

Derangements and Quantum Colourings

Chris Godsil*
Combinatorics & Optimization
University of Waterloo

October 14, 2019

1 Introduction

These notes are meant to accompany a talk “Derangements and Quantum Colourings”. They are largely based on the following sources:

1. Cameron, Newman, Montanaro, Severini, Winter. “On the quantum chromatic number of a graph”. arXiv:quant-ph/0608016. (2006)
2. Mancińska, Roberson. “Quantum homomorphisms”. arXiv:1212.1742. (2016)
3. Mancińska, Roberson. “Oddities of quantum colorings”. arXiv:1801.03542. (2018)
4. David E. Roberson (2013). Variations on a Theme: Graph Homomorphisms. UWSpace. <http://hdl.handle.net/10012/7814>

2 Colourings

A *graph homomorphism* from X to Y is a map $\varphi : V(X) \rightarrow V(Y)$ such that if u and v are adjacent in X , then $\varphi(u)$ and $\varphi(v)$ are adjacent in Y . We write

*Research supported by Natural Sciences and Engineering Council of Canada, Grant No. RGPIN-9439

$X \rightarrow Y$ to denote that there is a homomorphism from X to Y . The most familiar examples are colourings: an m -colouring of X is a homomorphism from X to K_m . For a second class of examples, any automorphism of X is a homomorphism from X to itself.

We note that if $X \rightarrow Y$, then $\chi(X) \leq \chi(Y)$.

If $\varphi : X \rightarrow Y$ is a homomorphism and $y \in V(Y)$, the set

$$\varphi^{-1} := \{x \in V(X) : \varphi(x) = y\}$$

is the *fibre* of φ at y . Since all graphs we consider are free of loops, and fibre of a homomorphism from X to Y is a coclique in X . It follows that an m -colouring of X determines a partition of $V(X)$ with exactly m cells (where each cell is a coclique).

We choose to represent a partition π by a *characteristic matrix*, this is the 01-matrix with i -th column equal to the characteristic vector of the i -th cell of π . If M is the characteristic matrix of a partition, then the columns of M sum to $\mathbf{1}$. If M is the characteristic matrix of a colouring, each column is the characteristic vector of a coclique.

A *quantum m -colouring* of a graph X is a $|V(X)| \times m$ matrix with $d \times d$ projections as entries, such that

- (a) If $u \in V(X)$, then $\sum_{i=1}^m M_{u,i} = I_d$.
- (b) If $uv \in E(X)$, then $M_{u,i}M_{v,i} = 0$ for $i = 1, \dots, m$.

You may verify that, if $d = 1$, this is just a classical m -colouring. The least integer m such that X admits a quantum m -colouring is the *quantum chromatic number* of X , denoted $\chi_q(X)$. (Note that d is allowed to vary.) We have

$$\chi_q(X) \leq \chi(X).$$

It is an interesting exercise to verify that $\chi_q(X) = 2$ if and only if $\chi(X) = 2$.

2.1 Theorem. *If X admits a quantum m -colouring, then it admits a quantum m -colouring where all projections have the same rank.* \square

If the common rank is r , then $mr = d$.

Using the previous theorem, it is not too difficult to prove that $\chi_q(K_n) = n$.

3 Orthogonality Graphs

Suppose S is a subset of an inner product space, e.g., the unit vectors in \mathbb{C}^d . The *orthogonality graph* based on S have vertex set S , and two elements of S are adjacent if they are orthogonal. If the inner product space has dimension d , then cliques of size d are orthogonal bases and the clique number of the graph is d . We use $S(d)$ to denote the orthogonality graphs formed from the unit vectors in \mathbb{C}^d . A complex vector or matrix is *flat* if all its entries have the same absolute value. We use $S^b(d)$ to denote the subgraph of $S(d)$ induced by the flat vectors. The cliques in $S^b(d)$ correspond to flat $d \times d$ matrices; the character table of an abelian group of order d is flat and unitary (so $\omega(S^b(d)) = d$).

