

# THE MINIMALLY NON-IDEAL BINARY CLUTTERS WITH A TRIANGLE

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ABSTRACT. It is proved that the lines of the Fano plane and the odd circuits of  $K_5$  constitute the only minimally non-ideal binary clutters that have a triangle.

## 1. INTRODUCTION

A clutter  $\mathbb{F}$  over a finite ground set  $E(\mathbb{F})$  is a family of subsets of  $E(\mathbb{F})$  where no subset is contained in another one. We say that  $R \subseteq E(\mathbb{F})$  is a *cover* of  $\mathbb{F}$  if, for all  $S \in \mathbb{F}$ ,  $S \cap R \neq \emptyset$ . The *blocker*  $b(\mathbb{F})$  of  $\mathbb{F}$  is the clutter, over the same ground set, of all (inclusion-wise) minimal covers of  $\mathbb{F}$ . It is well known that for any clutter  $\mathbb{F}$ ,  $b(b(\mathbb{F})) = \mathbb{F}$  [4, 8]. We say that  $\mathbb{F}$  is *binary* if, for all  $S \in \mathbb{F}$  and  $R \in b(\mathbb{F})$ ,  $|S \cap R|$  is odd. By definition, if a clutter is binary, then so is its blocker. Take disjoint subsets  $I, J \subseteq E(\mathbb{F})$ . Then  $\mathbb{F}/I \setminus J$  denotes the clutter over ground set  $E(\mathbb{F}) - (I \cup J)$  that consists of the minimal sets in  $\{S - I : S \in \mathbb{F}, S \cap J = \emptyset\}$ .<sup>1</sup> We say that  $\mathbb{F}/I \setminus J$  is a *minor* of  $\mathbb{F}$ ; it is a *proper minor* if  $I \cup J \neq \emptyset$ . It can be readily checked that if  $\mathbb{F}$  is binary, then so are all its minors [17]. We say clutters  $\mathbb{F}_1$  and  $\mathbb{F}_2$  are *isomorphic*, and denote it by  $\mathbb{F}_1 \cong \mathbb{F}_2$ , if relabeling the ground set of  $\mathbb{F}_1$  yields  $\mathbb{F}_2$ .

Denote by  $\mathbb{L}_7$  the clutter of the lines of the Fano plane, that is,

$$\mathbb{L}_7 \cong \{ \{1, 2, 3\}, \{1, 4, 5\}, \{1, 6, 7\}, \{3, 4, 6\}, \{3, 5, 7\}, \{2, 5, 6\}, \{2, 4, 7\} \}.$$

It can be readily checked that  $\mathbb{L}_7 = b(\mathbb{L}_7)$  and that  $\mathbb{L}_7$  is binary. A *cycle* in a graph is a non-empty edge subset where every vertex is incident to an even number of the edges, and a *circuit* is a minimal cycle. Denote by  $\mathbb{O}_5$  the clutter, over ground set  $E(K_5)$ , of odd circuits of the complete graph  $K_5$ . The two clutters  $\mathbb{O}_5, b(\mathbb{O}_5)$  are also binary.

A clutter  $\mathbb{F}$  is *ideal* if the polyhedron  $\{x \in \mathbb{R}_+^{E(\mathbb{F})} : x(S) \geq 1 \forall S \in \mathbb{F}\}$  has only integral extreme points. If a clutter is ideal, then so are all its minors [18]. A clutter is *minimally non-ideal (mni)* if it is not ideal and every proper minor is ideal. For instance, the three clutters  $\mathbb{L}_7, \mathbb{O}_5$  and  $b(\mathbb{O}_5)$  are mni. Notice that every non-ideal clutter has an mni clutter as a minor. Seymour ([18], page 200) proposed in 1977 the following conjecture:

**The  $f$ -flowing conjecture.**  $\mathbb{L}_7, \mathbb{O}_5$  and  $b(\mathbb{O}_5)$  are the only mni binary clutters.

A *triangle* in clutter  $\mathbb{F}$  is a set  $S \in \mathbb{F}$  such that  $|S| = 3$ . Observe that both  $\mathbb{L}_7$  and  $\mathbb{O}_5$  have triangles. As  $b(b(\mathbb{O}_5)) = \mathbb{O}_5$ , the  $f$ -flowing conjecture implies immediately the following:

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<sup>1</sup>Given sets  $A, B$  we denote by  $A - B$  the set  $\{a \in A : a \notin B\}$  and, for element  $a$ , we write  $A - a$  instead of  $A - \{a\}$ .

**The weak  $f$ -flowing conjecture.** If  $\mathbb{F}$  is an mni binary clutter, then  $\mathbb{F}$  or  $b(\mathbb{F})$  has a triangle.

The following is the main result of the paper:

**Theorem 1.**  $\mathbb{L}_7$  and  $\mathbb{O}_5$  are the only mni binary clutters that have a triangle.

In other words we prove that the *weak  $f$ -flowing conjecture* implies the  *$f$ -flowing conjecture*.

**1.1. Review of existing results.** All matroids considered in this paper are binary, and a basic knowledge of these matroids is assumed; we follow for the most part the notation used in Oxley [12] (second edition). Take a matroid  $M$  over ground set  $E(M)$ . Recall that a circuit is a minimal dependent set of  $M$  and a cocircuit is a minimal dependent set of the dual  $M^*$ . A *cycle* is the symmetric difference of circuits, and a *cocycle* is the symmetric difference of cocircuits. It is well-known that a nonempty cycle is a disjoint union of circuits ([12], Theorem 9.1.2). Let  $\Sigma \subseteq E(M)$ . The pair  $(M, \Sigma)$  is called a *signed matroid*. A subset  $\Gamma \subseteq E(M)$  is a *signature* of  $(M, \Sigma)$  if  $\Sigma \Delta \Gamma$  is a cocycle of  $M$ . Observe that the symmetric difference of an odd number of signatures is another signature. For a signature  $\Gamma$ , the operation of replacing  $(M, \Sigma)$  by  $(M, \Gamma)$  is called *resigning*. A subset  $S \subseteq E(M)$  is said to be *odd* (resp. *even*) if  $|S \cap \Sigma|$  is odd (resp. even). An element  $f \in E(M)$  is *odd* (resp. *even*) if  $\{f\}$  is odd (resp. even). Observe that resigning a signed matroid preserves the parity of every cycle. Signed matroids are key objects as they represent binary clutters:

**Proposition 2** ([8, 11], also see [3, 6]). *A clutter  $\mathbb{F}$  is binary if, and only if, the sets of  $\mathbb{F}$  are the odd circuits of a signed matroid. Moreover, assuming  $\mathbb{F}$  is the clutter of odd circuits of a signed matroid, then  $b(\mathbb{F})$  is precisely the set of minimal signatures.*

Take disjoint subsets  $I, J \subseteq E(M)$ . If  $I$  contains an odd circuit, we define  $(M, \Sigma)/I \setminus J := (M/I \setminus J, \emptyset)$ . Otherwise, by Proposition 2, there is a signature  $\Sigma'$  that is disjoint from  $I$ , and we define  $(M, \Sigma)/I \setminus J := (M/I \setminus J, \Sigma' - J)$ . We call  $(M, \Sigma)/I \setminus J$  a *minor* of  $(M, \Sigma)$ . Notice that minors are defined only up to resigning. We have the following relation between minors of binary clutters and minors of signed matroids:

**Remark 3** (see [3]). *Let  $\mathbb{F}$  be a binary clutter represented as the signed matroid  $(M, \Sigma)$ . Take disjoint subsets  $I, J \subseteq E(\mathbb{F})$ . Then  $\mathbb{F}/I \setminus J$  is represented as the signed matroid  $(M, \Sigma)/I \setminus J$ .*

We denote by  $F_7$  the Fano matroid. The lines of the Fano plane  $\mathbb{L}_7$  is represented as the signed matroid  $(F_7, E(F_7))$ . The following is an implicit result of Cornuéjols and Guenin [3], pages 349-350 (that we do not rely on):

**Theorem 4.** *Let  $\mathbb{F}$  be an mni binary clutter represented as the signed matroid  $(M, \Sigma)$ . If  $M$  has no  $F_7$  minor, then  $\mathbb{F} \cong \mathbb{O}_5$  or  $\mathbb{F} \cong b(\mathbb{O}_5)$ .*

The proof of this theorem relies on a connectivity result. To explain this result, let  $M$  be a matroid whose rank function is  $r : 2^{E(M)} \rightarrow \{0, 1, 2, \dots\}$ . The *connectivity function*  $\lambda_M : 2^{E(M)} \rightarrow \{0, 1, 2, \dots\}$  is defined, for each  $X \subseteq E(M)$ , as  $\lambda_M(X) := r(X) + r(\bar{X}) - r(E(M))$ .<sup>2</sup> Take an integer  $k \geq 1$ . We say that

<sup>2</sup> $\bar{X} := E(M) - X$

$X \subseteq E(M)$  is  $k$ -separating if  $\lambda_M(X) \leq k - 1$ . A  $k$ -separation is a pair  $(X, \bar{X})$ , where  $X$  is  $k$ -separating and  $\min\{|X|, |\bar{X}|\} \geq k$ . We say  $M$  is  $(k + 1)$ -connected if, for each  $r \in [k]$ ,  $M$  has no  $r$ -separation.<sup>3</sup> A matroid is *internally 4-connected* if it is 3-connected, and for every 3-separation  $(X, \bar{X})$ , either  $|X| = 3$  or  $|\bar{X}| = 3$ . A key step in the proof of Theorem 4 is the following tool, that will also be essential for us:

**Theorem 5** (Cornuéjols and Guenin [3], Remark 5.3, Propositions 6.1 and 7.1). *Let  $\mathbb{F}$  be an mni binary clutter represented as the signed matroid  $(M, \Sigma)$ . Then  $M$  is internally 4-connected.*

For a graph  $G$ ,  $\text{cycle}(G)$  will denote the cycle matroid of  $G$ , i.e. the matroid whose cycles are exactly the cycles of the graph  $G$ . Then by definition,  $\mathbb{O}_5$  is represented as the signed matroid  $(\text{cycle}(K_5), E(K_5))$ . Another essential tool for our work is the following result:

**Theorem 6** (Guenin [5], also see Schrijver [14]). *Let  $\mathbb{F}$  be an mni binary clutter represented as the signed matroid  $(M, \Sigma)$ . If  $M$  is graphic, then  $\mathbb{F} \cong \mathbb{O}_5$ .*

The proof of Theorem 4 relies on Theorems 5 and 6, as well as Seymour's characterization of regular matroids [15]. It seems very hard to extend Theorem 4 directly to the general case of mni binary clutters, as the existence of an  $F_7$  minor in  $M$  is not sufficient to imply the existence of an  $(F_7, E(F_7))$  minor in  $(M, \Sigma)$ . We will, however, make use of Theorems 5 and 6 in the proof of our main result, Theorem 1. In the next section we sketch the proof, and give an outline of the remainder of the paper.

## 2. A PROOF SKETCH OF THE MAIN RESULT

Take a matroid  $M$ . For  $R \subseteq E(M)$ , we write  $M|R := M \setminus (E(M) - R)$ . We say that  $\{e_1, \dots, e_6\} \subseteq E(M)$  is an *induced  $K_4$*  of  $M$  if  $M|\{e_1, \dots, e_6\}$  is isomorphic to  $\text{cycle}(K_4)$ .

**2.1. The clutter theoretic part.** Let  $\mathbb{F}$  be an mni binary clutter with a triangle. By Proposition 2,  $\mathbb{F}$  can be represented as a signed matroid. By using a seminal result of Lehman on mni clutters, we find a suitable representation of  $\mathbb{F}$  to work with:

**Theorem 7.** *Let  $\mathbb{F}$  be an mni binary clutter with a triangle. Then  $\mathbb{F}$  is the clutter of odd circuits of a signed matroid  $(M, E(M))$  where the following statements hold:*

- a.  $M$  is internally 4-connected (Theorem 5),
- b. every element in  $E(M)$  is contained in exactly three triangles of  $M$ ,
- c. if  $|E(M)| \leq 12$ , then  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ ,
- d. if  $M$  is graphic, then  $\mathbb{F} \cong \mathbb{O}_5$  (Theorem 6),
- e. if  $M$  has an induced  $K_4$ , then  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ .

The proof of this theorem is provided in §3. After this point, we abandon the clutter  $\mathbb{F}$  and work solely with the signed matroid  $(M, E(M))$ . We will show the following:

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<sup>3</sup> $[k] := \{1, \dots, k\}$

**Theorem 8.** *Let  $M$  be an internally 4-connected binary matroid where every element is contained in exactly three triangles. Then one of the following statements holds:*

- i.  $|E(M)| \leq 11$ ,
- ii.  $M$  is graphic,
- iii.  $M$  has an induced  $K_4$ , or
- iv. the signed matroid  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor.

We will sketch the proof of this theorem shortly. Notice however that Theorem 1 is an immediate consequence of these two results:

*Proof of Theorem 1.* Let  $\mathbb{F}$  be an mni binary clutter with a triangle. By Theorem 7,  $\mathbb{F}$  is represented as a signed matroid  $(M, E(M))$ , where  $M$  is an internally 4-connected matroid and every element is contained in exactly three triangles. If either  $|E(M)| \leq 12$ ,  $M$  is graphic, or  $M$  has an induced  $K_4$ , then by Theorem 7 (c)-(e),  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ . Otherwise, by Theorem 8, the signed matroid  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor, so the mni  $\mathbb{F}$  has the non-ideal  $\mathbb{L}_7$  as a minor by Remark 3, implying in turn that  $\mathbb{F} \cong \mathbb{L}_7$ , thereby finishing the proof.  $\square$

**2.2. The matroid theoretic part.** We now sketch the proof of Theorem 8. Let  $M$  be a matroid where the following assumptions hold:

### Common hypotheses

- (h1)  $M$  is an internally 4-connected matroid,
- (h2) every element in  $E(M)$  is contained in exactly three triangles of  $M$ .

