2</sup>-function.
   Let H(x) denote the Hessian of f at point x.
   Then f is convex iff H(x) is PSD ∀x∈R<sup>n</sup>.
- **Proof:**  $\Rightarrow$  direction

Fix  $y \in \mathbb{R}^n$ . Consider function  $g_y(\alpha) = f(x+\alpha y)$ . Convexity of  $g_y$  follows from convexity of f. Thus  $g_y''(0) \ge 0$ . (convexity & 2<sup>nd</sup> derivative) Fact:  $g_y''(0) = y^T H(x) y$  (stated in Lecture 14) So  $y^T H(x) y \ge 0 \forall y \in \mathbb{R}^n \implies H(x)$  is PSD.

## **Convexity and Hessian**

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   Let H(x) denote the Hessian of f at point x.
   Then f is convex iff H(x) is PSD ∀x∈R<sup>n</sup>.
- **Proof:**  $\Leftarrow$  direction. Fix  $x,y \in \mathbb{R}^n$ . Define  $g : [0,1] \to \mathbb{R}$  by  $g(\alpha) = f(x+\alpha(y-x))$ . For all  $\alpha$ ,  $g''(\alpha) = (y-x)^T H(x+\alpha(y-x))(y-x) \ge 0$ . (The equality was stated in Lecture 14. The inequality holds since H is PSD.) So g is convex. (By Prop "Convexity and Second Derivative")  $f((1, \alpha)x + \alpha x) = g(\alpha) \le (1, \alpha)g(0) + \alpha g(1)$ 
  - $f((1-\alpha)x+\alpha y) = g(\alpha) \le (1-\alpha)g(0) + \alpha g(1)$  $= (1-\alpha)f(x) + \alpha f(y).$

### Hessian Example

#### • Example:

Let M be a symmetric nxn matrix. Let  $z{\in}\mathbb{R}^n.$ 

Define  $f : \mathbb{R}^n \to \mathbb{R}$  by  $f(x) = x^T M x - x^T z$ .

Note  $f(x) = \Sigma_i \Sigma_j M_{i,j} x_i x_j - \Sigma_i x_i z_i$ .

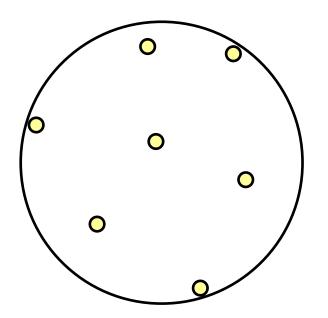
Taking partial derivatives  $\nabla f(x)_i = 2 \sum_j M_{i,j} x_j - z_i$ . So  $\nabla f(x) = 2Mx - z$ .

Recall the Hessian at x is  $H(x) = \nabla(\nabla f(x))$ .

So H(x) = 2M.

# Smallest Ball Problem

- Let {p<sub>1</sub>,...,p<sub>n</sub>} be points in R<sup>d</sup>.
   Find (unique!) ball of smallest volume (not an ellipsoid!) that contains all the p<sub>i</sub>'s.
- In other words, we want to solve:  $\min \ \{ r : \exists y \in \mathbb{R}^d \text{ s.t. } p_i \in B(y,r) \ \forall i \ \}$



# Smallest Ball Problem

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- In other words, we want to solve:  $\min \{ r : \exists y \in \mathbb{R}^d \text{ s.t. } p_i \in B(y,r) \forall i \}$
- We will formulate this as a convex program.
- In fact, our convex program will be of the form min { f(x) : Ax=b, x≥0 }, where f is convex. Minimizing a convex function over an (equality form) polyhedron
- To solve this, we will need optimality conditions for convex programs.

## (Mini)-KKT Theorem

**Theorem:** Let  $f: \mathbb{R}^n \to \mathbb{R}$  be a convex,  $C^2$  function. Let  $x \in \mathbb{R}^n$  be a feasible solution to the convex program min {  $f(x) : Ax=b, x \ge 0$  } Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t. 1)  $\nabla f(x)^T + y^T A \ge 0$ , 2) For all j, if  $x_j > 0$  then  $\nabla f(x)_j + y^T A_j = 0$ .

 Proven by Karush in 1939 (his Master's thesis!), and by Kuhn and Tucker in 1951.