We use $S_{\mathbb{R}}(d)$ and $S_{\mathbb{R}}^b(d)$ to denote the real analogs of $S(d)$ and $S^b(d)$ (note that the latter graph is finite). These real graphs are induced subgraphs of the complex versions.

If M is the characteristic matrix of an m -colouring of X , then two rows of M indexed by adjacent vertices are orthogonal. This the map that sends a vertex u to $e_u^T M$ is a homomorphism from X to $S(m)$. If define the *orthogonal rank* $\xi(X)$ of X to be the least integer k such that there is a homomorphism from X to $S(k)$. We have

$$\xi(X) \leq \chi(X).$$

3.1 Theorem (Gleason). *If $d \geq 3$, then $\chi(S_{\mathbb{R}}(d)) > d$.* □

The surprising thing about Gleason's theorem is that is a simple consequence of an important result clarifying the nature of quantum measurements.

3.2 Theorem. *Assume $n = 2^k$. We have $\omega(S_{\mathbb{R}}^b(n)) = n$. If $k \geq 4$, then $\chi(S_{\mathbb{R}}^b(n)) > n$; further $\chi(S_{\mathbb{R}}^b(n))$ increases exponentially with n .*

Here the claim about the clique number is equivalent to the existence of Hadamard matrices of order a power of two. The claims are the chromatic number are much deeper, in particular the final claim follows from work of Frankl and Rödl.

4 Derangements

A *derangement* is a permutation of a set with no fixed point. If D denotes the set of derangements in $\text{Sym } n$, then D is closed under inverses and does

not contain the identity, so we may use D as the connection set for the Cayley graph $X(\text{Sym } n, D)$; we denote this graph by $\mathcal{D}(n)$. We summarize some relevant properties of $\mathcal{D}(n)$.

4.1 Theorem. *We have:*

- (a) *The maximum size of a clique in $\mathcal{D}(n)$ is n ; cliques of size n correspond to $n \times n$ Latin squares.*
- (b) *The maximum size of a coclique is $(n - 1)!$; the cocliques of size $(n - 1)!$ are cosets of the stabilizer of a point.*
- (c) *The chromatic number of $\mathcal{D}(n)$ is n .*

4.2 Corollary. *For any graph X we have $\chi(X) \leq n$ if and only if $X \rightarrow \mathcal{D}(n)$. \square*

Two graphs X and Y are *homomorphically equivalent* if $X \rightarrow Y$ and $Y \rightarrow X$. The previous corollary may restated as the statement that $\mathcal{D}(n)$ and K_n are homomorphically equivalent.

Let us represent elements of $\text{Sym } n$ by permutation matrices. The space of $n \times n$ complex matrices is an inner product space, with inner product

$$\langle M, N \rangle := \text{tr}(M^* N).$$

If M is a permutation matrix, $M^* = M^{-1}$ and we see that if M and N are permutation matrices and $M^{-1}N$ represents a derangement, then $\langle M, N \rangle = 0$. Thus $\mathcal{D}(n)$ is an orthogonality graph.

5 Rank-1 Quantum Colourings and Unitary Derangements

Suppose M defines a quantum m -colouring of X , where the entries of M have rank one. Then the entries of M must be of order $m \times m$. If

$$P_1, \dots, P_m$$

are the projections in row i of M , then there are unit vectors x_1, \dots, x_m such that

$$P_r = x_r x_r^*$$

Since $P_r P_s = 0$ if $r \neq s$, the vectors x_1, \dots, x_m are pairwise orthogonal, and therefore they form the columns of a unitary matrix, R say. If S is the unitary matrix corresponding to row j of M , then the condition $M_{i,r} M_{j,r} = 0$ holds for each r if and only if the diagonal entries of $R^* S$ are all zero. Since R and S are unitary, so is $R^* S$.