Since  $M$  is internally 4-connected,  $M$  is a simple (and cosimple) matroid. In particular, the three triangles containing an element are otherwise pairwise disjoint. Take an element  $\Omega \in E(M)$ . Denote the three triangles of  $M$  containing  $\Omega$  by  $\{\Omega, f, f'\}$ ,  $\{\Omega, g, g'\}$ ,  $\{\Omega, h, h'\}$ . Since  $M$  is simple,  $M/\Omega$  does not have a loop, and  $\{f, f'\}$ ,  $\{g, g'\}$ ,  $\{h, h'\}$  are the non-trivial parallel classes of  $M/\Omega$ . It follows that the simplification  $\text{si}(M/\Omega)$  is obtained from  $M/\Omega$  by deleting one element from each one of  $\{f, f'\}$ ,  $\{g, g'\}$ ,  $\{h, h'\}$ . If  $f, g, h$  are the elements left in  $\text{si}(M/\Omega)$ , we write  $\Lambda(\Omega) := \{f, g, h\}$ .

The proof of Theorem 8 relies on the following four propositions, as well as two theorems not done by us:

**Proposition 9.** *Suppose (h1)-(h2) hold and let  $\Omega \in E(M)$ . If  $\Lambda(\Omega)$  is a cocycle of  $\text{si}(M/\Omega)$ , then  $M$  has an induced  $K_4$ .*

*Proof.* Suppose that  $\Lambda(\Omega)$  is a cocycle of  $\text{si}(M/e)$ . Denote by  $\{\Omega, f, f'\}$ ,  $\{\Omega, g, g'\}$ ,  $\{\Omega, h, h'\}$  the triangles of  $M$  containing  $\Omega$  where  $\Lambda(\Omega) = \{f, g, h\}$ . Since  $\{f, g, h\}$  is a cocycle of  $\text{si}(M/\Omega)$ ,  $D := \{f, f', g, g', h, h'\}$  is a cocycle of  $M/\Omega$  and hence of  $M$ . As  $f$  is in three triangles of  $M$ , it is contained in a triangle  $C$  that is different from  $\{\Omega, f, f'\}$ . For  $D$  is a cocycle,  $|C \cap D|$  is even, and because  $f \in C \cap D$ ,  $|C \cap D| = 2$ . Moreover,  $C \cap D \neq \{f, f'\}$ , for otherwise  $C \Delta \{f, f', \Omega\}$  would be a cycle of cardinality two, which cannot be

the case as  $M$  is simple. Hence, we may assume that  $C \cap D = \{f, g\}$  or  $C \cap D = \{f, g'\}$ . In either cases,  $C \cup \{\Omega, f, f', g, g'\}$  is an induced  $K_4$  of  $M$ , as required.  $\square$

We require a few preliminaries to prove the next proposition. A *graft* is a pair  $(G, T)$ , where  $G$  is a graph and  $T \subseteq V(G)$  is of even cardinality. Vertices in  $T$  are called *terminals*. Take a subset  $J \subseteq E(G)$ . Denote by  $\text{odd}(J) \subseteq V(G)$  the vertices incident with an odd number of non-loop edges in  $J$ . If  $\text{odd}(J) = T$ , then we call  $J$  a  $T$ -*join*. Start with the vertex-edge incidence matrix of  $G$ , and add the vertex-incidence vector of  $T$  as a column; call this matrix  $A$ . Let  $M$  be the (binary) matroid whose binary representation is  $A$ , and denote by  $t$  the element of  $M$  corresponding to column  $T$ . Then  $C \subseteq E(M)$  is a cycle of  $M$  if, and only if, one of the following holds:

- $t \notin C$  and  $C$  is a cycle of  $G$ ,
- $t \in C$  and  $C - t$  is a  $T$ -join of  $G$ .

We call  $M$  the *graft matroid* of  $(G, T)$ . By convention,  $t$  will always be the element of  $M$  corresponding to the terminals  $T$ . Notice that if  $|T| \leq 2$ , then the graft matroid of  $(G, T)$  is graphic. The next folklore remark states that graft matroids are precisely those matroids that are one deletion away from being graphic (see for instance [12], Lemma 10.3.8):

**Remark 10.** *Take a binary matroid  $M$  and an element  $t \in E(M)$  such that  $M \setminus t = \text{cycle}(G)$ , for some graph  $G$ . If  $C$  is a cycle of  $M$  containing  $t$ , then  $M$  is the graft matroid of the graft  $(G, \text{odd}(C - t))$ .*

We are now ready for the next proposition:

**Proposition 11.** *Suppose (h1)-(h2) hold and let  $\Omega \in E(M)$ . If  $M \setminus \Omega$  is graphic, then  $M$  is graphic or has an induced  $K_4$ .*

*Proof.* Suppose  $M \setminus \Omega$  is graphic. By Remark 10, there is a graft  $(G, T)$  whose graft matroid is  $M$ , where  $t = \Omega$ . If  $|T| \leq 2$ , then  $M$  is graphic, so we are done. Otherwise,  $|T| \geq 4$ . Denote the three triangles of  $M$  containing  $\Omega$  by  $\{\Omega, f, f'\}$ ,  $\{\Omega, g, g'\}$ ,  $\{\Omega, h, h'\}$ . Then  $\{f, f'\}$ ,  $\{g, g'\}$  and  $\{h, h'\}$  are  $T$ -joins of  $G$ . Since  $M$  is simple, we see that  $G$  does not have parallel edges. As a result,  $|T| = 4$  and  $\{f, f', g, g', h, h'\}$  is an induced  $K_4$  of  $M$ , as required.  $\square$

**Proposition 12.** *Suppose (h1)-(h2) hold and let  $\Omega \in E(M)$ . If  $\Lambda(\Omega)$  is contained in a circuit of  $si(M/\Omega)$ , then either  $M$  has an induced  $K_4$  or  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor.*

We prove this proposition in §5.4.

**Proposition 13.** *Suppose (h1)-(h2) hold and let  $e_1, e_2, e_3, e_4$  be distinct elements of  $M$  such that, for every  $i \in [4]$ ,  $si(M/e_i)$  is internally 4-connected and is the cycle matroid of a graph where the three edges of  $\Lambda(e_i)$  are incident to the same vertex. Then either  $|E(M)| \leq 11$ , or there is an element  $\Omega \in E(M)$  such that  $M \setminus \Omega$  is graphic.*

This proposition is proved in §7. We will also need the following result of Seymour [19] that characterizes, under appropriate connectivity conditions, when three distinct elements of a matroid are contained in a circuit:

**Theorem 14.** *Let  $M$  be an internally 4-connected binary matroid, and let  $f, g, h$  be distinct elements. Then one of the following statements holds:*

- a.  $\{f, g, h\}$  is contained in a circuit of  $M$ ,
- b.  $\{f, g, h\}$  is a cocycle of  $M$ , or
- c.  $M$  is the cycle matroid of a graph where edges  $f, g, h$  are incident to the same vertex.

The following result of Chun and Oxley [2] on internally 4-connected matroids is the last needed ingredient:

**Theorem 15.** *Let  $M$  be an internally 4-connected binary matroid where every element is in exactly three triangles. Then there exist distinct elements  $e_1, e_2, e_3, e_4 \in E(M)$  such that, for each  $j \in [4]$ ,  $\text{si}(M/e_j)$  is internally 4-connected.*

We are now ready to prove Theorem 8:

*Proof of Theorem 8.* Suppose (h1)-(h2) hold. By Theorem 15, there exist distinct elements  $e_1, e_2, e_3, e_4$  of  $M$  such that, for each  $j \in [4]$ ,  $\text{si}(M/e_j)$  is internally 4-connected. For  $j \in [4]$ ,

- if  $\Lambda(e_j)$  is contained in a circuit of  $\text{si}(M/e_j)$ , then by Proposition 12, either  $M$  has an induced  $K_4$  and so (iii) holds, or  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor and so (iv) holds,
- if  $\Lambda(e_j)$  is a cocycle of  $\text{si}(M/e_j)$ , then by Proposition 9,  $M$  has an induced  $K_4$ , so (iii) holds.

Otherwise, it follows from Theorem 14 that, for each  $j \in [4]$ ,  $\text{si}(M/e_j)$  is the cycle matroid of a graph where the three edges in  $\Lambda(e_j)$  are incident to the same vertex. By Proposition 13, either  $|E(M)| \leq 11$  and so (i) holds, or there is an element  $\Omega \in E(M)$  such that  $M \setminus \Omega$  is graphic. By Proposition 11, either  $M$  is graphic and so (ii) holds, or  $M$  has an induced  $K_4$  and so (iii) holds. In all cases, one of (i)-(iv) holds, and so we are done.  $\square$

**2.3. Outline of the paper.** In §3 we review Lehman's theorem on mni clutters and prove Theorem 7. In §4 we introduce *signed grafts* and present two instances that have  $(F_7, E(F_7))$  as a minor. In §5 we leverage these results to prove Proposition 12. In §6 we introduce *even cycle matroids* and prove several relevant results, which in turn lead to a proof of Proposition 13 in §7.

### 3. LEHMAN AND THEOREM 7

Let  $\mathbb{F}$  be a clutter. We denote by  $M(\mathbb{F})$  the 0, 1 matrix whose columns are indexed by  $E(\mathbb{F})$  and whose rows are the incidence vectors of the sets in  $\mathbb{F}$ . The clutter of the minimum cardinality sets in  $\mathbb{F}$  is denoted by  $\bar{\mathbb{F}}$ . For an integer  $k \geq 1$ , a 0, 1 matrix is *k-regular* if each row and each column has exactly  $k$  ones. Lehman [9] (see Seymour [16]) proved a structural result on mni clutters; we only need the binary version of his result:

**Theorem 16.** *Let  $\mathbb{F}$  be an mni binary clutter. Then  $\mathbb{K} := b(\mathbb{F})$  is also mni, and the following statements hold:*

- a.  $M(\bar{\mathbb{F}})$  and  $M(\bar{\mathbb{K}})$  are square matrices,
- b. for some integers  $r \geq 3$  and  $s \geq 3$ ,  $M(\bar{\mathbb{F}})$  is  $r$ -regular and  $M(\bar{\mathbb{K}})$  is  $s$ -regular,

- c. for  $n := |E(\mathbb{F})|$ ,  $rs - n$  is an even integer such that  $2 \leq rs - n \leq \min\{r - 1, s - 1\}$ , and  
d. after possibly rearranging the rows of  $M(\overline{\mathbb{K}})$ , we have

$$M(\overline{\mathbb{F}})M(\overline{\mathbb{K}})^\top = J + (rs - n)I.$$

Here,  $J$  is the all-ones matrix and  $I$  is the identity matrix.

Notice that if  $\mathbb{F}$  is an mni binary clutter with a triangle, then  $r = 3$  and  $3s - n = 2$ . The following is therefore an easy consequence (see for instance [10]):

**Remark 17.** Let  $\mathbb{F}$  be an mni binary clutter with a triangle, and let  $\mathbb{K} := b(\mathbb{F})$ . Denote by  $s$  the minimum cardinality of a set in  $\mathbb{K}$ . Then,

- a. if  $s = 3$  then  $\mathbb{F} \cong \mathbb{L}_7$ , and  
b. if  $s = 4$  then  $\mathbb{F} \cong \mathbb{O}_5$ .

Bridges and Ryser [1] showed that the two matrices satisfying the equation in (d) commute:

**Theorem 18.** Take square 0, 1 matrices  $A, B$  such that for some integer  $d \geq 1$ ,  $AB = J + dI$ . Then  $AB = BA$ .

We are now ready to prove Theorem 7:

*Proof of Theorem 7.* Let  $\mathbb{F}$  be an mni binary clutter with a triangle, and set  $n := |E(\mathbb{F})|$ . Let  $\mathbb{K} := b(\mathbb{F})$  and denote by  $s$  the minimum cardinality of a set in  $\mathbb{K}$ . By Theorem 16, after possibly rearranging the rows of  $M(\overline{\mathbb{K}})$ ,

$$(\star) \quad r = 3 \quad \text{and} \quad s \geq 3 \quad \text{and} \quad 3s - n = 2 \quad \text{and} \quad M(\overline{\mathbb{F}})M(\overline{\mathbb{K}})^\top = J + 2I.$$

Note further that  $\overline{\mathbb{F}}$  is precisely the clutter of the triangles of  $\mathbb{F}$ , and since  $M(\overline{\mathbb{F}})$  is 3-regular, every element of  $E(\mathbb{F})$  is contained in exactly 3 triangles of  $\mathbb{F}$ . Label the rows of  $M(\overline{\mathbb{F}})$  as  $S_1, \dots, S_n \in \overline{\mathbb{F}}$ , and the rows of  $M(\overline{\mathbb{K}})$  as  $R_1, \dots, R_n \in \overline{\mathbb{K}}$ . Then the last equation implies, for all  $i, j \in [n]$ , that

$$|S_i \cap R_j| = \begin{cases} 3 & \text{if } i = j, \\ 1 & \text{if } i \neq j. \end{cases}$$

For each  $i \in [n]$ , we say that  $S_i$  and  $R_i$  are *mates* of one another. Thus, a triangle of  $\mathbb{F}$  is contained in its mate, and it intersects all the other triangle mates exactly once.