# (Mini)-KKT Theorem

**Theorem:** Let  $f:\mathbb{R}^n \to \mathbb{R}$  be a convex, C<sup>2</sup> function. Let  $x \in \mathbb{R}^n$  be a feasible solution to the convex program min {  $f(x) : Ax=b, x \ge 0$  } Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t. 1)  $\nabla f(x)^T + y^T A \ge 0$ , 2) For all j, if  $x_i > 0$  then  $\nabla f(x)_i + y^T A_i = 0$ .

#### Special Case: (Strong LP Duality)

Let  $f(x) = -c^T x$ . (So the convex program is max {  $c^T x : Ax=b, x \ge 0$  }) Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t.

1) 
$$-c^{T} + y^{T}A \ge 0$$
,

2) For all j, if  $x_j > 0$  then  $-C_j + y^T A_j = 0$ .

y is feasible for Dual LP

Complementary Slackness holds **Theorem:** Let  $f:\mathbb{R}^n \to \mathbb{R}$  be a convex, C<sup>2</sup> function.

Let  $x{\in}\mathbb{R}^n$  be a feasible solution to the convex program

min { f(x) : Ax=b, x ≥0 }

Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t.

1)  $\nabla f(\mathbf{x})^T + \mathbf{y}^T \mathbf{A} \ge \mathbf{0}$ ,

2) For all j, if  $x_j > 0$  then  $\nabla f(x)_j + y^T A_j = 0$ .

**Proof:**  $\Leftarrow$  direction. Suppose such a y exists. Then

 $(\nabla f(x)^T + y^T A) x = 0.$  (Just like complementary slackness) For any feasible  $z \in \mathbb{R}^n$ , we have

 $(\nabla f(\mathbf{x})^{\mathsf{T}} + \mathbf{y}^{\mathsf{T}} \mathsf{A}) \mathbf{z} \geq 0.$ 

Subtracting these, and using Ax=Az=b, we get

 $\nabla f(x)^{T}(z-x) \geq 0 \quad \forall \text{ feasible } z.$ 

So x is optimal. (By earlier proposition "Minimizing over a Convex Set")

**Theorem:** Let  $f:\mathbb{R}^n \to \mathbb{R}$  be a convex, C<sup>2</sup> function.

Let  $x \in \mathbb{R}^n$  be a feasible solution to the convex program

min { f(x) : Ax=b, x ≥0 }

Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t.

1)  $\nabla f(\mathbf{x})^{\mathsf{T}} + \mathbf{y}^{\mathsf{T}} \mathbf{A} \ge \mathbf{0}$ ,

2) For all j, if  $x_i > 0$  then  $\nabla f(x)_i + y^T A_i = 0$ .

**Proof:**  $\Rightarrow$  direction. Suppose x is optimal. Let c=- $\nabla f(x)$ . Then  $\nabla f(x)^T(z-x) \ge 0 \Rightarrow c^T z \le c^T x$  for all feasible points z.

By our earlier proposition "Minimizing over a Convex Set"

**Theorem:** Let  $f:\mathbb{R}^n \to \mathbb{R}$  be a convex, C<sup>2</sup> function.

Let  $x \in \mathbb{R}^n$  be a feasible solution to the convex program

min { f(x) : Ax=b, x ≥0 }

Then x is optimal iff  $\exists y \in \mathbb{R}^m$  s.t.

1)  $\nabla f(\mathbf{x})^{T} + \mathbf{y}^{T} \mathbf{A} \geq \mathbf{0}$ ,

2) For all j, if  $x_j > 0$  then  $\nabla f(x)_j + y^T A_j = 0$ . **Proof:**  $\Rightarrow$  direction. Suppose x is optimal. Let  $c = -\nabla f(x)$ . Then  $\nabla f(x)^T(z-x) \ge 0 \Rightarrow c^T z \le c^T x$  for all feasible points z. So x is optimal for the LP max {  $c^T x : Ax = b, x \ge 0$  }. So there is an optimal solution y to dual LP min {  $b^T y : A^T y \ge c$  }. So  $\nabla f(x)^T + y^T A = -c^T + y^T A \ge 0 \Rightarrow$  (1) holds.

Furthermore, x and y are both optimal so C.S. holds.

⇒ whenever  $x_j>0$ , the j<sup>th</sup> dual constraint is tight ⇒  $y^T A_j = c_j$  ⇒ (2) holds.