We define a *unitary derangement* to be a unitary matrix with all diagonal entries zero. Any permutation matrix is unitary, and it is a unitary derangement if and only if the permutation it represents is a derangement. The inverse of a unitary derangement is its conjugate-transpose, and so it is again a unitary derangement. Hence we may define a Cayley graph $\mathcal{UD}(n)$ on the unitary group $U(d)$, with connection set the set of unitary derangements. Note that the derangement graph $\mathcal{D}(n)$ is an induced subgraph of $\mathcal{UD}(n)$.

5.1 Theorem. *A graph X has a rank-1 quantum n -colouring if and only if $X \rightarrow \mathcal{UD}(n)$.* \square

5.2 Lemma. *If the matrices M_1, \dots, M_n form a clique in $\mathcal{UD}(n)$, let \mathcal{M} denote the $n \times n$ matrix of projections with*

$$\mathcal{M}_{i,j} = M_i e_j (M_i e_j)^*.$$

Then \mathcal{M} is a rank-1 quantum n -colouring of K_n . \square

In the context of quantum computing, rank-1 quantum n -colouring of K_n are referred as *quantum Latin squares*. (The choice of term was **not** motivated by any analogy to derangement graphs.) If L is an $n \times n$ Latin square with entries from $\{1, \dots, n\}$ we can convert L to a quantum Latin square: if $L_{i,j} = r$, replace the entry r by the projection $e_e e_r^T$.

If z is a unit vector in \mathbb{C}^n , then the unitary matrices M with i -th row equal to z form a coclique, for if $M e_i = N e_i = z$ then

$$(M^* N)_{i,i} = e_i^T M^* N e_i = z z^* \neq 0$$

and $M^* N$ is not a derangement.

You may find it interesting to prove that $\chi_q^{(1)(X)} = 3$ if and only if $\chi(X) = 3$.

6 Three Homomorphisms

We use homomorphisms to relate some of the parameters at hand. One observation is in order.

6.1 Lemma. *If W is a flat unitary matrix and D_1 and D_2 are diagonal matrices (all of the same order), then*

$$\langle D_1, W^* D_2 W \rangle = \text{tr}(D_1) \text{tr}(D_2). \quad \square$$

6.2 Theorem. *We have homomorphisms as follows:*

$$K_n \rightarrow S^{\flat}(n) \rightarrow \mathcal{UD}(n) \rightarrow S(n).$$

Proof. The n -cliques in $S^{\flat}(n)$ are exactly the flat unitary matrices of order $n \times n$. This takes care of the first homomorphism.

For the second, if $z \in \mathbb{C}^n$, let D_z be the diagonal matrix with

$$(D_z)_{i,i} = z_i.$$

If $z \in S^{\flat}(n)$, then D_z is unitary and the map

$$z \mapsto D_z W$$

takes elements of $S^{\flat}(n)$ to unitary matrices. Consider the matrix

$$Q = (D_y W)^* D_z W.$$

We have

$$Q_{i,i} = \text{tr}(e_i e_i^T Q) = \langle e_i e_i^T, (D_y W)^* D_z W \rangle = \langle e_i e_i^T, W^* D_y^* D_z W \rangle$$

and, applying the lemma (with $D_1 = e_i e_i^T$), we deduce that

$$\langle e_i e_i^T, W^* D_y^* D_z W \rangle = \text{tr}(W^* D_y^* D_z W) = \text{tr}(D_y^* D_z) = \langle y, z \rangle.$$

Accordingly if y and z are orthogonal, then Q is a derangement.