**1.** Take an element  $e \in E(\mathbb{F})$ , denote by  $S, S', S''$  the triangles of  $\mathbb{F}$  containing  $e$ , and by  $R, R', R''$  their respective mates in  $\overline{\mathbb{K}}$ . Then

- $R \cap R' = R' \cap R'' = R'' \cap R = \{e\}$ ,
- $R \cup R' \cup R'' = E(\mathbb{F})$ , and
- $S \cap S' = S' \cap S'' = S'' \cap S = \{e\}$ .

*Subproof.* It follows from  $(\star)$  and Theorem 18 that  $M(\overline{\mathbb{K}})^\top M(\overline{\mathbb{F}}) = J + 2I$ . Denote by  $c_e$  the column of  $M(\overline{\mathbb{F}})$  corresponding to  $e$ , and for each  $f \in E(\mathbb{F})$ , denote by  $c'_f$  the column of  $M(\overline{\mathbb{K}})$  corresponding to  $f$ . Then the matrix equation implies that  $c_e^\top c'_e = 3$  and, for each  $f \in E(\mathbb{F}) - e$ , that  $c_e^\top c'_f = 1$ ; the first and second lines follow. Since  $S \subseteq R, S' \subseteq R'$  and  $S'' \subseteq R''$ , the third line follows.  $\diamond$

Since  $\mathbb{F}$  is a binary clutter, we get from Proposition 2 that  $\mathbb{F}$  is the clutter of odd circuits of a signed matroid  $(M, \Sigma)$ .

**2.**  $E(M)$  is a signature of  $(M, \Sigma)$ .

*Subproof.* Take  $e \in E(\mathbb{F})$  and let  $R, R', R''$  be the mates of the triangles of  $\mathbb{F}$  containing  $e$ . Since  $R, R', R''$  belong to  $b(\mathbb{F})$ , they are signatures of  $(M, \Sigma)$  by Proposition 2. So their symmetric difference  $R \triangle R' \triangle R''$  is also a signature. However, (1) implies that  $R \triangle R' \triangle R'' = E(M)$ , so  $E(M)$  is a signature.  $\diamond$

Thus,  $\mathbb{F}$  is the clutter of odd circuits of the signed matroid  $(M, E(M))$ . It follows from Theorem 5 that  $M$  is internally 4-connected, so (a) holds.

**3.** Every element of  $E(M)$  is contained in exactly 3 triangles of  $M$ , so (b) holds.

*Subproof.* Since  $\mathbb{F}$  is the clutter of odd circuits of  $(M, E(M))$ , the triangles of  $\mathbb{F}$  are precisely the triangles of  $M$ . Since every element of  $E(\mathbb{F})$  is contained in exactly 3 triangles of  $\mathbb{F}$ , the claim follows.  $\diamond$

**4.** If  $|E(M)| \leq 12$ , then  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ , so (c) holds.

*Subproof.* By  $(\star)$ ,  $3s - 2 = n = |E(\mathbb{F})| = |E(M)| \leq 12$  and  $s \geq 3$ , so  $s \in \{3, 4\}$  and by Remark 17, we get that  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ .  $\diamond$

It follows from Theorem 6 that if  $M$  is graphic, then  $\mathbb{F} \cong \mathbb{O}_5$ , so (d) holds. It remains to prove (e). To this end, assume that  $M$  has an induced  $K_4$ , that is, there are elements  $e_1, \dots, e_6 \in E(M)$  such that  $M|_{\{e_1, \dots, e_6\}} \cong \text{cycle}(K_4)$ . As the triangles of  $M$  are precisely the triangles of  $\mathbb{F}$ , we may assume that  $S_1, S_2, S_3, S_4$  are the four triangles of  $M|_{\{e_1, \dots, e_6\}}$ .

**5.** For all distinct  $i, j \in [4]$ ,  $R_i \cap R_j \subseteq \{e_1, \dots, e_6\}$ .

*Subproof.* As  $S_i, S_j$  are distinct triangles of  $K_4$ , there is an  $e \in \{e_1, \dots, e_6\}$  such that  $S_i \cap S_j = \{e\}$ . It now follows from (1) that  $R_i \cap R_j = \{e\} \subseteq \{e_1, \dots, e_6\}$ .  $\diamond$

**6.** For all  $i \in [4]$ ,  $R_i \cap \{e_1, \dots, e_6\} = S_i$  and  $|R_i - \{e_1, \dots, e_6\}| = s - 3$ .

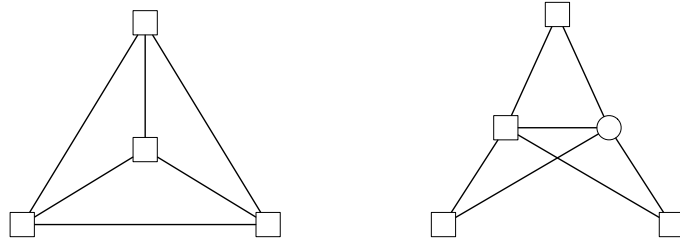
*Subproof.* Since  $R_i$  is the mate of  $S_i$ , we have  $S_i \subseteq R_i$ . As  $R_i$  intersects every other triangle exactly once, and  $|S_i \cap S_j| = 1$  for each  $j \in [4] - i$ , we get that  $R_i \cap \{e_1, \dots, e_6\} = S_i$ .  $\diamond$

Putting (5) and (6) together, we get that  $|E(M)| \geq 6 + 4(s - 3)$ . From  $(\star)$  we have that  $s \geq 3$ , and also that  $|E(M)| = |E(\mathbb{F})| = n = 3s - 2$ , so  $3s - 2 \geq 6 + 4(s - 3)$ , implying in turn that  $s \in \{3, 4\}$ . It now follows Remark 17 that  $\mathbb{F} \cong \mathbb{L}_7$  or  $\mathbb{F} \cong \mathbb{O}_5$ , thereby proving (e). This finishes the proof of Theorem 7.  $\square$

## 4. QUADRUMS AND TRIFOLDS

**4.1. Representations of the Fano matroid.** A *plain quadrum* is the graft  $(K_4, V(K_4))$ . A *plain trifold* is the graft where, the graph has vertex set  $[5]$  and edges  $\{1, 2\}, \{1, 3\}, \{1, 4\}, \{1, 5\}, \{2, 3\}, \{2, 4\}, \{2, 5\}$ , and the terminals are  $\{2, 3, 4, 5\}$ . Drawings of the plain quadrum and the plain trifold are given in Figure 1.



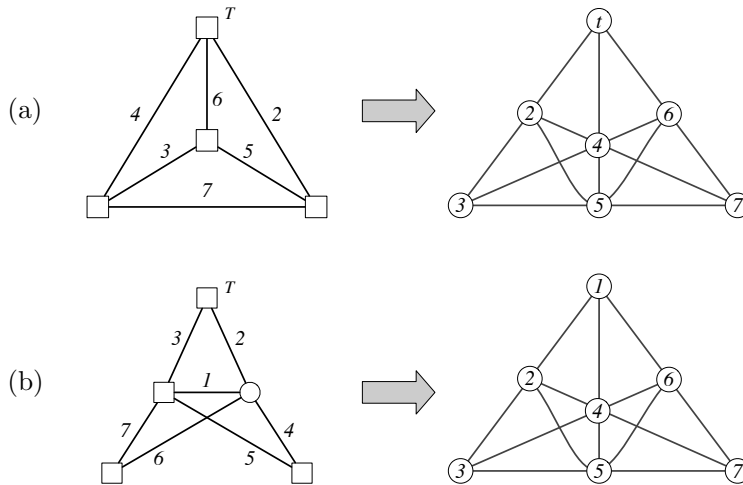


**Figure 1.** Left: plain quadrum, right: plain trifold. Square vertices are terminals.

**Remark 19.** Let  $(G, T)$  be a graft, and  $N$  its graft matroid. Then the following statements hold:

- a. if  $(G, T)$  is a plain quadrum, then  $N \cong F_7$ ,
- b. if  $(G, T)$  is a plain trifold, then  $N/t \cong F_7$ .

*Proof.* Notice that a matroid is determined by the set of its circuits. **(a)** Consider Figure 2 (a). We assign  $t$  and each edge of the plain quadrum to an element of  $F_7$ . It now suffices to observe that the circuits of  $N$  correspond to the circuits of  $F_7$ , i.e. to the lines and the line complements of the Fano plane. **(b)** Consider Figure 2 (b). We assign each edge of the plain trifold to an element of  $F_7$ . Observe that the circuits of  $N/t$ , which are the circuits and  $T$ -joins of  $G$ , correspond to the circuits of  $F_7$ . □

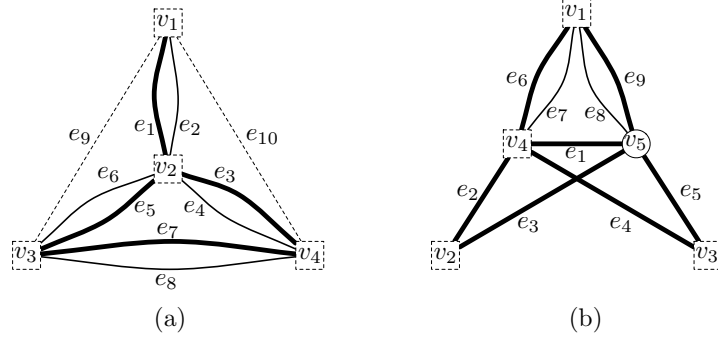


**Figure 2.** The Fano matroid in disguise

**4.2. Signed grafts: quadrum and trifolds.** A signed graft is a triple  $(G, T, \Gamma)$ , where  $(G, T)$  is a graft and  $\Gamma \subseteq E(G) \cup \{t\}$ . Note that we assign parity to each edge as well as to the set of terminals.

A quadrum is the signed graft  $(G, T, \Gamma)$  where  $(G, T)$  is a plain quadrum and  $\Gamma = E(G) \cup \{t\}$ . A super quadrum is the signed graft displayed in Figure 3 (a) which is obtained as follows: start with a plain quadrum, take a set  $S$  of four edges that contain a triangle, the element  $t$  and the two edges outside  $S$  can have either parities, and replace each edge of  $S$  by a pair of parallel edges of distinct parities.

A *trifold* is the signed graft  $(G, T, \Gamma)$  where  $(G, T)$  is a plain trifold and  $\Gamma = E(G)$ . A *super trifold* is the signed graft displayed in Figure 3 (b) which is obtained as follows: start with a plain trifold, take two triangles and the two edges  $S$  disjoint from them, the edges outside  $S$  are odd, replace each edge of  $S$  by a pair of parallel edges of distinct parities, and element  $t$  can have either parities.



**Figure 3.** (a) Super quadrum, (b) super trifold. Square vertices are terminals. Bold edges are odd. Thin edges are even. Dashed edges can be either odd or even. In (a) and (b),  $t$  can be odd or even.

4.3. **Segway to  $(F_7, E(F_7))$ .** Let  $G$  be a graph. For a vertex  $v \in V(G)$ , we denote by  $\delta_G(v)$  the set of non-loop edges of  $G$  that are incident with  $v$ . Take a signed graft  $(G, T, \Gamma)$  and a terminal  $v \in T$ . Let  $B := \delta_G(v) \cup \{t\}$ . We say that  $(G, T, \Gamma \Delta B)$  is obtained from  $(G, T, \Gamma)$  by *resigning on the terminal*  $v$ .

**Remark 20.** Let  $(G, T, \Gamma)$  be a signed graft, and  $N$  the graft matroid of the graft. If  $(G, T, \Gamma')$  is obtained from  $(G, T, \Gamma)$  by resigning on a terminal, then  $\Gamma'$  is a signature of the signed matroid  $(N, \Gamma)$ .

*Proof.* It suffices to show that, for each terminal  $v \in T$ , the set  $B := \delta_G(v) \cup \{t\}$  is a cocycle of  $N$ . To this end, let  $C$  be a cycle of  $N$ . If  $t \notin C$ , then  $C$  is a cycle of  $G$ , and so  $|C \cap B| = |C \cap \delta_G(v)|$  even. Otherwise,  $t \in C$  and  $C - t$  is a  $T$ -join of  $G$ . Since  $v \in T$ ,  $|(C - t) \cap \delta_G(v)|$  must be odd, implying in turn that  $|C \cap B|$  is even. In both cases, for every cycle  $C$  of  $N$ ,  $|C \cap B|$  is even, which means that  $B$  is a cocycle of  $N$ .  $\square$

**Proposition 21.** Let  $(G, T, \Gamma)$  be a signed graft, and  $N$  the graft matroid of the graft. If  $(G, T, \Gamma)$  is a quadrum, a super quadrum, a trifold, or a super trifold, then  $(N, \Gamma)$  has  $(F_7, E(F_7))$  as a minor.