## Smallest Ball Problem

- Let  $P = \{p_1, ..., p_n\}$  be points in  $\mathbb{R}^d$ . Let Q be dxn matrix s.t.  $Q_i = p_i$ . Let  $z \in \mathbb{R}^n$  satisfy  $z_i = p_i^T p_i$ . Define  $f : \mathbb{R}^n \to \mathbb{R}$  by  $f(x) = x^T Q^T Q x - x^T z$ .
- Claim 1: f is convex.
- Consider the convex program min {  $f(x) : \sum_j x_j = 1, x \ge 0$  }.
- Claim 2: This program has an optimal solution x.
- Claim 3: Let  $p^* = Qx$  and  $r = \sqrt{-f(x)}$ . Then  $P \subset B(p^*, r)$ .
- Claim 4: B(p<sup>\*</sup>,r) is the smallest ball containing P.

- Let  $P=\{p_1,...,p_n\}$  be points in  $\mathbb{R}^d$ . Let Q be dxn matrix s.t.  $Q_i=p_i$ . Let  $z\in\mathbb{R}^n$  satisfy  $z_i=p_i^Tp_i$ . Define  $f:\mathbb{R}^n\to\mathbb{R}$  by  $f(x)=x^TQ^TQx-x^Tz$ .
- Claim 1: f is convex.
- Proof:

By our earlier example, the Hessian of f at x is  $H(x) = 2 \cdot Q^{T}Q$ .

This is positive semi-definite.

By our proposition "Convexity and Hessian", f is convex.

- Let  $P=\{p_1,...,p_n\}$  be points in  $\mathbb{R}^d$ . Let Q be dxn matrix s.t.  $Q_i=p_i$ . Let  $z\in\mathbb{R}^n$  satisfy  $z_i=p_i^Tp_i$ . Define  $f:\mathbb{R}^n\to\mathbb{R}$  by  $f(x)=x^TQ^TQx-x^Tz$ .
- Consider the convex program min { f(x) :  $\sum_j x_j = 1, x \ge 0$  }.
- Claim 2: This program has an optimal solution.
- Proof:

The objective function is continuous.

The feasible region is a bounded polyhedron, and hence compact.

By Weierstrass' Theorem, an optimal solution exists.

- Let Q be dxn matrix s.t.  $Q_i = p_i$ . Let  $z \in \mathbb{R}^n$  satisfy  $z_i = p_i^T p_i$ . Define  $f : \mathbb{R}^n \to \mathbb{R}$  by  $f(x) = x^T Q^T Q x - x^T z$ .
- Let x be an optimal solution of min {  $f(x) : \sum_j x_j = 1, x \ge 0$  } Let  $p^* = Qx$  and  $r^2 = -f(x) = \sum_j x_j p_j^T p_j - p^{*T} p^*$ .
- **Claim 3:** The ball B(p<sup>\*</sup>,r) contains P.
- **Proof:** Note that  $\nabla f(x)^T = 2x^TQ^TQ z^T$ . By KKT,  $\exists y \in \mathbb{R}$  s.t.  $2p_j^Tp^* - p_j^Tp_j + y \ge 0 \forall j$ . Furthermore, equality holds  $\forall j$  s.t.  $x_j > 0$ . So  $y = \sum_j x_j y = \sum_j x_j p_j^Tp_j - 2\sum_j x_j p_j^Tp^* = \sum_j x_j p_j^Tp_j - 2p^{*T}p^*$ . So  $y + p^{*T}p^* = \sum_j x_j p_j^Tp_j - p^{*T}p^* = -f(x) \Rightarrow r = \sqrt{y + p^{*T}p^*}$

- Let  $p^* = Qx$  and  $r^2 = -f(x) = \sum_j x_j p_j^T p_j p^{*T} p^*$ .
- **Claim 3:** The ball B(p<sup>\*</sup>,r) contains P.
- **Proof:** Note that  $\nabla f(x) = 2x^TQ^TQ z$ . By KKT,  $\exists y \in \mathbb{R}$  s.t.  $2p_j^Tp^* - p_j^Tp_j + y \ge 0 \forall j$ . Furthermore, equality holds  $\forall j$  s.t.  $x_j > 0$ .  $\Rightarrow r = \sqrt{y + p^{*T}p^*}$ 
  - It remains to show that  $B(p^*,r)$  contains P. This holds iff  $||p_j-p^*|| \le r \forall j$ . Now  $||p_j-p^*||^2 = (p_j-p^*)^T(p_j-p^*)$   $= p^{*T}p^*-2p_j^Tp^*+p_j^Tp_j$  $\le y+p^{*T}p^* = r^2 \forall j$ .

- **Claim 4:** B(p<sup>\*</sup>,r) is the smallest ball containing P.
- **Proof:** See Textbook.