The third homomorphism is again simple. As

$$\langle M e_1, N e_1 \rangle = (M^* N)_{1,1}$$

we may use the map $M \mapsto M e_1$ as the homomorphism. \square

6.3 Corollary. *For any graph X ,*

$$\chi(X) \geq \xi^{\flat}(X) \geq \chi_q^{(1)}(X) \geq \xi(X). \quad \square$$

7 Separating χ and χ_q

Let q be an odd prime power. The vertices of the Erdős-Rényi graph $ER(q)$ are the 1-dimensional subspace of the 3-dimensional vectors over $GF(q)$; two subspaces spanned by nonzero vectors x and y are adjacent if $x^T y = 0$. (Note: this is not an Erdős-Rényi random graph.) We see that $ER(q)$ has $q^2 + q + 1$ vertices and each vertex has $q + 1$ neighbours but, unfortunately perhaps, there are $q + 1$ vertices with loops on them.

The graph we use is $ER(3)$, on 13 vertices. Each vertex is represented by a vector length three with entries 0, 1 and -1 . We normalize the vectors by assuming that the first non-zero entry is 1. Now we view these vectors as vectors over \mathbb{R} , and work with the orthogonality graph on these vectors. Denote it by Y . Clearly $\xi(Y) \leq 3$.

Cameron et al. prove the following, using properties of the quaternions.

7.1 Lemma. *There is a homomorphism from $S_{\mathbb{R}}(4)$ into the subgraph of $\mathcal{UD}(4)$ induced by the real orthogonal matrices.* \square

7.2 Corollary. *If $\xi_{\mathbb{R}}(x) \leq 4$, then $\chi_q^{(1)}(X) \leq 4$.* \square

A direct computation shows that $\chi(Y) = 4$. Consider the cone \hat{Y} over Y . Here $\xi_{\mathbb{R}}(\hat{Y}) \leq 4$, whence $\chi_q(\hat{Y}) \leq 4$. However $\chi(\hat{Y})$ must be five. Thus we have established that χ and χ_q can differ and, also, that a graph and its cone may have the same quantum chromatic number. We have not ruled out the possibility that $\chi_q(Y) = 3$, this is done in the oddities paper.

8 Derangements of Index k

We have seen that rank-1 colourings give rise to unitary derangements. What of rank- k colourings?

A $d \times d$ projection P of rank k can be written as $P = UU^*$, where U is $d \times k$ and its columns are an orthonormal basis for $\text{im}(P)$. So $U^*U = I_k$. If the matrix M represents a rank- k quantum m -colouring of X , there are $d \times k$ matrices $U_{a,i}$ (for $a \in V(X)$ and $i = 1 \dots, m$) such that

$$U_{a,i}^* U_{a,i} = I_k, \quad U_{a,i} U_{a,i}^* = M_{a,i}.$$

We see that if $i \neq j$, then $U_{a,i}^* U_{a,j} = 0$ and if $ab \in E(X)$, then $U_{a,i}^* U_{b,i} = 0$. Let \mathcal{U} be the matrix with ai -entry equal to $U_{a,i}$. Since $mk = d$, each row of \mathcal{U} is a $d \times d$ unitary matrix. If $ab \in E(X)$, then

$$(U_{a,1} \ \dots \ U_{a,m})^* (U_{b,1} \ \dots \ U_{b,m})$$

is a unitary matrix of order $mk \times mk$ with k diagonal blocks of zeros.

We define a unitary matrix M to be a *unitary derangement of index k* if it has order $mk \times mk$ and

$$M \circ (I_m \otimes J_k) = 0.$$

(If $k = 1$ we recover our previous derangements.) We can apply this term to permutation matrices, since they are unitary, and we will refer to them simply as *derangements of index k* . Since the set of $mk \times mk$ unitary derangements with index k is closed under conjugate transpose and does not contain the identity, we can use it as the connection set for a Cayley graph for the full unitary group; if $n = km$, we denote it by $\mathcal{UD}_k(n)$.

8.1 Theorem. *Let $n = mk$. A graph X has a rank- k quantum m -colouring if and only if there is a homomorphism $X \rightarrow \mathcal{UD}_k(n)$. \square*

If M is a unitary derangement (of index one) and Q is unitary of order $k \times k$, then $M \otimes Q$ is a unitary derangement of index k .