*Proof.* **Case 1:**  $(G, T, \Gamma)$  is a quadrum. By definition,  $(G, T)$  is a plain quadrum and  $\Gamma = E(N)$ . Remark 19 states that  $N \cong F_7$ . Hence,  $(N, \Gamma) \cong (F_7, E(F_7))$ . **Case 2:**  $(G, T, \Gamma)$  is a super quadrum. Label the vertices and edges of  $G$  as in Figure 3 (a). After possibly resigning on terminals  $v_2, v_3, v_4$ , we may assume by Remark 20 that  $\Gamma = \{e_1, e_3, e_5, e_7, e_9, e_{10}, t\}$ . Let  $(N', \Gamma) := (N, \Gamma) \setminus \{e_2, e_4, e_6, e_8\}$  and  $G' := G \setminus \{e_2, e_4, e_6, e_8\}$ . Then  $N'$  is the graft matroid of  $(G', T)$ , and  $(G', T, \Gamma)$  is a quadrum. It therefore follows from Case 1 that  $(N, \Gamma) \setminus \{e_2, e_4, e_6, e_8\} = (N', \Gamma) \cong (F_7, E(F_7))$ . **Case 3:**  $(G, T, \Gamma)$  is a trifold. By definition,  $(G, T)$  is a plain trifold and  $\Gamma = E(N) - t$ . Remark 19 states that  $N/t \cong F_7$ , implying in turn that  $(N, \Gamma)/t \cong (F_7, E(F_7))$ . **Case 4:**  $(G, T, \Gamma)$  is a super trifold. Label the vertices and edges of  $G$  as in Figure 3 (b). After possibly

resigning on terminal  $v_1$ , we may assume by Remark 20 that  $\Gamma = \{e_1, e_2, e_3, e_4, e_5, e_6, e_9\}$ . Let  $(N', \Gamma) := (N, \Gamma) \setminus \{e_7, e_8\}$  and  $G' := G \setminus \{e_7, e_8\}$ . Then  $N'$  is the graft matroid of  $(G', T)$ , and  $(G', T, \Gamma)$  is a trifold. It therefore follows from Case 3 that  $(N, \Gamma) \setminus \{e_7, e_8\}/t = (N', \Gamma)/t \cong (F_7, E(F_7))$ .  $\square$

## 5. PROPOSITION 12

**5.1. Starting the proof.** Suppose (h1)-(h2) hold, that is,  $M$  is an internally 4-connected matroid where every element is in exactly three triangles. Let  $\Omega \in E(M)$ . We would like to show that if  $\Lambda(\Omega)$  is contained in a circuit of  $\text{si}(M/\Omega)$ , then either  $M$  has an induced  $K_4$  or  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor. Let us make the following assumptions:

### Further hypotheses

- (h3)  $\Omega \in E(M)$  is contained in the triangles  $\{\Omega, f, f'\}$ ,  $\{\Omega, g, g'\}$ ,  $\{\Omega, h, h'\}$  where  $\Lambda(\Omega) = \{f, g, h\}$ ,
- (h4)  $M_\Omega := M/\Omega \setminus \{f', g', h'\}$  and  $(M_\Omega, \Sigma_\Omega) := (M, E(M))/\Omega \setminus \{f', g', h'\}$ ,
- (h5)  $C$  is a circuit in  $M_\Omega$  of minimum cardinality that contains  $\{f, g, h\}$ ,
- (h6)  $M$  does not have an induced  $K_4$ .

Note that  $M_\Omega = \text{si}(M/\Omega)$ . We leave the following as an easy exercise for the reader:

**Remark 22.** Take a binary matroid  $N$ , an element  $e \in E(N)$ , and a subset  $D \subseteq E(N)$ . Then the following statements hold:

- a. if  $D$  is a circuit of  $N/e$ , then exactly one of  $D, D \cup \{e\}$  is a circuit of  $N$ ,
- b. if  $D$  is a cycle of  $N/e$ , then at least one of  $D, D \cup \{e\}$  is a cycle of  $N$ ,
- c. if  $D$  is a cycle of  $N$  and  $e \in D$ , then  $D - e$  is a cycle of  $N/e$ , and
- d. if  $D$  is a cycle of  $N$  and  $e \notin D$ , then  $D$  is a cycle of  $N/e$ .

A key milestone in the proof is the proposition below:

**Proposition 23.** Suppose (h1)-(h6) hold. Let  $D$  be a circuit of  $M_\Omega$  containing  $\{f, g, h\}$ . Then,

- a. the elements of  $D - \{f, g, h\}$  are in series in  $M|D \cup \{f', g', h', \Omega\}$ ,
- b.  $D - \{f, g, h\} \neq \emptyset$ , and for each  $t \in D - \{f, g, h\}$ ,  $M|(D \cup \{f', g', h', \Omega\})/(D - \{f, g, h, t\})$  is the graft matroid of a plain trifold,<sup>4</sup> and
- c. if  $|D|$  is odd, then  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor.

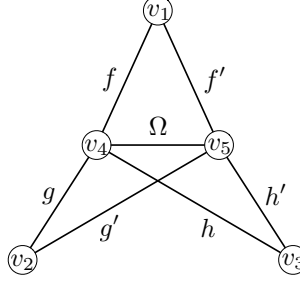
*Proof.* Let  $M' := M|D \cup \{f', g', h', \Omega\}$ .

(a) Note that  $M'/\Omega$  consists of the circuit  $D$  together with edges  $f', g', h'$ . Since  $\{f, f'\}, \{g, g'\}, \{h, h'\}$  are parallel classes of  $M'/\Omega$ , it follows that the elements of  $D - \{f, g, h\}$  are in series in  $M'/\Omega$ , so they are also in series in  $M'$ . (b) If  $\{f, g, h\}$  or  $\{f, g, h, \Omega\}$  is a circuit of  $M$ , then  $\{\Omega, f, f', g, g', h\}$  would be an induced  $K_4$  of  $M$ , which cannot occur by (h6). Hence, neither  $\{f, g, h\}$  nor  $\{f, g, h, \Omega\}$  is a circuit of  $M$  – this has two consequences. (1) Since one of  $D, D \cup \{\Omega\}$  is a circuit of  $M$  by Remark 22 (a), we get that  $D \neq \{f, g, h\}$ .

<sup>4</sup> $M|I/J$  is short-hand notation for  $(M|I)/J$ .

(2) The set  $\{f, g, h, \Omega\}$  is independent in the matroid  $M$ , and this in turn implies that  $M|\{f, f', g, g', h, h', \Omega\}$  is the cycle matroid of the graph  $G$  displayed below on vertices  $\{v_1, v_2, v_3, v_4, v_5\}$  and edges

$$\Omega = \{v_4, v_5\}, f = \{v_1, v_4\}, f' = \{v_1, v_5\}, g = \{v_2, v_4\}, g' = \{v_2, v_5\}, h = \{v_3, v_4\}, h' = \{v_3, v_5\}.$$



Let  $t \in D - \{f, g, h\}$  and

$$N := M'/(D - \{f, g, h, t\}).$$

By (a), the elements of  $D - \{f, g, h\}$  are in series in  $M'$ , so  $M'/(D - \{f, g, h, t\}) \setminus t = M' \setminus (D - \{f, g, h\})$ , implying in turn that  $N \setminus t = M|\{f, f', g, g', h, h', \Omega\} = \text{cycle}(G)$ . It therefore follows from Remark 10 that, for some  $T \subseteq V(G)$  of even cardinality,  $N$  is the graft matroid of the graft  $(G, T)$ . Since one of  $D, D \cup \{\Omega\}$  is a circuit of  $M$  by Remark 22 (a), it follows that one of  $\{f, g, h, t\}, \{f, g, h, t, \Omega\}$  is a circuit of  $N$ , implying in turn that one of  $\{f, g, h\}, \{f, g, h, \Omega\}$  is a  $T$ -join of  $G$ . This means that  $T = \{v_1, v_2, v_3, v_4\}$  or  $T = \{v_1, v_2, v_3, v_5\}$ . Either way, we see that  $(G, T)$  is a plain trifold. (c) Assume that  $|D|$  is odd. Let  $\Gamma := \{\Omega, f, f', g, g', h, h'\}$ . Notice that  $(G, T, \Gamma)$  is a trifold. Thus, by Proposition 21,  $(N, \Gamma)$  has an  $(F_7, E(F_7))$  minor. It therefore suffices to show that  $(N, \Gamma)$  is a minor of  $(M', E(M'))$ , which itself is a minor of  $(M, E(M))$ . Since  $|D|$  is odd,  $D - \{f, g, h\}$  has an even number of elements, all of which are in series in  $M'$  by (a), so  $D - \{f, g, h\}$  is a cocycle of  $M'$ . As a result,  $E(M') \Delta (D - \{f, g, h\}) = \Gamma$  is a signature of  $(M', E(M'))$ . However,  $(M', \Gamma)/(D - \{f, g, h, t\}) = (N, \Gamma)$ , so  $(N, \Gamma)$  is a minor of  $(M', E(M'))$ , as required.  $\square$

We may therefore make the following assumptions:

#### Further hypotheses

(h7) every circuit of  $M_\Omega$  containing  $\{f, g, h\}$  has even cardinality.

In particular,  $|C|$  is even and so  $C - \{f, g, h\} \neq \emptyset$ . Let  $S$  be a triangle of  $M_\Omega$  containing an element of  $C - \{f, g, h\}$ . We say that  $S$  is  $f$ -splitting if either  $S \cap \{f, g, h\} = \{f\}$  or the following statements hold:

- $|S \cap C| = 1$ , and
- $S \Delta C$  is the union of two disjoint circuits of  $M_\Omega$ , one of which contains  $f$  and the other contains  $g, h$ .

Similarly, we have  $g$ -splitting and  $h$ -splitting triangles.

**Corollary 24.** *Suppose (h1)-(h7) hold. Then every triangle of  $M_\Omega$  containing an element of  $C - \{f, g, h\}$  is a splitting triangle.*

*Proof.* Take an element  $e \in C - \{f, g, h\}$  and a triangle  $S$  of  $M_\Omega$  such that  $e \in S$ . Notice that  $|S \cap \{f, g, h\}| \leq 1$ . So if  $S \cap \{f, g, h\} \neq \emptyset$ , then  $S$  is a splitting triangle. We may therefore assume that  $S \cap \{f, g, h\} = \emptyset$ . Clearly,  $1 \leq |S \cap C| \leq 2$ . Note that  $|S \cap C| = 1$ ; for if not, then  $S \Delta C$  would be an odd-length circuit of  $M_\Omega$  containing  $\{f, g, h\}$ , which cannot occur by (h7). Consider now the odd-length cycle  $S \Delta C$ , which is either a circuit or the disjoint union of two circuits. However, it follows from (h7) that  $S \Delta C$  is the union of two disjoint circuits, both of which contain elements from  $\{f, g, h\}$ . This implies that  $R$  is a splitting triangle.  $\square$

The rest of this section is organized as follows: we will show that

- unless  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor, every element of  $C - \{f, g, h\}$  is in three otherwise disjoint triangles of  $M_\Omega$ , one of which is  $f$ -splitting, the second one is  $g$ -splitting, and the third one is  $h$ -splitting (§5.3),
- the circuit  $C$ , together with its splitting triangles, gives rise to a so-called Type I or a Type II configuration in  $(M_\Omega, \Sigma_\Omega)$  (§5.4),
- a Type I configuration gives a super trifold minor in  $(M, E(M))$ , and a Type II configuration gives a super quadrum minor in  $(M, E(M))$  (§5.2),

and by Proposition 21, the last step leads to an  $(F_7, E(F_7))$  minor, thereby finishing the proof of Proposition 12.

**5.2. Type I and Type II configurations.** In  $M_\Omega$ , take an element  $p \in E(M_\Omega) - C$  that is spanned by  $C$ . Then  $C \cup \{p\}$  contains exactly three circuits, one of which is  $C$ , the other two contain  $p$  and their symmetric difference is  $C$ ; notice that the other two circuits either separate  $f, g, h$  or not. We say that  $p$  is  $f$ -splitting if there is a circuit in  $C \cup \{p\}$  that contains  $f$  and none of  $g, h$ . Observe that if  $S$  is an  $f$ -splitting triangle, then each element of  $S - C$  is  $f$ -splitting. Similarly, we have  $g$ -splitting and  $h$ -splitting elements. If  $p$  is  $e$ -splitting, for some  $e \in \{f, g, h\}$ , we denote by  $\Theta(p)$  the circuit contained in  $C \cup \{p\}$  such that  $\Theta(p) \cap \{f, g, h\} = \{e\}$ .

In this section, we identify two configurations of splitting elements and show that their presence implies the existence of a super trifold or super quadrum minor in  $(M, E(M))$ .

We say  $\{p_1, p_2\} \subseteq E(M_\Omega)$  is a *Type I configuration* if the following statements hold:

- $p_1$  and  $p_2$  are  $e$ -splitting, for some  $e \in \{f, g, h\}$ ,
- $C - (\Theta(p_1) \cup \Theta(p_2) \cup \{f, g, h\}) \neq \emptyset$ ,
- $\Theta(p_1) \Delta \Theta(p_2)$  is an odd cycle of  $(M_\Omega, \Sigma_\Omega)$ ,
- $|\Theta(p_1)|$  and  $|\Theta(p_2)|$  are odd.

We will show that a Type I configuration leads to a super trifold in  $(M, E(M))$ . To prove this, however, we need an ingredient. Recall that by Remark 22 (a), if  $D$  is a circuit of  $M/\Omega$ , then exactly one of  $D, D \cup \{\Omega\}$  is a circuit of  $M$ ; the following proposition characterizes when  $D$  is the circuit in  $M$ :

**Remark 25.** *Suppose (h1)-(h7) hold. Let  $D$  be a circuit of  $M/\Omega$ . Then  $D$  is a circuit of  $M$  if, and only if, the parity of  $|D|$  is equal to the parity of  $D$  in  $(M, E(M))/\Omega$ . In particular, if  $D$  is a circuit of  $M_\Omega$ , then the following statements are equivalent:*

- i.  $D$  is a circuit of  $M$ ,

ii.  $|D|$  and  $|D \cap \Sigma_\Omega|$  have the same parity.