9 Grassmann Graphs

The *Grassmann graph* $Gr(d, k)$ is the graph with the k -dimensional subspaces of \mathbb{C}^d as vertices, with two subspaces adjacent if they are orthogonal. We may, and will, choose to represent the vertices of $Gr(d, k)$ by $d \times d$ projections of rank k . If P and Q are two such projections, then

$$\|P - Q\|^2 = \langle P - Q, P - Q \rangle = \text{tr}(P + Q - PQ - QP) = 2k - 2\langle P, Q \rangle.$$

Hence P and Q are at maximum distance if and only if they are orthogonal. (Since P and Q are positive semidefinite, $\langle P, Q \rangle \geq 0$.) Consequently we may view $Gr(d, k)$ as an analog of the Kneser graph $K_{d,k}$. Since the fractional chromatic number of a graph is determined by homomorphisms into Kneser graphs this suggests, correctly, that homomorphisms to Grassmann graphs will provide a quantum analog to fractional chromatic number.

9.1 Theorem. *There is a homomorphism $\mathcal{UD}_k(mk) \rightarrow Gr(mk, k)$.*

Proof. Define

$$D = \begin{pmatrix} I_k & 0 \\ 0 & 0 \end{pmatrix}.$$

If M is a $md \times md$ unitary matrix, then

$$MDM^*$$

represents orthogonal projection onto the column space of MD , i.e., onto the span of the first k columns of M .

10 Quantum Homomorphisms

Since colourings can be usefully viewed as a special case of graph homomorphisms, it seems natural to look for quantum homomorphisms.

The definition is an extension of the definition of a quantum colouring, as it should be. A *quantum homomorphism* from X to Y is a $|V(X)| \times |V(Y)|$ matrix with entries $d \times d$ projections, such that:

- (a) For each vertex a of X , we have $\sum_{y \in V(Y)} M_{a,y} = I_d$.
- (b) If a and b are adjacent vertices in X and y and z are vertices in Y that are **not** adjacent, then $M_{a,y}M_{b,z} = 0$.

You might check that, if $d = 1$, we recover the usual definition of a graph homomorphism.

It is important to note that a quantum homomorphism from X to Y is **not** a function from $V(X)$ to $V(Y)$. One symptom of this issue is that it is not obvious how we might compose quantum homomorphisms. Before discussing this, we offer a second definition of quantum homomorphism.

A *measurement on Y* is an assignment of a $d \times d$ projection to each vertex in Y , such that if P_u denote the projection indexed by u in $V()$, then

$$\sum_{u \in V(Y)} P_u = I_d.$$

Given this condition, projections associated to different vertices are orthogonal. For a physicist a measurement indexed by Y would be referred to

as a projective measurement, where the outcome of an actual measurement would be a vertex of Y . Two measurements $\{P_u\}_{u \in V(Y)}$ and $\{Q_u\}_{u \in V(Y)}$ are *compatible* if when u and v are vertices in Y that are not adjacent,

$$P_u Q_v = 0.$$

(So $P_u Q_u = 0$ for each vertex u .) The vertices of the measurement graph $\mathcal{M}_d(Y)$ are the measurements on Y , where two measurements are adjacent if they are compatible. We have immediately:

10.1 Lemma. *There is a quantum homomorphism $X \xrightarrow{q} Y$ if and only if $X \rightarrow \mathcal{M}_d(Y)$ for some d .* □

We turn to quantum composition. Assume M and N respectively represents quantum homomorphisms from X to Y and Y to Z . We define $M \star N$ to be the matrix with rows indexed by $V(X)$, columns indexed by $V(Z)$ and with

$$(M \star N)_{x,z} = \sum_{y \in V(Y)} M_{x,y} \otimes N_{y,z}, \quad (x \in V(X), z \in V(Z)).$$

You need to verify that $M \star N$ represents a quantum homomorphism from X to Z , and that this product is associative. (It then follows that graphs and quantum homomorphisms are the objects and arrows of a category.)