*Proof.* Let  $D$  be a circuit of  $M/\Omega$ . Assume that  $D$  is a circuit of  $M$ . Then the parity of  $|D|$  is equal to the parity of  $D$  in  $(M, E(M))$ , which is equal to the parity of  $D$  in  $(M, E(M))/\Omega$ . Conversely, assume that the parity of  $|D|$  is equal to the parity of  $D$  in  $(M, E(M))/\Omega$ . Suppose, for a contradiction, that  $D$  is not a circuit of  $M$ . By Remark 22 (a),  $D \cup \{\Omega\}$  is a circuit of  $M$ , and moreover, the parity of  $D \cup \{\Omega\}$  in  $(M, E(M))$  is equal to the parity of  $D$  in  $(M, E(M))/\Omega$ , which by assumption is equal to the parity of  $|D|$ . However, the parity of  $D \cup \{\Omega\}$  in  $(M, E(M))$  is equal to the parity of  $|D \cup \{\Omega\}| = |D| + 1$ , a contradiction.  $\square$

**Proposition 26.** *Suppose (h1)-(h7) hold. If there is a Type I configuration, then  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor.*

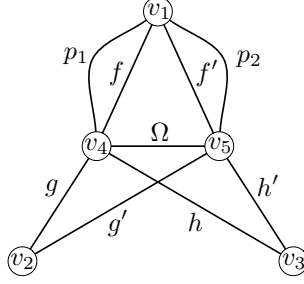
*Proof.* Assume that there is a Type I configuration  $\{p_1, p_2\}$ . After possibly interchanging the roles of  $f, g, h$ , we may assume that  $p_1, p_2$  are  $f$ -splitting, and after possibly interchanging the roles of  $p_1, p_2$ , we may assume that  $\Theta(p_1)$  is an odd circuit and  $\Theta(p_2)$  an even circuit of  $(M_\Omega, \Sigma_\Omega)$ . Since  $|\Theta(p_1)|$  and  $|\Theta(p_2)|$  are odd, it follows from Remark 25 that  $\Theta(p_1)$  is an odd circuit of  $(M, E(M))$ , and together with Remark 22 (a), that  $\Theta(p_2) \cup \{\Omega\}$  is an even circuit of  $(M, E(M))$ . Now take an element  $t \in C - (\Theta(p_1) \cup \Theta(p_2) \cup \{f, g, h\})$ . Consider the following minor of  $(M, E(M))$ :

$$(N, \Gamma) := (M, E(M)) / (C \cup \{f', g', h', \Omega, p_1, p_2\}) / (C - \{f, g, h, t\}).$$

We will show that  $(N, \Gamma)$  corresponds to a super trifold. By Proposition 23 (b),  $N \setminus \{p_1, p_2\}$  is the graft matroid of a plain trifold  $(G', T)$ , where  $V(G') = \{v_1, v_2, v_3, v_4, v_5\}$ ,  $E(G')$  consists of

$$\Omega = \{v_4, v_5\}, f = \{v_1, v_4\}, f' = \{v_1, v_5\}, g = \{v_2, v_4\}, g' = \{v_2, v_5\}, h = \{v_3, v_4\}, h' = \{v_3, v_5\},$$

and either  $T = \{v_1, v_2, v_3, v_4\}$  or  $T = \{v_1, v_2, v_3, v_5\}$ . (See below for an illustration.) Since the triangles  $\{\Omega, f, f'\}, \{\Omega, g, g'\}, \{\Omega, h, h'\}$  are odd in the signed matroid  $(M, E(M))$ , they are odd also in the minor  $(N, \Gamma) \setminus \{p_1, p_2\} = (N \setminus \{p_1, p_2\}, \Gamma - \{p_1, p_2\})$ . We may therefore assume that  $\{\Omega, f, f', g, g', h, h'\} \subseteq \Gamma - \{p_1, p_2\}$ . Notice that we do not know whether or not  $t$  belongs to  $\Gamma - \{p_1, p_2\}$ . Since  $\Theta(p_1)$  is a circuit of  $M$  containing  $\{f, p_1\}$  and all of its other edges belong to  $C - \{f, g, h, t\}$ , it follows that  $\{f, p_1\}$  is a circuit of  $N$ . Similarly, since  $\Theta(p_2) \cup \{\Omega\}$  is a circuit of  $M$  containing  $\{f, p_2, \Omega\}$  and all of its other edges belong to  $C - \{f, g, h, t\}$ , we get that  $\{f, p_2, \Omega\}$  is a triangle of  $N$ , which in turn implies that  $\{f', p_2\}$  is a circuit of  $N$ . As a consequence,  $N$  is the graft matroid of the graft  $(G, T)$  obtained from  $(G', T)$  after adding edge  $p_1$  parallel to  $f$ , and edge  $p_2$  parallel to  $f'$ . Since  $\Theta(p_1)$  is an odd circuit of  $(M, E(M))$ , we get that  $\{f, p_1\}$  is an odd circuit of  $(N, \Gamma)$ , so  $p_1 \notin \Gamma$ . Similarly, as  $\Theta(p_2) \cup \{\Omega\}$  is an even circuit of  $(M, E(M))$ , we get that  $\{f, p_2, \Omega\}$  is an even triangle of  $(N, \Gamma)$ , and as  $\{f, f', \Omega\}$  is an odd triangle, we have that  $\{f', p_2\}$  is also an odd circuit of  $(N, \Gamma)$ . Hence,  $p_2 \notin \Gamma$ . Therefore, the signed graft  $(G, T, \Gamma)$  is a super trifold.



It follows from Proposition 21 that  $(N, \Gamma)$ , and therefore  $(M, E(M))$ , has an  $(F_7, E(F_7))$  minor, as desired.  $\square$

We say  $\{p_1, p'_1, p_2, p_3\} \subseteq E(M_\Omega)$  is a *Type II configuration* if the following statements hold:

- $p_1$  and  $p'_1$  are  $e_1$ -splitting,  $p_2$  is  $e_2$ -splitting, and  $p_3$  is  $e_3$ -splitting, for a permutation  $e_1, e_2, e_3$  of  $f, g, h$ ,
- $\Theta(p_1) \cap \Theta(p'_1) \cap \Theta(p_2) \cap \Theta(p_3) \neq \emptyset$ ,
- $\Theta(p_1) \Delta \Theta(p'_1)$  is an odd cycle of  $(M_\Omega, \Sigma_\Omega)$ .

We will show that a Type II configuration leads to a super quadrum in  $(M, E(M))/\Omega$ , and therefore, in  $(M, E(M))$ :

**Proposition 27.** *Suppose (h1)-(h7) hold. If there is a Type II configuration, then  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor.*

*Proof.* Assume that there is a Type II configuration  $\{p_1, p'_1, p_2, p_3\}$ . By symmetry, we may assume that  $p_1, p'_1$  are  $f$ -splitting,  $p_2$  is  $g$ -splitting, and  $p_3$  is  $h$ -splitting. Take an element  $t \in \Theta(p_1) \cap \Theta(p'_1) \cap \Theta(p_2) \cap \Theta(p_3)$ . Observe that  $(M', \Sigma') := (M, E(M))/\Omega$  is obtained from  $(M_\Omega, \Sigma_\Omega)$  after adding edges  $f', g', h'$  parallel of different parity to  $f, g, h$ , respectively. Consider now the minor

$$(N, \Gamma) := (M', \Sigma')|(C \cup \{p_1, p'_1, p_2, p_3\})/(C - \{f, g, h, t\}).$$

We will show that  $(N, \Gamma)$  corresponds to a super quadrum.

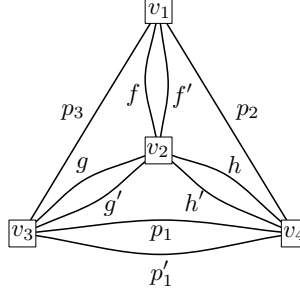
Since  $C$  is a circuit of  $M'$ , it follows that  $\{f, g, h, t\}$  is a circuit of  $N$ . Start with the graft  $(G'', T)$  on vertices  $\{v_1, v_2, v_3, v_4\}$  and edges  $f = \{v_2, v_1\}, g = \{v_2, v_3\}, h = \{v_2, v_4\}$ , where  $T = \{v_1, v_2, v_3, v_4\}$ . Note that  $N|\{f, g, h, t\}$  is the graft matroid of  $(G'', T)$ . Since  $\Theta(p_1)$  is a circuit of  $M_\Omega$ , it is also a circuit of  $M'$ , and as it contains  $\{f, p_1, t\}$  and all of its other edges belong to  $C - \{f, g, h, t\}$ , it follows that  $\{f, p_1, t\}$  is a cycle of  $N$ . Similarly,  $\{g, p_2, t\}$  and  $\{h, p_3, t\}$  are also cycles of  $N$ . As a consequence,  $N|\{f, g, h, t, p_1, p_2, p_3\}$  is the graft matroid of the plain quadrum  $(G', T)$  obtained from  $(G'', T)$  after adding  $p_1 = \{v_3, v_4\}, p_2 = \{v_4, v_1\}$  and  $p_3 = \{v_1, v_3\}$ .

Notice that  $N$  has no loop, because  $(C - \{f, g, h, t\}) \cup \{e\}$  contains no circuit of  $M'$ , for each  $e \in E(N)$ . Therefore, since  $\{f, f'\}, \{g, g'\}, \{h, h'\}$  are odd circuits in  $(M', \Sigma')$ , they are also odd circuits in  $(N, \Gamma)$ . Moreover, as  $\Theta(p_1) \Delta \Theta(p'_1)$  is an odd cycle of  $(M_\Omega, \Sigma_\Omega)$ , it is also an odd cycle of  $(M', \Sigma')$ , and because it contains  $\{p_1, p'_1\}$  and all of its other edges belong to  $C - \{f, g, h, t\}$ , it follows that  $\{p_1, p'_1\}$  is an odd circuit of  $(N, \Gamma)$ . Thus,  $N$  is the graft matroid of the graft  $(G, T)$  obtained from  $(G', T)$  after adding edges  $f', g', h', p'_1$  parallel

to  $f, g, h, p_1$ , respectively. (See below for an illustration.) Moreover,

$$|\Gamma \cap \{f, f'\}| = |\Gamma \cap \{g, g'\}| = |\Gamma \cap \{h, h'\}| = |\Gamma \cap \{p_1, p'_1\}| = 1,$$

and we do not know whether or not  $p_2, p_3, t$  belong to  $\Gamma$ . This means that  $(G, T, \Gamma)$  is a super quadrum.



From Proposition 21 we get that  $(N, \Gamma)$ , and therefore  $(M, E(M))$ , has an  $(F_7, E(F_7))$  minor, as desired.  $\square$

**5.3. Splitting triangles.** Take a signed matroid  $(N, \Gamma)$ . For  $R \subseteq E(N)$ , we write  $(N, \Gamma)|R := (N, \Gamma) \setminus (E(N) - R)$ . We say that  $\{e_1, \dots, e_6\}$  is in *induced odd  $K_4$*  of  $(N, \Gamma)$  if  $N|\{e_1, \dots, e_6\}$  is an induced  $K_4$  in which every triangle is odd in  $(N, \Gamma)|\{e_1, \dots, e_6\}$ . A consequence of Remark 25 is the following:

**Corollary 28.** *Suppose (h1)-(h7) hold. Then  $(M, E(M))/\Omega$  does not have an induced odd  $K_4$ .*

*Proof.* Suppose, for a contradiction, that  $\{e_1, \dots, e_6\}$  is an induced odd  $K_4$  of  $(M, E(M))/\Omega$ , whose odd triangles are  $\{e_1, e_2, e_3\}, \{e_1, e_4, e_5\}, \{e_2, e_4, e_6\}, \{e_3, e_5, e_6\}$ . It then follows from Remark 25 that these are also triangles of  $M$ , implying in turn that  $\{e_1, \dots, e_6\}$  is an induced  $K_4$  of  $M$ , thereby contradicting (h6).  $\square$

**Remark 29.** *Suppose (h1)-(h7) hold. If  $S$  is a triangle of  $M_\Omega$ , then  $|S \cap \{f, g, h\}| \leq 1$ .*

*Proof.* Let  $S$  be a triangle of  $M_\Omega$ . It follows from Proposition 23 (b) that  $|S \cap \{f, g, h\}| \leq 2$ . Suppose, for a contradiction, that  $|S \cap \{f, g, h\}| = 2$ . We may assume that  $S \cap \{f, g, h\} = \{f, g\}$ . By Remark 22 (a), one of  $S, S \cup \{\Omega\}$  is a circuit of  $M$ , so one of  $S, S \Delta \{f, f'\}$  is a triangle of  $M$ . But then  $S \cup \{\Omega, f, f', g, g'\}$  is an induced  $K_4$  of  $M$ , a contradiction to (h6).  $\square$

We are now ready to prove the main result of this section:

**Proposition 30.** *Suppose (h1)-(h7) hold, and  $(M, E(M))$  does not have  $(F_7, E(F_7))$  as a minor. Then for every element  $e \in C - \{f, g, h\}$ , there exist triangles  $S_f, S_g, S_h$  of  $M_\Omega$  such that  $S_f \cap S_g = S_g \cap S_h = S_h \cap S_f = \{e\}$ , and for each  $z \in \{f, g, h\}$ ,*

- $S_z$  is  $z$ -splitting, and
- if  $S_z \cap \{f, g, h\} = \emptyset$ , then  $S_z$  is odd in  $(M_\Omega, \Sigma_\Omega)$ .

*Proof.* Take an element  $e \in C - \{f, g, h\}$ . Denote by  $T_1, T_2, T_3$  the three triangles of  $M$  containing  $e$ , whose existence is guaranteed by (h2). Recall that  $T_1 \cap T_2 = T_2 \cap T_3 = T_3 \cap T_1 = \{e\}$ . Since  $e \notin \{\Omega, f, f', g, g', h, h'\}$ ,  $\Omega \notin T_1 \cup T_2 \cup T_3$ . Therefore, since  $M$  is a simple matroid and  $M/\Omega$  is a loopless matroid whose non-trivial



parallel classes are precisely  $\{f, f'\}, \{g, g'\}, \{h, h'\}$ , it follows that  $T_1, T_2, T_3$  are also triangles of  $M/\Omega$ ; note that they are odd triangles of  $(M, E(M))/\Omega$ . For each  $i \in [3]$ , let  $S_i$  be the triangle corresponding to  $T_i$  in the simplification  $M_\Omega$ . We will show that, after a relabeling,  $S_1, S_2, S_3$  are the desired three triangles.

**1.**  $S_1 \cap S_2 = S_2 \cap S_3 = S_3 \cap S_1 = \{e\}$ . Moreover, for each  $i \in [3]$ ,  $S_i$  is a splitting triangle, and if  $S_i \cap \{f, g, h\} = \emptyset$ , then  $S_i$  is odd in  $(M_\Omega, \Sigma_\Omega)$ .

*Subproof.* Suppose, for a contradiction, that  $\{e\} \subsetneq S_1 \cap S_2$ . Since  $\{e\} = T_1 \cap T_2$ , we may assume that  $f \in S_1 \cap S_2$ ,  $f \in T_1$  and  $f' \in T_2$ . However, since  $\{f, f', \Omega\}$  is a triangle of  $M$ , it follows that  $T_1 \cup T_2 \cup \{\Omega\}$  is an induced  $K_4$  of  $M$ , a contradiction to (h6). Thus,  $S_1 \cap S_2 = \{e\}$  and similarly,  $S_2 \cap S_3 = S_3 \cap S_1 = \{e\}$ . Take an index  $i \in [3]$ . Clearly, if  $S_i \cap \{f, g, h\} = \emptyset$ , then  $S_i = T_i$  and therefore  $S_i$  is an odd triangle of  $(M_\Omega, \Sigma_\Omega)$ . Moreover, since  $e \in S_i$ , we get from Corollary 24 that  $S_i$  is a splitting triangle.  $\diamond$

It therefore suffices to show that no two of  $S_1, S_2, S_3$  split the same element of  $\{f, g, h\}$ . Suppose, for a contradiction, that  $S_1, S_2$  are  $f$ -splitting. Since these triangles are  $f$ -splitting, it follows that  $S_1 \cap \{g, h\} = S_2 \cap \{g, h\} = \emptyset$ , and by (1),  $S_1 \cap S_2 = \{e\}$ .

Fix an index  $i \in [2]$ . Let us carefully label the elements of  $S_i - e$ . If  $f \in S_i$ , then let  $p_i := f$  and  $q_i$  the element in  $S_i - \{e, f\}$ . Otherwise,  $f \notin S_i$ . Because  $S_i$  is  $f$ -splitting,  $S_i \cap C = \{e\}$  and  $S_i \triangle C$  is the union of two disjoint circuits of  $M_\Omega$ . That is, the elements of  $S_i - \{e\}$  are  $f$ -splitting, and for a labeling  $p_i, q_i$  of these elements,  $S_i \triangle C$  is the disjoint union of  $\Theta(p_i)$  and  $C \triangle \Theta(q_i)$ .

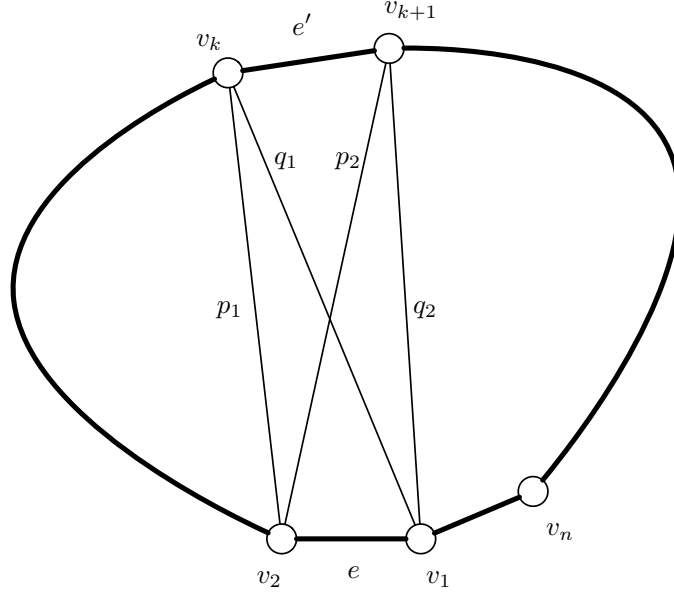
Since  $M_\Omega$  is a simple matroid, when  $f \notin \{p_1, p_2\}$ , we get that  $\Theta(p_1) - p_1 \neq \Theta(p_2) - p_2$ . We may therefore assume that  $f \neq p_2$  and, if  $f \neq p_1$ ,  $(\Theta(p_2) - p_2) - (\Theta(p_1) - p_1) \neq \emptyset$ .

**2.**  $M_\Omega|(C \cup \{p_1, q_1, p_2, q_2\})$  is the cycle matroid of a simple graph  $G$  described as follows: for some integers  $n, k$  such that  $n - 2 \geq k \geq 3$ ,

- $V(G) = \{v_1, \dots, v_n\}$ ,
- $C = \{\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}\}$ ,
- $e = \{v_1, v_2\}$ ,  $p_1 = \{v_2, v_k\}$ ,  $q_1 = \{v_1, v_k\}$ ,  $p_2 = \{v_2, v_{k+1}\}$ ,  $q_2 = \{v_1, v_{k+1}\}$ , and
- $f \in \{\{v_2, v_3\}, \dots, \{v_{k-1}, v_k\}\}$  and  $g, h \in \{\{v_{k+1}, v_{k+2}\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}\}$ .

*Subproof.* Let  $n := |C|$ . Clearly,  $M_\Omega|C$  is the cycle matroid of the simple graph  $G_1$  on vertices  $\{v_1, \dots, v_n\}$  whose edges are  $\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}$ . We assume that  $e = \{v_1, v_2\}$  and the edges of  $\Theta(q_1) - q_1$  appear consecutively on the graph circuit. Then there is an integer  $k \geq 3$  such that  $\Theta(q_1) - q_1 = \{\{v_1, v_2\}, \{v_2, v_3\}, \dots, \{v_{k-1}, v_k\}\}$ . Note that  $f \in \{\{v_2, v_3\}, \dots, \{v_{k-1}, v_k\}\}$ , and  $f = p_1$  if and only if  $k = 3$ .

Let  $G_2$  be the graph obtained from  $G_1$  after adding the edge  $q_1 = \{v_1, v_k\}$ , and if  $k > 3$ , the edge  $p_1 = \{v_2, v_k\}$ . Note that  $M_\Omega|(C \cup \{p_1, q_1\})$  is the cycle matroid of the simple graph  $G_2$ . Consider the set  $\Theta(p_2) - p_2$ . As  $f \in \Theta(p_2)$ , we have  $(\Theta(p_2) - p_2) \cap \{\{v_2, v_3\}, \dots, \{v_{k-1}, v_k\}\} \neq \emptyset$ , and as  $(\Theta(p_2) - p_2) - (\Theta(p_1) - p_1) \neq \emptyset$  when  $k > 3$ , we have  $(\Theta(p_2) - p_2) \cap \{\{v_k, v_{k+1}\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}\} \neq \emptyset$ . After possibly rearranging the edges of  $G_2$  within series classes  $\{\{v_2, v_3\}, \dots, \{v_{k-1}, v_k\}\}$  and  $\{\{v_k, v_{k+1}\}, \dots, \{v_{n-1}, v_n\}, \{v_n, v_1\}\}$ ,



**Figure 4.** An illustration of graph  $G$ , where the edges in  $C$  are bold.

we may assume that the edges of  $\Theta(p_2) - p_2$  appear consecutively on the circuit  $C$ . So there are indices  $i, j \in [n]$  such that

$$\Theta(p_2) - p_2 = \{\{v_i, v_{i+1}\}, \dots, \{v_{j-1}, v_j\}\}$$

where  $k - 1 \geq i \geq 2$  and  $n - 1 \geq j \geq k + 1$ .

Let  $G_3$  be the graph obtained from  $G_2$  after adding the edge  $p_2 = \{v_i, v_j\}$ . Note that  $M_\Omega|(C \cup \{p_1, q_1, p_2\})$  is the cycle matroid of the simple graph  $G_3$ . We will show that  $i = 2$  and  $j = k + 1$ . Consider the following circuit of  $G_3$ :

$$\{p_2\} \cup \{\{v_i, v_{i+1}\}, \dots, \{v_{k-1}, v_k\}\} \cup \{q_1\} \cup \{\{v_1, v_n\}, \{v_n, v_{n-1}\}, \dots, \{v_{j+1}, v_j\}\}.$$

This circuit contains edges  $f, g, h$  and has  $n - (i - 1) - (j - k) + 2$  many edges. It therefore follows from the minimality of  $C$  in (h5) that  $n - (i - 1) - (j - k) + 2 \geq n$ , implying in turn that  $i = 2$  and  $j = k + 1$ . Now let  $G$  be the graph obtained from  $G_3$  after adding the edge  $q_2 = \{v_1, v_{k+1}\}$ . It is clear that  $M_\Omega|(C \cup \{p_1, q_1, p_2, q_2\})$  is the cycle matroid of the simple graph  $G$ , which is the desired graph.  $\diamond$

By Remark 29, edges  $g, h$  do not lie in a triangle of  $G$ , so in fact  $n - 3 \geq k$ . Let  $e' := \{v_k, v_{k+1}\} \in E(G) = E(M_\Omega)$  and note that  $\{e, p_1, q_1, p_2, q_2, e'\}$  is an induced  $K_4$  of  $M_\Omega$ . Let

$$(N, \Gamma) := (M_\Omega, \Sigma_\Omega)|\{e, p_1, q_1, p_2, q_2, e'\}.$$

The triangle  $S_2 = \{e, p_2, q_2\}$ , being disjoint from  $\{f, g, h\}$ , is odd in  $(N, \Gamma)$ , and if  $f \neq p_1$ , then the triangle  $S_1 = \{e, p_1, q_1\}$  would also be odd in  $(N, \Gamma)$ . Since  $M$  has no induced  $K_4$  by (h6), it follows from Corollary 28 that exactly two of  $\{e, p_1, q_1\}, \{e', p_1, p_2\}, \{e', q_1, q_2\}$  are even in  $(N, \Gamma)$ . Thus, if  $f \neq p_1$ , then  $\{e\}$  is a signature for  $(N, \Gamma)$ , and if  $f = p_1$ , then one of  $\{e\}, \{e, f\}, \{e, q_1\}$  is a signature for  $(N, \Gamma)$ .

3.  $f \neq p_1$ .

*Subproof.* Suppose, for a contradiction, that  $f = p_1$ . We will show that  $\{p_2, q_1\}$  is a Type I configuration. Recall that  $p_2, q_1$  are  $f$ -splitting, and since  $n - 3 \geq k$ , it follows that  $C - (\Theta(p_2) \cup \Theta(q_1) \cup \{f, g, h\}) \neq \emptyset$ . Moreover,  $|\Theta(p_2)| = |\Theta(q_1)| = 3$ . If  $\{e, q_1\}$  is a signature for  $(N, \Gamma)$ , then  $\{e, f, q_1, p_2, q_2, e'\} \Delta \{f, f'\}$  is an induced odd  $K_4$  of  $(M, E(M))/\Omega$ , which cannot occur by Corollary 28. Thus, one of  $\{e\}, \{e, f\}$  is a signature for  $(N, \Gamma)$ . Either way, we see that  $\Theta(p_2) \Delta \Theta(q_1)$  is odd cycle of  $(M_\Omega, \Sigma_\Omega)$ . Thus,  $\{p_2, q_1\}$  is a Type I configuration. But then Proposition 26 implies that  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor, a contradiction to our hypothesis.  $\diamond$

Recall that  $p_1, q_1, p_2, q_2$  are  $f$ -splitting, and since  $n - 3 \geq k$ ,

$$C - (\Theta(p_2) \cup \Theta(q_1) \cup \{f, g, h\}) = C - (\Theta(p_1) \cup \Theta(q_2) \cup \{f, g, h\}) \neq \emptyset.$$

We also know that  $\{e\}$  is a signature for  $(N, \Gamma)$ . It can be readily seen that if  $|\Theta(p_1) - p_1|$  is odd, then  $\{p_2, q_1\}$  is a Type I configuration, and otherwise,  $\{p_1, q_2\}$  is a Type I configuration. Either way, we get from Proposition 26 that  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor, thereby contradicting our hypothesis. Therefore,  $S_1$  and  $S_2$  cannot both be  $f$ -splitting. Similarly, no two of  $S_1, S_2, S_3$  split the same element. Among these triangles, let  $S_f$  be the  $f$ -splitting one,  $S_g$  the  $g$ -splitting one, and  $S_h$  the  $h$ -splitting one. These are the desired triangles, and the proof of Proposition 30 is finished.  $\square$

**5.4. Proof of Proposition 12.** We may assume that (h1)-(h7) hold. To remind the reader why, assume that (h1)-(h2), as well as the setup conditions (h3)-(h5), hold. If  $M$  has an induced  $K_4$ , then we are done. Otherwise, (h6) holds. We will prove that  $(M, E(M))$  has  $(F_7, E(F_7))$  as a minor, thereby finishing the proof of Proposition 12. Suppose otherwise. It then follows from Proposition 23 (c) that (h7) holds.

1.  $|C| \geq 6$ .

*Subproof.* Suppose otherwise. By (h7),  $|C|$  is even. Thus, for some  $t \in E(M_\Omega)$ ,  $C = \{f, g, h, t\}$ . By Proposition 30, there is an  $h$ -splitting triangle  $\{t, h, p\}$ , where  $p$  is  $h$ -splitting. But then  $\{p, f, g\}$  is a triangle of  $M_\Omega$ , thereby contradicting Remark 29.  $\diamond$

2. There is an  $f$ -splitting triangle  $S$  where  $|S \cap C| = 1$  and  $S$  is odd in  $(M_\Omega, \Sigma_\Omega)$ .

*Subproof.* Suppose otherwise. By (1), there are distinct elements  $e_1, e_2, e_3 \in C - \{f, g, h\}$ . Fix an index  $i \in [3]$ . Then by Proposition 30 and our contrary assumption,  $e_i$  is contained in an  $f$ -splitting triangle  $S_i$  such that  $S_i \cap C = \{e_i, f\}$ . By Remark 22 (a), one of  $S_i, S_i \cup \{\Omega\}$  is a circuit of  $M$ , implying in turn that one of  $S_i, S_i \Delta \{f, f'\}$  is a triangle of  $M$ . By (h2),  $f$  and  $f'$  are each in exactly 3 triangles of  $M$ , a common one being  $\{\Omega, f, f'\}$ . Hence, it cannot be that each one of  $S_1, S_2, S_3$  is a triangle of  $M$  or that each one of  $S_1 \Delta \{f, f'\}, S_2 \Delta \{f, f'\}, S_3 \Delta \{f, f'\}$  is a triangle of  $M$ . We may therefore assume that  $S_1, S_2 \Delta \{f, f'\}$  are triangles of  $M$ . In other words,  $S_1, S_2 \cup \{\Omega\}$  are circuits of  $M$ , so by Remark 25,  $S_1$  is an odd triangle and  $S_2$  is an even triangle of  $(M_\Omega, \Sigma_\Omega)$ . Let  $p_1$  be the element in  $S_1 - \{e_1, f\}$  and  $p_2$  the element in  $S_2 - \{e_2, f\}$ . Then  $p_1, p_2$  are  $f$ -splitting elements for which  $\Theta(p_1) = S_1$  and  $\Theta(p_2) = S_2$ . Since  $e_3 \in C - (S_1 \cup S_2 \cup \{f, g, h\})$ ,

it follows that  $\{p_1, p_2\}$  is a Type I configuration. By Proposition 26,  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor, a contradiction.  $\diamond$

Write  $S = \{e, p_1, p'_1\}$  where  $C \cap S = \{e\}$  and  $e \in \Theta(p'_1)$ . Note that  $\Theta(p'_1) \Delta \Theta(p_1) = S$ , so it is an odd cycle of  $(M_\Omega, \Sigma_\Omega)$ . As  $M_\Omega$  is simple, there is an element  $t \in \Theta(p_1) - \{f, p_1\}$ . Note that  $t \in \Theta(p'_1)$ . By Proposition 30,  $t$  is contained in a  $g$ -splitting triangle  $S_2$  and an  $h$ -splitting triangle  $S_3$ . Pick the  $g$ -splitting element  $p_2 \in S_2$  for which  $t \in \Theta(p_2)$  and the  $h$ -splitting element  $p_3 \in S_3$  for which  $t \in \Theta(p_3)$ . Then  $\{p_1, p'_1, p_2, p_3\}$  is a Type II configuration. By Proposition 27,  $(M, E(M))$  has an  $(F_7, E(F_7))$  minor, which is a contradiction. This finishes the proof of Proposition 12.  $\square$

## 6. EVEN CYCLE MATROIDS

Let  $G$  be a graph and  $\Gamma \subseteq E(G)$ . The signed matroid  $(\text{cycle}(G), \Gamma)$  is identified as  $(G, \Gamma)$  and is simply referred to as a *signed graph*. Zaslavsky [22] proved that the even cycles of  $(G, \Gamma)$  are the cycles of a (binary) matroid that we call the *even cycle matroid* of  $(G, \Gamma)$  and denote by  $\text{ecycle}(G, \Gamma)$ . Notice that every signature of  $(G, \Gamma)$  is a cocycle of  $\text{ecycle}(G, \Gamma)$ .

Given a graph  $H$  and a new edge label  $e$ , denote by  $H + e$  any graph obtained from  $H$  after adding  $e$  as a loop. The following folklore result states that matroids one contraction away from being graphic are even cycle matroids:

**Remark 31.** *Take a binary matroid  $M$  and an element  $e \in E(M)$  such that  $M/e = \text{cycle}(H)$ , for some graph  $H$ . If  $\Gamma$  is a cocycle of  $M$  containing  $e$ , then  $M = \text{ecycle}(H + e, \Gamma)$ .*

*Proof.* Let  $\Gamma$  be a cocycle of  $M$  containing  $e$ . Let  $C \subseteq E(M)$ . We need to show that  $C$  is a cycle of  $M$  if, and only if,  $C$  is an even cycle of  $(H + e, \Gamma)$ .  $(\Rightarrow)$  Suppose first that  $C$  is a cycle of  $M$ . Since  $|C \cap \Gamma|$  is even, it suffices to show that  $C$  is a cycle of  $H + e$ . If  $e \in C$ , then  $C - e$  is a cycle of  $M/e$  by Remark 22 (c), so it is also a cycle of  $H$ , implying in turn that  $C$  is a cycle of  $H + e$ . Otherwise,  $e \notin C$ , so  $C$  is a cycle of  $M/e$  by Remark 22 (d), implying in turn that  $C$  is a cycle of  $H$ , and therefore, of  $H + e$ .  $(\Leftarrow)$  Suppose conversely that  $C$  is an even cycle of  $(H + e, \Gamma)$ . Assume first that  $e \notin C$ . Then  $C$  is a cycle of  $H$ , so it is also a cycle of  $M/e$ , implying by Remark 22 (b) that either  $C$  or  $C \cup \{e\}$  is a cycle of  $M$ . Since  $|C \cap \Gamma|$  is even,  $e \in \Gamma$  and  $e \notin C$ , it follows that  $|(C \cup \{e\}) \cap \Gamma|$  is odd, and because  $\Gamma$  is a cocycle of  $M$ ,  $C \cup \{e\}$  cannot be a cycle of  $M$ . Thus,  $C$  is a cycle of  $M$ . Assume in the remaining case that  $e \in C$ . Then  $C - e$  is a cycle of  $H$ , so it is also a cycle of  $M/e$ , and thus by Remark 22 (b), one of  $C - e, C$  is a cycle of  $M$ . Since  $e \in C \cap \Gamma$  and  $|C \cap \Gamma|$  is even, it follows that  $|(C - e) \cap \Gamma|$  is odd. Therefore, because  $\Gamma$  is a cocycle of  $M$ ,  $C - e$  cannot be a cycle of  $M$ , and as a result,  $C$  is a cycle of  $M$ .  $\square$

**6.1. Even cycle matroids and connectivity.** Let  $G$  be a graph. For a subset  $X \subseteq E(G)$ , we denote by  $V_G(X)$  the ends of the edges in  $X$ , and by  $G[X]$  the subgraph on vertices  $V_G(X)$  and edges  $X$ .

Let  $(G, \Gamma)$  be a signed graph, where  $G$  is connected. If  $(G, \Gamma)$  has no odd circuit, then  $\text{ecycle}(G, \Gamma) = \text{cycle}(G)$  and therefore, any spanning tree of  $G$  is a basis for  $\text{ecycle}(G, \Gamma)$ . Otherwise, when  $(G, \Gamma)$  has an odd

circuit,  $T \cup \{e\}$  is a basis for  $\text{ecycle}(G, \Gamma)$ , where  $T$  is a spanning tree of  $G$ , and  $e \in E(G) - T$  is chosen so that  $T \cup \{e\}$  contains an odd circuit of  $(G, \Gamma)$ .

The next remark describes the connectivity function for even cycle matroids.

**Remark 32** ([7]). *Let  $(G, \Gamma)$  be a signed graph, where  $G$  is connected. Take a non-empty and proper subset  $X \subseteq E(G)$  where both  $G[X]$  and  $G[\bar{X}]$  are connected. Then*

$$\lambda_{\text{ecycle}(G, \Gamma)}(X) \leq \lambda_{\text{cycle}(G)}(X) + 1 = |V_G(X) \cap V_G(\bar{X})|.$$

*Proof.* The equation is a routine exercise (see [12] Lemma 8.1.7 for a proof). To prove the inequality, let  $E := E(G)$ ,  $M := \text{cycle}(G)$  and  $M' := \text{ecycle}(G, \Gamma)$ . Denote by  $r, r'$  the rank functions of  $M, M'$ , respectively. If  $(G, \Gamma)$  has no odd circuit, then  $M = M'$ , so  $r = r'$ , implying in turn that  $\lambda_M = \lambda_{M'}$ . We may therefore assume that  $(G, \Gamma)$  has an odd circuit. What we argued above implies that  $r'(E) = r(E) + 1$ ,  $r'(X) \in \{r(X), r(X) + 1\}$  and  $r'(\bar{X}) \in \{r(\bar{X}), r(\bar{X}) + 1\}$ , so

$$\lambda_{M'}(X) = r'(X) + r'(\bar{X}) - r'(E) \leq r(X) + 1 + r(\bar{X}) + 1 - r(E) - 1 = \lambda_M(X) + 1,$$

as required.  $\square$

A connected graph on at least 3 vertices is *2-connected* if it remains connected after deleting any vertex. A 2-connected graph on at least 4 vertices is *3-connected* if it remains connected after deleting any pair of vertices. For a graph  $G$ , denote the set of all loops by  $\text{loops}(G)$ . Given a signed graph  $(G, \Gamma)$ , denote by  $\text{si}(G, \Gamma)$  the signed graph obtained after deleting all even loops, deleting all odd loops except for one, and deleting all but one edge from each class of parallel edges in  $G$  of the same parity in  $(G, \Gamma)$ .

**Proposition 33.** *Let  $(G, \Gamma)$  be a signed graph that has an odd loop  $e$ , and let  $N := \text{ecycle}(G, \Gamma)$ . Then,*

- a. *if  $N$  is simple and cosimple,  $e$  is the unique loop of  $G$ , parallel edges of  $G$  have distinct parities in  $(G, \Gamma)$ , and  $G$  does not have edges in series,*
- b. *if  $N$  is internally 4-connected and  $|E(N)| \geq 8$ , then  $G \setminus e$  is 3-connected,*
- c.  *$\text{si}(N) = \text{ecycle}(\text{si}(G, \Gamma))$ ,*
- d. *if  $\text{si}(N)$  is internally 4-connected and  $|E(\text{si}(N))| \geq 8$ , then  $G \setminus \text{loops}(G)$  is 3-connected.*

*Proof.* **(a)** Suppose  $N$  is simple and cosimple. Then  $N$  has no cycle of size at most 2 and no two elements in series. In particular,  $(G, \Gamma)$  does not have an even loop or an even cycle of size two, and  $G$  does not have two edges in series. Since  $e$  is an odd loop, there cannot be another odd loop. As a result,  $e$  is the unique loop of  $G$  and parallel edges of  $G$  have distinct parities in  $(G, \Gamma)$ . **(b)** Suppose  $N$  is internally 4-connected. In particular,  $N$  is simple and cosimple. Thus, (a) implies that in  $G$ , no two edges are in series,  $e$  is the unique loop, and every parallel class has size at most two. Since  $|E(G \setminus e)| \geq 7$ , we get that  $G \setminus e$  has at least 4 vertices. It suffices to show that when  $G \setminus e$  is connected, then it is 3-connected. We first show that  $G \setminus e$  is 2-connected. Suppose, for a contradiction, that there is a non-trivial partition  $X, Y$  of  $E(G \setminus e)$  such that  $|V_{G \setminus e}(X) \cap V_{G \setminus e}(Y)| = 1$ . Since  $|X| + |Y| \geq 7$ , after possibly interchanging the roles of  $X$  and  $Y$ , we may assume that  $|X| \geq 2$ . Let  $\bar{X} := Y \cup \{e\}$ . Then  $|\bar{X}| \geq 2$ . Assuming the end of  $e$  belongs to  $V_{G \setminus e}(Y)$ , we see that  $G, G[X], G[\bar{X}]$  are

connected. Hence, Remark 32 implies that  $\lambda_N(X) \leq |V_G(X) \cap V_G(\bar{X})| = 1$ , so  $(X, \bar{X})$  is a 2-separation of  $N$ , a contradiction as  $N$  is 3-connected. It remains to show that  $G \setminus e$  is 3-connected. Suppose, for a contradiction, that there is a non-trivial partition  $X, Y$  of  $E(G \setminus e)$  such that  $|V_{G \setminus e}(X) \cap V_{G \setminus e}(Y)| = 2$ ,  $|V_{G \setminus e}(X)| \geq 3$  and  $|V_{G \setminus e}(Y)| \geq 3$ . Since  $G \setminus e$  is 2-connected, it follows that  $(G \setminus e)[X], (G \setminus e)[Y]$  are connected, implying in turn that  $|X| \geq 2$  and  $|Y| \geq 2$ . In fact, since  $G \setminus e$  does not have two edges in series, we have  $|X| \geq 3$  and  $|Y| \geq 3$ . Because  $|X| + |Y| \geq 7$ , we may assume that  $|X| \geq 4$ . Let  $\bar{X} := Y \cup \{e\}$ . Then  $|\bar{X}| \geq 4$ . Assuming the end of  $e$  belongs to  $V_{G \setminus e}(Y)$ , we see that  $G, G[X], G[\bar{X}]$  are connected. Thus, by Remark 32, we get that  $\lambda_N(X) \leq |V_G(X) \cap V_G(\bar{X})| = 2$ , so  $(X, \bar{X})$  is a 3-separation of  $N$ , a contradiction as  $N$  is internally 4-connected. (c) is immediate. (d) Let  $(G', \Gamma') := \text{si}(G, \Gamma)$ . We may assume that  $e$  is also an odd loop of  $(G', \Gamma')$ . By (c),  $\text{si}(N) = \text{ecycle}(G', \Gamma')$ , so from (b) we get that  $G' \setminus e$  is 3-connected. As  $G$  is obtained from  $G'$  by adding loops and edges parallel to existing ones, we get that  $G \setminus \text{loops}(G)$  is also 3-connected.  $\square$

**6.2. Even cycle matroids that are graphic.** Here we characterize, under relevant conditions, when an even cycle matroid is graphic. A complete and technical answer to this problem was obtained by Shih in his PhD thesis [20] but was never published in a refereed journal – our arguments will not rely on this characterization. We will need the following seminal result of Whitney [21]:

**Theorem 34.** *Let  $G, G'$  be graphs over the same edge set such that  $\text{cycle}(G) = \text{cycle}(G')$ . If  $G \setminus \text{loops}(G)$  is 3-connected, then  $G \setminus \text{loops}(G) = G' \setminus \text{loops}(G')$  and  $\text{loops}(G) = \text{loops}(G')$ .*

Let  $(G, \Gamma)$  be a signed graph, and take a vertex  $v \in V(G)$ . We say  $v$  is a *blocking vertex* if every non-loop odd circuit of  $(G, \Gamma)$  uses  $v$ . It follows from Proposition 2 that  $v$  is a blocking vertex if, and only if, there is a signature contained in  $\delta_G(v) \cup \text{loops}(G)$ .

**Remark 35.** *Let  $(G, \Gamma)$  be a signed graph. If  $(G, \Gamma)$  has a blocking vertex, then  $\text{ecycle}(G, \Gamma)$  is graphic.*

*Proof.* Let  $v$  be a blocking vertex. After possibly resigning, we may assume that  $\Gamma \subseteq \delta_v(G) \cup \text{loops}(G)$ . We may also assume that every odd loop is incident to  $v$  and every even loop is incident to another vertex. Let  $H$  be the graph obtained from  $G$  after splitting  $v$  into vertices  $v_1, v_2$  such that every edge in  $\delta_G(v) \cap \Gamma$  is incident with  $v_1$ , every edge in  $\delta_G(v) - \Gamma$  is incident with  $v_2$ , and every odd loop has ends  $v_1, v_2$ . It can be readily checked that  $\text{ecycle}(G, \Gamma) = \text{cycle}(H)$ .  $\square$

Provided an odd loop and 3-connectedness, we can guarantee the converse also holds:

**Proposition 36.** *Let  $(G, \Gamma)$  be a signed graph that has an odd loop and  $G \setminus \text{loops}(G)$  is 3-connected. If  $\text{ecycle}(G, \Gamma)$  is graphic, then  $(G, \Gamma)$  has a blocking vertex.*

*Proof.* Set  $E := E(G)$  and let  $e \in E$  be an odd loop of  $(G, \Gamma)$ . Let  $H$  be a graph with edge set  $E$  such that  $\text{ecycle}(G, \Gamma) = \text{cycle}(H)$ . As  $e$  is not an even loop of  $(G, \Gamma)$ ,  $e$  is not a loop of  $H$ ; let  $v_1, v_2$  be the ends of  $e$  in  $H$ . Since the even circuits of  $(G, \Gamma)$  are precisely the circuits of  $H$  we have, for  $C \subseteq E$ , the following correspondence:

- $C$  is an odd circuit of  $(G, \Gamma)$  if, and only if,  $C$  is a  $v_1v_2$ -path of  $H$ ,
- $C$  is an even circuit of  $(G, \Gamma)$  if, and only if,  $C$  is a circuit of  $H$ .

Let  $G'$  be the graph obtained from  $H$  after identifying vertices  $v_1$  and  $v_2$ ; call the identified vertex  $v$ . Let  $\Gamma' := \delta_H(v_1)$ . Then the correspondence above implies that  $\text{cycle}(G') = \text{cycle}(G)$  and  $\text{ecycle}(G', \Gamma') = \text{ecycle}(G, \Gamma)$ . Since  $G \setminus \text{loops}(G)$  is 3-connected, it follows from Theorem 34 that  $G' \setminus \text{loops}(G') = G \setminus \text{loops}(G)$  and  $\text{loops}(G) = \text{loops}(G')$ . After changing the ends of the loops of  $G'$ , if necessary, we may assume that  $G' = G$ . Since  $\text{ecycle}(G, \Gamma') = \text{ecycle}(G, \Gamma)$ ,  $\Gamma'$  is a signature of  $(G, \Gamma)$  and as  $\Gamma' \subseteq \delta_G(v) \cup \text{loops}(G)$ , we see that  $v$  is a blocking vertex of  $(G, \Gamma)$ .  $\square$

**6.3. Blocking pairs.** Let  $(G, \Gamma)$  be a signed graph. Take disjoint  $I, J \subseteq E(G)$ . If  $I$  contains an odd circuit, we define  $(G, \Gamma)/I \setminus J := (G/I \setminus J, \emptyset)$ . Otherwise, by Proposition 2, there is a signature  $\Gamma'$  that is disjoint from  $I$ , and we define  $(G, \Gamma)/I \setminus J := (G/I \setminus J, \Gamma' - J)$ . We call  $(G, \Gamma)/I \setminus J$  a *minor* of  $(G, \Gamma)$ . Notice that minors are defined only up to resigning, and since  $\text{cycle}(G)/I \setminus J = \text{cycle}(G/I \setminus J)$ , the signed graph  $(G, \Gamma)/I \setminus J$  represents  $(\text{cycle}(G), \Gamma)/I \setminus J$ . We also have the following relationship:

**Remark 37** ([13], page 21). *Take a signed graph  $(G, \Gamma)$  and disjoint  $I, J \subseteq E(G)$ . Then  $\text{ecycle}(G, \Gamma)/I \setminus J = \text{ecycle}((G, \Gamma)/I \setminus J)$ .*

Take vertices  $u, v$  of  $G$ . We say  $u$  and  $v$  form a *blocking pair* if every non-loop odd circuit of  $(G, \Gamma)$  uses either  $u$  or  $v$ . We see from Proposition 2 that  $u$  and  $v$  form a blocking pair if, and only if, there is a signature contained in  $\delta_G(u) \cup \delta_G(v) \cup \text{loops}(G)$ .

**Proposition 38.** *Let  $(G, \Gamma)$  be a signed graph with an odd loop and without a blocking vertex, and let  $N := \text{ecycle}(G, \Gamma)$ . If  $e$  is a non-loop edge of  $G$  such that*

- $|E(\text{si}(N/e))| \geq 8$ ,
- $\text{si}(N/e)$  is internally 4-connected,
- $\text{si}(N/e)$  is graphic,

*then the ends of  $e$  form a blocking pair.*

*Proof.* Let  $(G', \Gamma') := (G, \Gamma)/e$ . By Remark 37,  $N/e = \text{ecycle}(G', \Gamma')$ . Notice that  $(G', \Gamma')$  also has an odd loop. Therefore, as  $\text{si}(N/e)$  is internally 4-connected and  $|E(\text{si}(N/e))| \geq 8$ , it follows from Proposition 33 (d) that  $G' \setminus \text{loops}(G')$  is 3-connected. Since  $\text{si}(N/e)$  is graphic, so is  $N/e$ , and so  $\text{ecycle}(G', \Gamma')$  is graphic. Putting these together, we get from Proposition 36 that  $(G', \Gamma')$  has a blocking vertex  $w$ , that is, every non-loop odd circuit of  $(G', \Gamma')$  uses  $w$ . Since  $(G, \Gamma)$  does not have a blocking vertex,  $w$  is the vertex in  $G' = G/e$  obtained after identifying the ends of  $e$  in  $G$ . Thus, every non-loop odd circuit of  $(G, \Gamma)$  uses one of the ends of  $e$ , implying that the ends of  $e$  form a blocking pair of  $(G, \Gamma)$ , as required.  $\square$

## 7. PROOF OF PROPOSITION 13

Suppose (h1)-(h2) hold and there are distinct elements  $e_1, e_2, e_3, e_4$  of  $M$  such that, for each  $i \in [4]$ ,  $\text{si}(M/e_i)$  is internally 4-connected and is the cycle matroid of a graph where the three edges  $\Lambda(e_i)$  are incident to the same

vertex. Assuming  $|E(M)| \geq 12$ , we need to show  $M$  is one deletion away from being graphic. Recall that, for each  $i \in [4]$ ,  $M/e_i$  is a loopless matroid with exactly three non-trivial parallel classes, and these classes have cardinality two, so  $|E(\text{si}(M/e_i))| = |E(M)| - 4 \geq 12 - 4$ :

**1.** For each  $i \in [4]$ ,  $|E(\text{si}(M/e_i))| \geq 8$ .

By (h1),  $M$  is simple, so we may assume that  $\{e_1, e_2, e_3\}$  is not a triangle of  $M$ . Denote the three triangles of  $M$  containing  $e_1$  by  $\{e_1, f, f'\}$ ,  $\{e_1, g, g'\}$ ,  $\{e_1, h, h'\}$  where  $\Lambda(e_1) = \{f, g, h\}$ ; the existence of these triangles is guaranteed by (h2). Recall that the non-trivial parallel classes of  $M/e_1$  are  $\{f, f'\}$ ,  $\{g, g'\}$ ,  $\{h, h'\}$ . As it is the case for  $\text{si}(M/e_1)$ , we know that  $M/e_1$  also is the cycle matroid of a graph  $H$  where  $f, g, h$  are incident to the same vertex, say  $v \in V(H)$ . Notice that  $H$  is a loopless graph with exactly three non-trivial parallel classes  $\{f, f'\}$ ,  $\{g, g'\}$ ,  $\{h, h'\}$ . In particular,  $v$  is the only vertex common to any two of  $f, g, h$ . Let  $\Gamma$  be a cocycle of  $M$  that contains  $e_1$ . By Remark 31,  $M = \text{ecycle}(H + e_1, \Gamma)$ . Clearly  $e_1$  is an odd loop of  $(H + e_1, \Gamma)$ , and therefore,  $\{f, f'\}$ ,  $\{g, g'\}$ ,  $\{h, h'\}$  are odd circuits of this signed graph. As a result, if  $(H + e_1, \Gamma)$  has a blocking vertex, then  $v$  must be the one, and if it has a blocking pair, then  $v$  must belong to the pair. If  $v$  is a blocking vertex, then  $M$  is graphic by Remark 35, and we are done. Otherwise,

**2.**  $(H + e_1, \Gamma)$  does not have a blocking vertex.

Consider the two edges  $e_2, e_3$  of  $H$ . Since  $\{e_1, e_2, e_3\}$  is not a triangle of  $M$ , edges  $e_2, e_3$  are not parallel. Since (1) and (2) hold, we may use Proposition 38 to conclude that, for  $j \in \{2, 3\}$ , the ends of  $e_j$  form a blocking pair of  $(H + e_1, \Gamma)$ . In particular,  $e_2 \cap e_3 = \{v\}$ . Write  $e_2 = \{v, u\}$  and  $e_3 = \{v, w\}$ .

**3.**  $H \setminus \{u, v, w\}$  is connected.

*Subproof.* Let  $H' := H/e_2$  and  $(H' + e_1, \Gamma') := (H + e_1, \Gamma)/e_2$ . By Remark 37,  $\text{ecycle}(H' + e_1, \Gamma') = M/e_2$ . Since (1) holds, we may use Proposition 33 (d) to conclude that  $H' \setminus \text{loops}(H')$  is 3-connected. In particular, if  $uv$  is the vertex of  $H'$  corresponding to the ends of  $e_2$ , the graph  $H' \setminus \text{loops}(H') \setminus \{uv, w\}$  is connected. As a result,  $H \setminus \{u, v, w\}$  is connected.  $\diamond$

Since  $\{v, u\}$  and  $\{v, w\}$  are blocking pairs, every non-loop odd circuit of  $(H + e_1, \Gamma)$  uses either  $v$  or both  $u, w$ . As the non-trivial parallel classes of  $H$  are incident with  $v$ , there is at most one edge with ends  $u, w$ .

**4.**  $H$  has an edge  $\Omega$  with ends  $u, w$ , and every non-loop odd circuit of  $(H + e_1, \Gamma)$  uses either  $v$  or the edge  $\Omega$ .

*Subproof.* Let  $C$  be a non-loop odd circuit  $C$  of  $(H + e_1, \Gamma)$  such that  $v \notin V(C)$ . Then  $\{u, w\} \subseteq V(C)$ . It suffices to show that  $C$  contains an edge whose ends are  $u$  and  $w$ . Suppose otherwise. Let  $x, y$  be the two neighbors of  $u$  in  $H[C]$  – note  $x, y \in V(H) - \{u, v, w\}$  by our contrary assumption. Thus by (3), there is an  $xy$ -path  $P \subseteq E(H)$  that is disjoint from  $\{u, v, w\}$ . Consider the two cycles  $C_1 := \{u, x\} \cup \{u, y\} \cup P$  and  $C_2 := C \Delta C_1$ . Since  $C_1$  is disjoint from the blocking pair  $\{v, w\}$ , and  $C_2$  is disjoint from the blocking pair  $\{v, u\}$ , it follows that both  $C_1, C_2$  are even in  $(H + e_1, \Gamma)$ , implying in turn that  $C = C_1 \Delta C_2$  is also even in  $(H + e_1, \Gamma)$ , contradicting our choice of  $C$ .  $\diamond$



Therefore,  $(H + e_1, \Gamma) \setminus \Omega$  has  $v$  as a blocking vertex. By Remark 37, we have  $M \setminus \Omega = \text{ecycle}((H + e_1, \Gamma) \setminus \Omega)$ , so it follows from Remark 35 that  $M \setminus \Omega$  is graphic, as required.  $\square$

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